

# BACKTRACKING

- Forward chaining goes from axioms forward into goals.
- Backward chaining goes from goals and works backward to prove them with existing axioms.

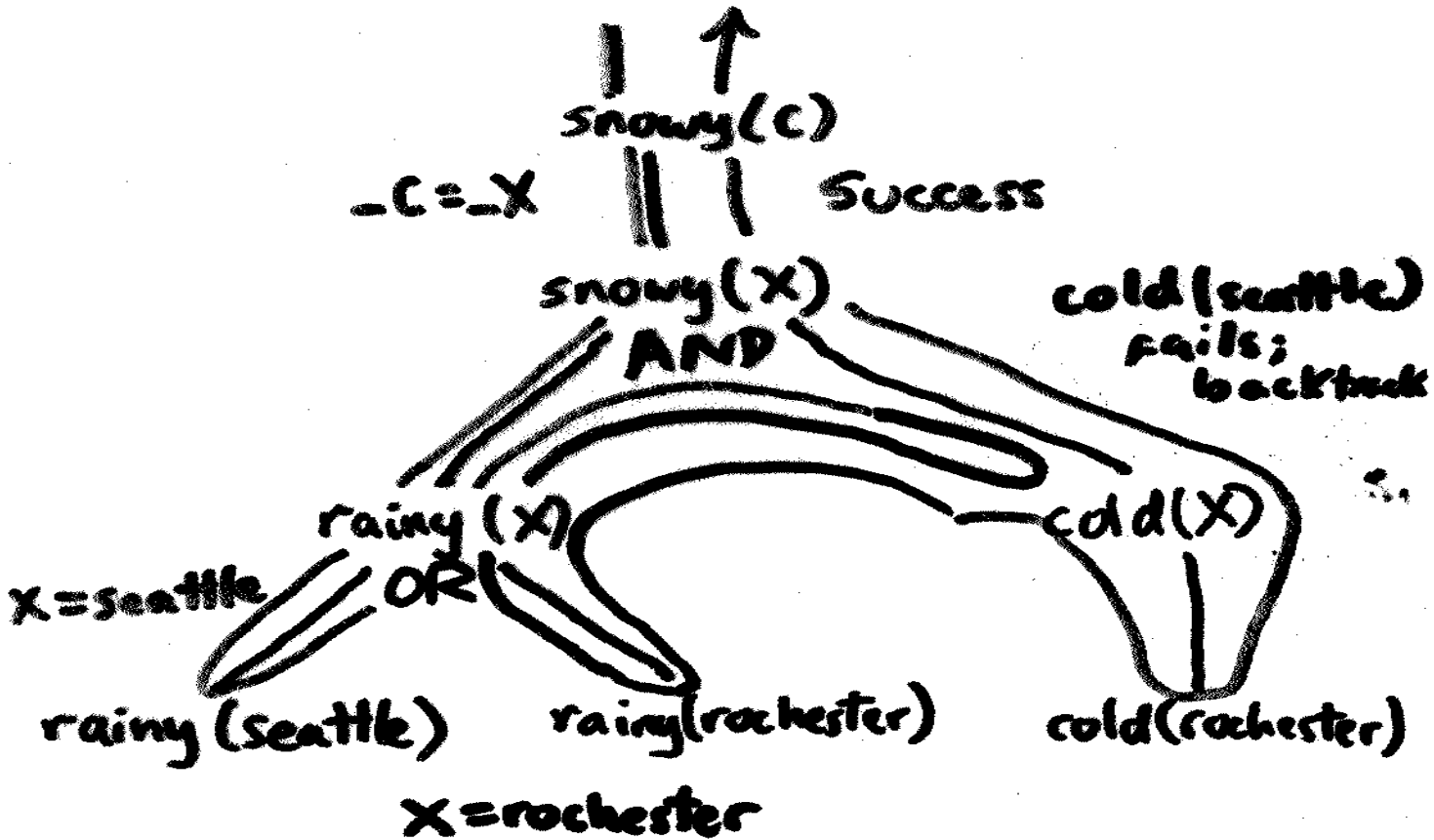
# BACKTRACKING

rainy(seattle).

rainy(rochester).

cold(rochester).

snowy(x) :- rainy(x), cold(x).



# IMPERATIVE CONTROL FLOW

Programmer has explicit control on backtracking process.

## CUT (!)

- As a goal it succeeds, but with side-effect:

Commits interpreter to choices made since unifying parent goal with left-hand side of current rule.

