

REFERENCE CELL IN PICT

```
new contents: ^Int
```

```
run contents!0
```

```
def set [v: Int c: Sig] =
```

```
  contents?_ =
```

```
  (contents!v | c![])
```

```
def get [res: ^Int] =
```

```
  contents?v =
```

```
  (contents!v | res!v)
```

USING PICT'S REFERENCE CELL

```
new done: ^[]  
new res: ^Int
```

```
run ( set! [S done] (chan done) )
```

```
  | done? [] =
```

```
    (get! [(chan res)])
```

```
  | res? i = print! i))
```

5

```
def res i: Int = print! i
```

```
def done [] = {get! [res]}
```

```
run set! [S done]
```

5

CONTINUATION-PASSING TRANSLATION

```
run (def f[x:Int res:/Int] = t![x x res]  
     y!f)
```

we create a channel (responsive) f ,
and we send it along y .

```
(new n:T x!n)
```

creates a channel n and sends it
along x . It is the same as:

```
x!(new n:T n)
```

Note it does not send the expression, it
evaluates the expression and sends the
value.

CPS - Continued

e.g.

$x! [23 \text{ (new } x:A \ x) \text{ (new } y:B \ y)]$

creates two new channels, packages into a single type (along with integer 23) and sends the result along x .

VALUE DECLARATIONS

e.g. $(\text{val } x = (\text{new } n:T [n \ n]) \ e)$

binds x to the result of executing $(\text{new } n:T [n \ n])$ and then executes e .

Note e blocks until the bindings in val expression have been established.

APPLICATION SYNTAX

$(v \ v_1 \ \dots \ v_n)$

e.g.

```
def double [s:String r:(String)]  
  = +$! [s s r]
```

we can write `(double s)` as
a value dropping explicit result channel `r`

```
run print!(double "soothe")
```

soothe soothe

FUNCTIONAL PROGRAMMING STYLE

`def f[a1:A1 a2:A2 a3:A3 r:T] = r!v`

can be replaced by a "function definition" that avoids explicitly naming `r`:

`def f(a1:A1 a2:A2 a3:A3):T = v`

ANONYMOUS ABSTRACTIONS

`\a`

is the same as:

`(def x a x)`



FOR LOOP EXAMPLE

```
def for (min: Int max: Int f: ! [Int / ()] d: / ())
```

```
  (def loop x: Int =
```

```
    if (<= x max) then
```

```
      (new c: ^ [])
```

```
      ( f! [x (chan c)]
```

```
      | c? [] = loop! (+ x 1) ))
```

```
    else
```

```
      done! []
```

```
  loop! min )
```

```
run (new done: ^ [])
```

```
  (for! [1 4
```

```
    \ [x c] = (print! x | c! [])
```

```
    (chan done)]
```

```
  | done? [] = print! "Done!") )
```

1
2
3
4
Done!

SEQUENCING

```
run  
( val [] = (pr "hello")  
  val [] = (pr "world")  
  () )
```

equivalent to:

```
run ( (pr "hello");  
      (pr "world");  
      () )
```

SEMANTICS OF VALUE DECLARATIONS & SEQUENCING

$$\llbracket (\text{val } p = v \ e) \rrbracket = (\text{def } c \ p = \llbracket e \rrbracket \\ \llbracket v \rightarrow c \rrbracket)$$

$$v; \Rightarrow \text{val } [] = v$$

USING REFERENCE CELL (WITH APP SYNTAX & SEQUENCING)

```
run ((set 5);  
      (prNL (int.toString (get))));
```

```
5  
      (set 5);  
      (set 8);  
      (prNL (int.toString (get))));  
( )
```

REFERENCE CELL IN PICT (REVISITED)

```
type RefInt = [  
  set = / [Int Sig]  
  get = / [!Int]  
]
```

```
def refInt (): RefInt =  
  (new contents: ^Int  
   run contents!0  
   [  
     set = \ [v: Int c: Sig] =  
       contents?_ = (contents!v | c![v])  
     get = \ [res: !Int] =  
       contents?v = (contents!v | res!v)  
   ]  
)
```



REVISITED REFERENCE CELL USAGE

```
val ref1 = (ref 1)
```

```
val ref2 = (ref 1)
```

```
run ((ref2.set 5);
```

```
    (ref1.set 3);
```

```
    (println (int.toString (ref1.get))));
```

```
    (println (int.toString (ref2.get))));
```

```
  )
```

3
5

LISTS

```
import "std/List"  
val l = (cons 6 (cons 7 (cons 8 nil)))  
run print! (car (cdr l))
```

7

FOLDING

(cons > 6 7 8 nil)

≡

(cons 6 (cons 7 (cons 8 nil)))

(f > a₁ a₂ ... a_n a)

≡

(f a₁ (f a₂ ... (f a_n a)))

RIGHT FOLDING

(f < a a₁ a₂ ... a_n) ≡

(f (f (f a a₁) a₂) ... a_n)

POLY MORPHISM

```
def print2nd [#X l: (List X) p: / [X /String]]  
= if (null l) then print! "Null list"  
  else if (null (cdr l)) then print! "Null tail"  
  else print! (p (car (cdr l)))
```

