Declarative Computation Model

Kernel language semantics
Basic concepts, the abstract machine (VRH 2.4.1-2.4.2)

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October 9, 2007

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Sequential declarative computation model

- The single assignment store
  - declarative (dataflow) variables
  - partial values (variables and values are also called *entities*)
- The kernel language syntax
- The kernel language semantics
  - The environment: maps textual variable names (variable identifiers) into entities in the store
  - Interpretation (execution) of the kernel language elements (statements) by the use of an abstract machine
  - Abstract machine consists of an execution stack of statements transforming the store

Kernel language syntax

The following defines the syntax of a statement, ⟨s⟩ denotes a statement

⟨s⟩ ::=	skip	*empty statement*
|	⟨x⟩ = ⟨y⟩	*variable-variable binding*
|	⟨x⟩ = ⟨v⟩	*variable-value binding*
|	⟨s_1⟩ ⟨s_2⟩	*sequential composition*
|	local ⟨x⟩ in ⟨s_1⟩ end	*declaration*
|	if ⟨x⟩ then ⟨s_1⟩ else ⟨s_2⟩ end	*conditional*
|	{ ⟨x⟩ ⟨y_1⟩ … ⟨y_n⟩ }	*procedural application*
|	case ⟨x⟩ of ⟨pattern⟩ then ⟨s_1⟩ else ⟨s_2⟩ end	*pattern matching*
⟨v⟩ ::=	proc { $ ⟨y_1⟩ … ⟨y_n⟩ } ⟨s_1⟩ end | …	*value expression*
⟨pattern⟩ ::=	...	

Examples

- `local X in X = 1 end`
-
```
local X Y T Z in
   X = 5
   Y = 10
   T = (X>=Y)
   if T then Z = X else Z = Y end
   {Browse Z}
end
```
-
```
local S T in
   S = proc {$ X Y} Y = X*X end
   {S 5 T}
   {Browse T}
end
```


Procedure abstraction

- Any statement can be abstracted to a procedure by selecting a number of the 'free' variable identifiers and enclosing the statement into a procedure with the identifiers as parameters
- `if X >= Y then Z = X else Z = Y end`
- Abstracting over all variables
```
proc {Max X Y Z}
   if X >= Y then Z = X else Z = Y end
end
```
- Abstracting over X and Z
```
proc {LowerBound X Z}
   if X >= Y then Z = X else Z = Y end
end
```


Computations (abstract machine)

- A computation defines how the execution state is transformed step by step from the initial state to the final state
- A *single assignment store* σ is a set of store variables, a variable may be unbound, bound to a partial value, or bound to a group of other variables
- An *environment E* is mapping from variable identifiers to variables or values in σ, e.g. $\{X \rightarrow x_1, Y \rightarrow x_2\}$
- A *semantic statement* is a pair
  (⟨s⟩ , *E*) where ⟨s⟩ is a statement
- *ST* is a stack of <u>semantic statements</u>

Computations (abstract machine)

- A computation defines how the execution state is transformed step by step from the initial state to the final state
- The *execution state* is a pair
  (ST , σ)
- ST is a stack of <u>semantic statements</u>
- A *computation* is a sequence of execution states
  $(ST_0 , \sigma_0) \rightarrow (ST_1 , \sigma_1) \rightarrow (ST_2 , \sigma_2) \rightarrow \ldots$

Semantics

- To execute a program (i.e., a statement) $\langle s \rangle$ the initial execution state is
  $([(\langle s \rangle , \varnothing)] , \varnothing)$
- ST has a single semantic statement $(\langle s \rangle , \varnothing)$
- The environment E is empty, and the store σ is empty
- [...] denotes the stack
- At each step the first element of ST is popped and execution proceeds according to the form of the element
- The final execution state (if any) is a state in which ST is empty

skip

- The semantic statement is
  (skip, E)
- Continue to next execution step

skip

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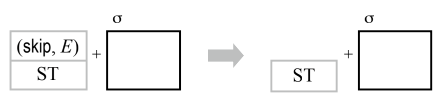
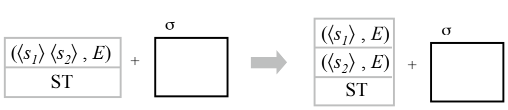
  (skip, E)
- Continue to next execution step