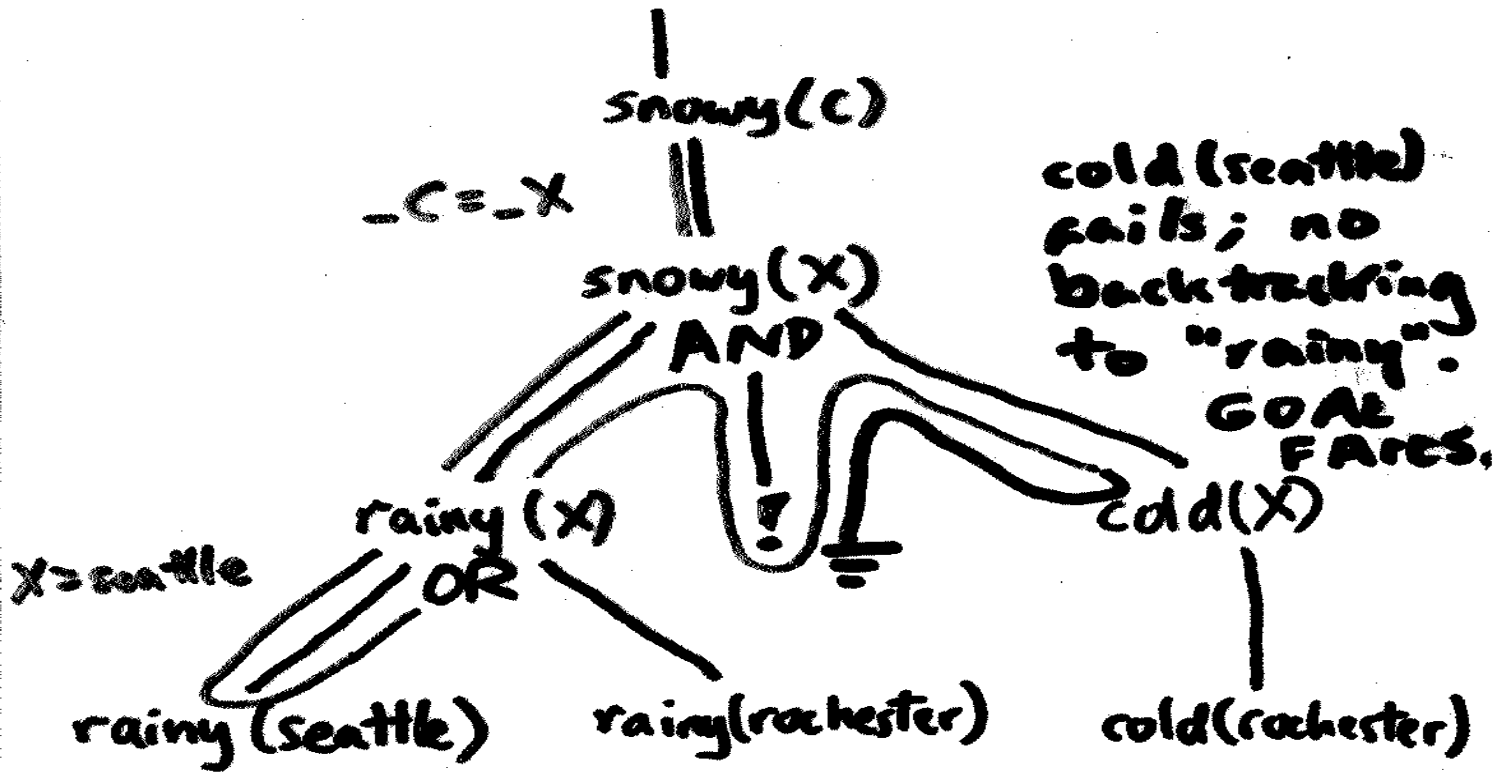
# BACKTRACKING

rainy(seattle).

rainy(rochester).

cold(rochester).