The # indicates it is a type parameter.

e.g.:

```
run print2nd! [#Int (cons 6 7 8 nil)  
              int.toString]
```

7

```
run print2nd! [#String (cons "one" "two" nil)  
                 \ (s:String) = s ]
```

two

ABSTRACT TYPES

```
val [# Weekday  
    m: Weekday t: Weekday w: Weekday  
    r: Weekday f: Weekday s: Weekday  
    n: Weekday  
    sameday: / [Weekday Weekday / Bool]  
    tomorrow: / [Weekday / Weekday]]
```

```
= [# Int  
   0 1 2 3 4 5 6  
   \ (d1: Int d2: Int) = (== d1 d2)  
   \ (d: Int) = (mod (+ d 1) 7)]
```

Now, we can use the abstract type, e.g.:

```
def weekend (d: Weekday): Bool =  
  (|| (sameday d s) (sameday d n))
```

USER-DEFINED TYPE CONSTRUCTORS

GENERIC REFERENCE CELL CONSTRUCTOR

```
type (Ref X) = [  
  set = / [X sig]  
  get = / [!X]  
]
```

Ref is a parametric type -- it describes a family of types.

$(\text{Ref } T) \equiv [\text{set} = / [T \text{ sig}] \text{ get} = / [!T]]$

GENERIC REFERENCE CELL IN PICT

```
def ref (#X init: X) : (Ref X) =
```

```
(new contents: ^X
```

```
run contents!init
```

```
[ set = \[v: X c: Sig] =
```

```
contents?_ = (contents!v | c![])
```

```
get = \[res: /X] =
```

```
contents?v = (contents!v | res!v)
```

1)

GENERIC REFERENCE CELL USAGE

```
val ref1 = (ref #int 0)
val ref2 = (ref #String "one")
run ((ref1.set 5);
     (prNL (ref2.get)));
    (prNL (int.toString (ref1.get))));
    ()
)
```

one
5

If the type parameter is omitted, PICT will infer it if possible.

e.g. $(\text{ref } \#int\ 0) \equiv (\text{ref } 0)$
 $(\text{ref } \#String\ "one") \equiv (\text{ref } "one")$

NOMADIC PICT

DECLARATIONS

type $T = T'$

new $c : T$ P

agent $a = P$ and ... and $a' = P'$ in Q

migrate to s P

def $f[\dots] = P$ and ... and $f'[\dots] = P'$ Q

type abbreviation
new channel name

agent creation

agent migration

process abstraction

PROCESSES

$(P \mid Q)$

$(D P)$

$()$

parallel composition

local declaration

null process

COMMUNICATION

$c!v$

output v on channel c
in current agent

$c?p = P$

input

$c?^*p = P$

replicated input

if v then P else Q conditional

if local $\langle a \rangle c!v$ then P else Q

test-and-send to agent a
on this site.

$\langle a \rangle c!v$

send to agent a on this side.

$\langle a@s \rangle c!v$

send to agent a on site s .

wait $c?p = P$ timeout $t \rightarrow Q$

input with timeout (secs)

terminate

kill agent

$c@a!v$

location-independent output
to channel c at agent a .

NOMADIC PICT EXAMPLE

getApplet?*[a s] →

agent b =

migrate to s →

(<aes>ack!b|@

in 0

LOCATION-INDEPENDENT COMMUNICATION

<ae?>c!v

c@a!v

MOBILE AGENT EXAMPLE

```
new answer: ^String
```

```
def spawn [s: Site prompt: String] =
```

```
  (agent b =
```

```
    (migrate to s
```

```
      answer@a! (sys.read prompt))
```

```
  in
```

```
    (1)
```

```
  ( spawn! [s1 "How are you?"]
```

```
    | spawn! [s2 "When do we start?"]
```

```
    | answer?# s = print! s
```

...

This code (part of agent a) spawns two agents at sites s1 and s2, and prints answers coming on its "answer" channel.

REFERENCE CELL IN NOMADIC PICT

type RefInt =

```
[ set = / [Agent Int Sig]  
  get = / [Agent /Int]  
]
```

def refInt (s: Site r: /RefInt) =

```
( new set: ^ [Agent Int Sig]  
  new get: ^ [Agent /Int]
```

agent refIntAg =

```
( new contents: ^ Int
```

```
  run contents! 0
```

```
  migrate to s
```

```
( set?# [a: Agent v: Int c: Sig] =  
  contents?_ = (contents!v | c! [I])  
  @a
```

```
  | get?# [a: Agent res: /Int] =  
  contents?v = (contents!v | res@a!v))
```

```
)  
r! [ set = \ [a: Agent v: Int c: Sig] =  
  set @ refIntAg ! [a v c]
```

```
  get = \ [a: Agent res: /Int] = get @ refIntAg [a res]
```

NONADIC PICT REFERENCE CELL USAGE

```
val cell1 = (refInt s1)  
val cell2 = (refInt s2)
```

```
agent a =
```

```
(  
  (cell2.set a 5);  
  (prNL (int.toString (cell1.get a))));  
  (prNL (int.toString (cell2.get a))));  
  ()  
)
```