Sequential composition

- The semantic statement is
  $(\langle s_1 \rangle \langle s_2 \rangle , E)$
- Push $(\langle s_2 \rangle , E)$ and then push $(\langle s_1 \rangle , E)$ on ST
- Continue to next execution step



Calculating with environments

- E is mapping from identifiers to entities (both store variables and values) in the store
- The notation $E(\langle y \rangle)$ retrieves the entity x associated with the identifier $\langle y \rangle$ from the store
- The notation $E + \{\langle y \rangle_1 \rightarrow x_1, \langle y \rangle_2 \rightarrow x_2, \ldots , \langle y \rangle_n \rightarrow x_n \}$
  - denotes a new environment E' constructed from E by adding the mappings
    $\{\langle y \rangle_1 \rightarrow x_1, \langle y \rangle_2 \rightarrow x_2, \ldots , \langle y \rangle_n \rightarrow x_n\}$
  - $E'(\langle z \rangle)$ is x_k if $\langle z \rangle$ is equal to $\langle y \rangle_k$, otherwise $E'(\langle z \rangle)$ is equal to $E(\langle z \rangle)$
- The notation $E|_{\{\langle y \rangle_1, \langle y \rangle_2, \ldots , \langle y \rangle_n\}}$ denotes the projection of E onto the set $\{\langle y \rangle_1, \langle y \rangle_2, \ldots , \langle y \rangle_n\}$, i.e., E restricted to the members of the set

Calculating with environments (2)

- $E = \{\mathsf{X} \rightarrow 1, \mathsf{Y} \rightarrow [2\ 3], \mathsf{Z} \rightarrow x_i\}$
- $E' = E + \{\mathsf{X} \rightarrow 2\}$
- $E'(\mathsf{X}) = 2$,
  $E(\mathsf{X}) = 1$
- $E|_{\{X,Y\}}$ restricts E to the 'domain' {X,Y},
  i.e., it is equal to $\{\mathsf{X} \rightarrow 1, \mathsf{Y} \rightarrow [2\ 3]\}$

Calculating with environments (3)

```
local X in
    X = 1               (E)
    local X in
        X = 2           (E')
        {Browse X}
    end                 (E)
    {Browse X}
end
```


Lexical scoping

- Free and bound identifier occurrences
- An identifier occurrence is *bound* with respect to a statement ⟨s⟩ if it is in the scope of a declaration inside ⟨s⟩
- A variable identifier is declared either by a 'local' statement, as a parameter of a procedure, or implicitly declared by a case statement
- An identifier occurrence is *free* otherwise
- In a running program every identifier is bound (i.e., declared)

Lexical scoping (2)

```
proc {P X}
    local Y in Y = 1 {Browse Y} end
    X = Y
end
```

Free Occurrences   Bound Occurrences

Lexical scoping (3)

```
local Arg1 Arg2 in
    Arg1 = 111*111
    Arg2 = 999*999
    Res = Arg1*Arg2
end
```

Free Occurrences   Bound Occurrences

This is not a runnable program!

Lexical scoping (4)

```
local Res in
    local Arg1 Arg2 in
        Arg1 = 111*111
        Arg2 = 999*999
        Res = Arg1*Arg2
    end
    {Browse Res}
end
```

Lexical scoping (5)

```
local P Q in
   proc {P}  {Q} end
   proc {Q} {Browse hello} end
   local Q in
       proc {Q} {Browse hi} end
       {P}
   end
end
```


Exercises

46. Translate the following function to the kernel language:

```
fun {AddList L1 L2}
  case L1 of H1|T1 then
    case L2 of H2|T2 then
        H1+H2|{AddList T1 T2}
    end
  else nil end
end
```

47. Translate the following function call to the kernel language:

```
{Browse {Max 5 7}}
```


Exercises

48. Explain the difference between static scoping and dynamic scoping. Give an example program that produces different results with static and dynamic scoping.
49. *Think of a reason why static scoping may be preferable to dynamic scoping. Think of a reason why dynamic scoping may be preferable to static scoping.