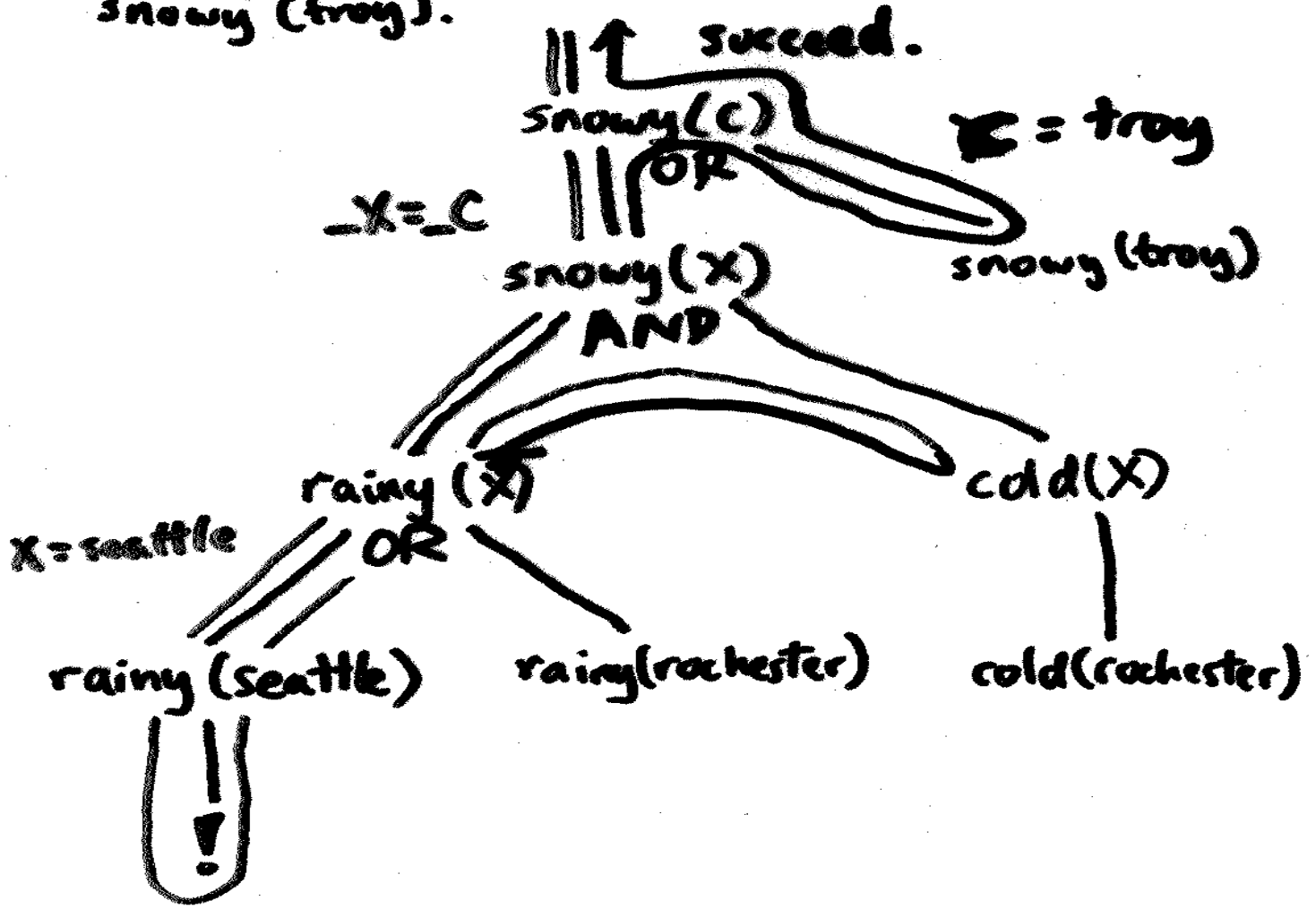
snowy(x) :- rainy(x), cold(x).





# BACKTRACKING

rainy (seattle):-!.  
rainy (rochester).  
cold (rochester).  
snowy (X) :- rainy (X), cold (X).  
snowy (troy).



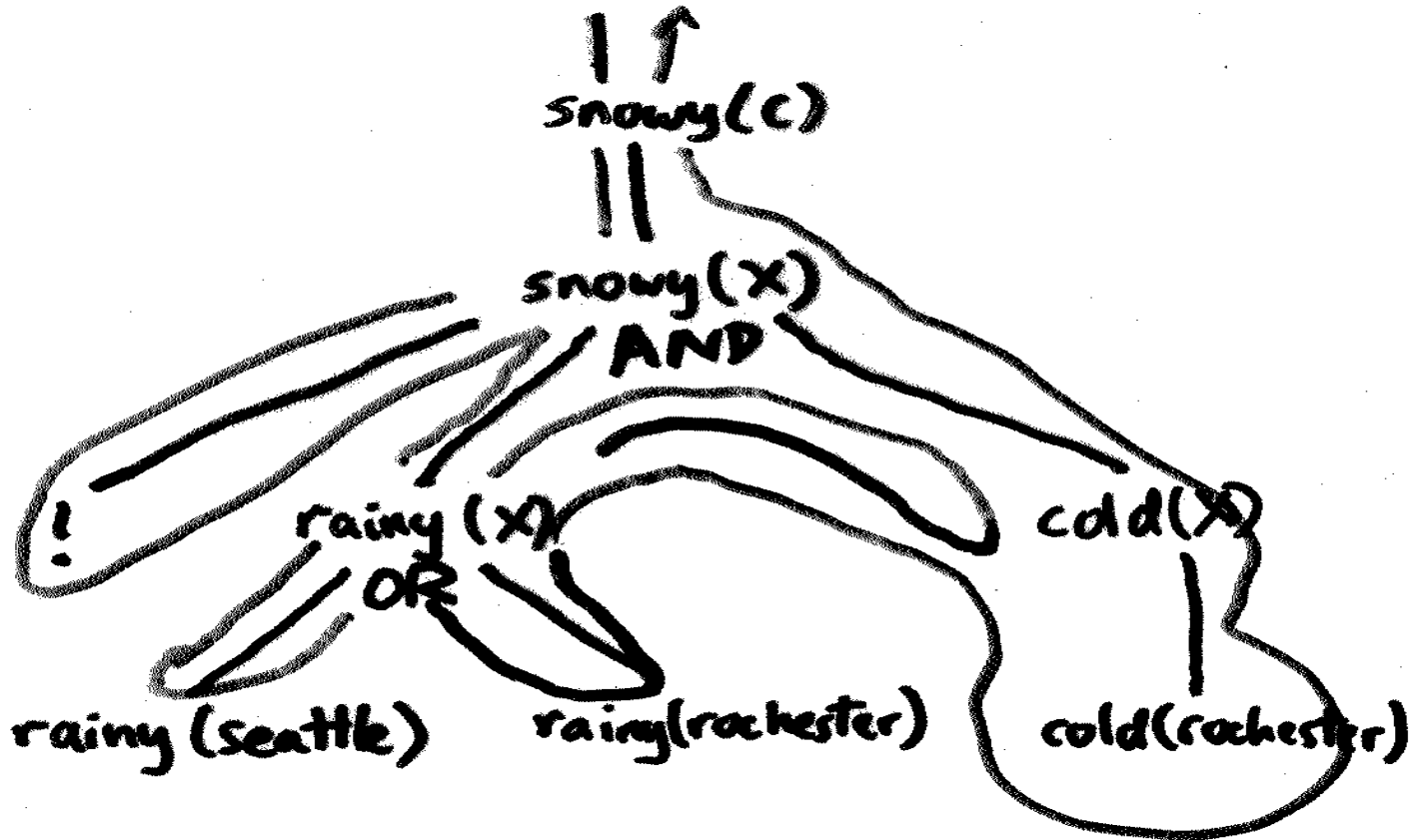
# BACKTRACKING

rainy(seattle).

rainy(rochester).

cold(rochester).

snowy(x) :- rainy(x), cold(x).



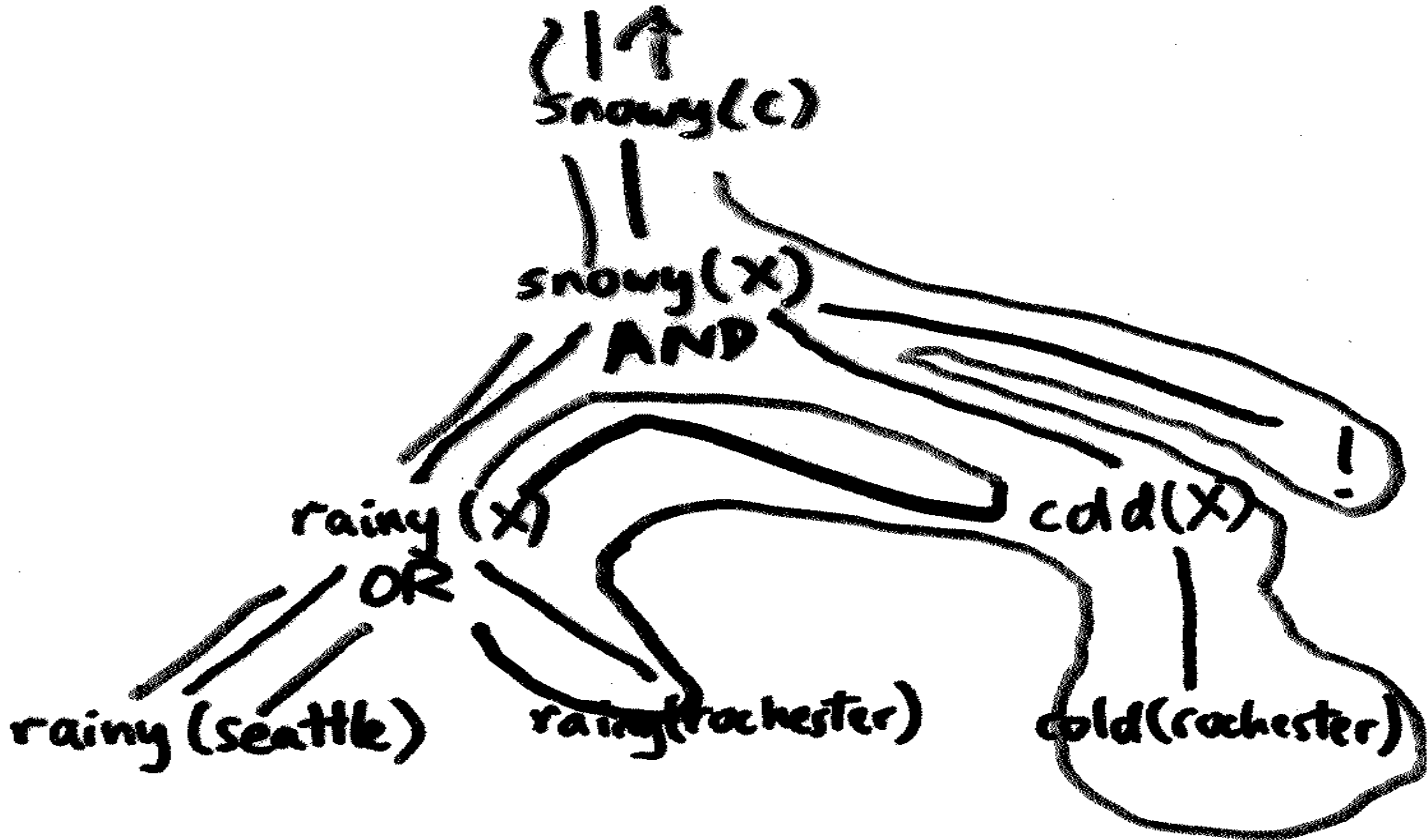
# BACKTRACKING

rainy(seattle).

rainy(rochester).

cold(rochester).

snowy(x) :- rainy(x), cold(x), !.





# FIRST-CLASS TERMS

call(P)

invoke predicate  
as a goal.

assert(P)

adds predicate to  
database

retract(P)

removes predicate  
from database

functor(T, F, A)

succeeds if T is  
a term w/ functor F  
and arity A.

not  $P$  is not  $\neg P$

- In Prolog, the database of facts and rules includes a list of things assumed to be true.
- It does not include anything assumed to be false.
- Unless our database contains everything that is true (the closed world assumption), the goal not  $P$  can succeed simply because our current knowledge is insufficient to prove  $P$ .

# NOT SEMANTICS

$\text{not}(P) :- \text{call}(P), !, \text{fail}.$   
 $\text{not}(P).$

Definition of not in terms of failure (fail) means that variable bindings are lost whenever not succeeds, e.g.:

? - not(not(snowy(x))).

X = \_G147

# MORE NOT VS $\neg$

? - snowy(x).

x = rochester

? - not(snowy(x)).

no

// it does NOT  
reply:

x = seattle

The meaning of not(snowy(x)) is:

$\neg \exists x [\text{snowy}(x)]$

rather than:

$\exists x [\neg \text{snowy}(x)]$

not(not(snowy(x)))

call(not(snowy(x)))

not(snowy(x))

!

fail

call(snowy(x))

snowy(x)

!

fail

x = rochester

true

# fail, true, ...

fail

fail current goal.

true

always succeed.

repeat

always succeed, provides  
infinite choice points

repeat.

repeat :- repeat.

**Exercise:**

what do the following queries do?

? - repeat.

? - repeat, true.

? - repeat, fail.

# NOT

$\text{not}(P) :- \text{call}(P), !, \text{fail}.$   
 $\text{not}(P).$

## IF-THEN ( $\rightarrow$ )

$\rightarrow(\text{If}, \text{Then}, \text{Else})$

$:- \text{If}, !, \text{Then}.$

$\rightarrow(\text{If}, \text{Then}, \text{Else}) :- \text{Else}.$

$\rightarrow(\text{If}, \text{Then}, \text{Else})$

$:- \text{If}, \text{Then}, !$

# CONDITIONALS AND LOOPS

statement :- condition, !, then-part.

statement :- else-part.

natural (1).  
natural (N) :- natural (M),  
N is M+1.

my-loop (N) :- natural (I), I <= N,  
write (I), nl,  
I = N, // test-  
!, fail. // cut-fail.

Also called generate-and-test.