

Declarative Concurrency

Lazy Execution (VRH 4.5)

Carlos Varela

Rensselaer Polytechnic Institute

Adapted with permission from:

Seif Haridi

KTH

Peter Van Roy

UCL

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Lazy evaluation

- The default functions in Oz are evaluated *eagerly* (as soon as they are called)
- Another way is lazy evaluation where a computation is done only when the result is needed

- Calculates the infinite list:

0 | 1 | 2 | 3 | ...

```
declare  
fun lazy {Ints N}  
  N|{Ints N+1}  
end
```

Lazy evaluation (2)

- Write a function that computes as many rows of Pascal's triangle as needed
- We do not know how many beforehand
- A function is *lazy* if it is evaluated only when its result is needed
- The function `PascalList` is evaluated when needed

```
fun lazy {PascalList Row}
  Row | {PascalList
        {AddList
         Row
         {ShiftRight Row}}}}
end
```

Lazy evaluation (3)

- Lazy evaluation will avoid redoing work if you decide first you need the 10th row and later the 11th row
- The function continues where it left off

```
declare  
L = {PascalList [1]}  
{Browse L}  
{Browse L.1}  
{Browse L.2.1}
```

```
L<Future>  
[1]  
[1 1]
```

Lazy execution

- Without laziness, the execution order of each thread follows textual order, i.e., when a statement comes as the first in a sequence it will execute, whether or not its results are needed later
- This execution scheme is called *eager execution*, or *supply-driven* execution
- Another execution order is that a statement is executed only if its results are needed somewhere in the program
- This scheme is called *lazy evaluation*, or *demand-driven* evaluation (some languages use lazy evaluation by default, e.g., Haskell)

Example

$$B = \{F1 X\}$$

$$C = \{F2 Y\}$$

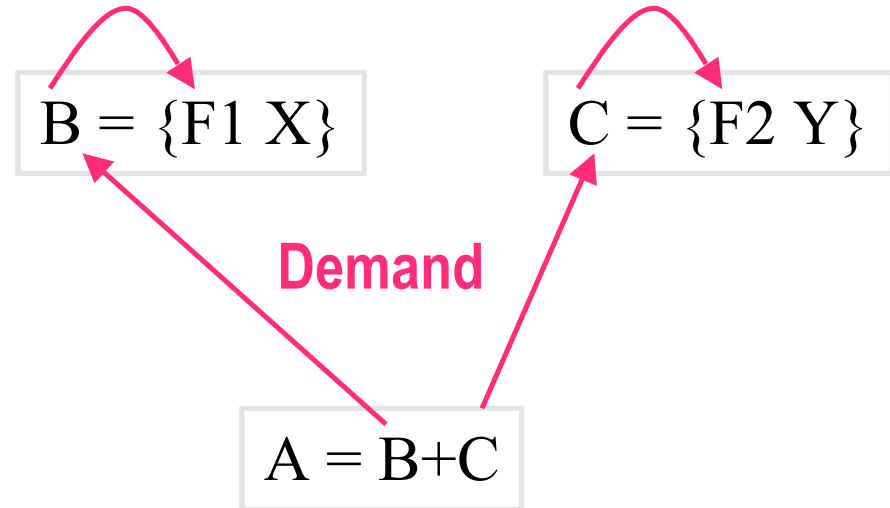
$$D = \{F3 Z\}$$

$$A = B+C$$

- Assume F1, F2 and F3 are lazy functions
- $B = \{F1 X\}$ and $C = \{F2 Y\}$ are executed only if and when their results are needed in $A = B+C$
- $D = \{F3 Z\}$ is not executed since it is not needed

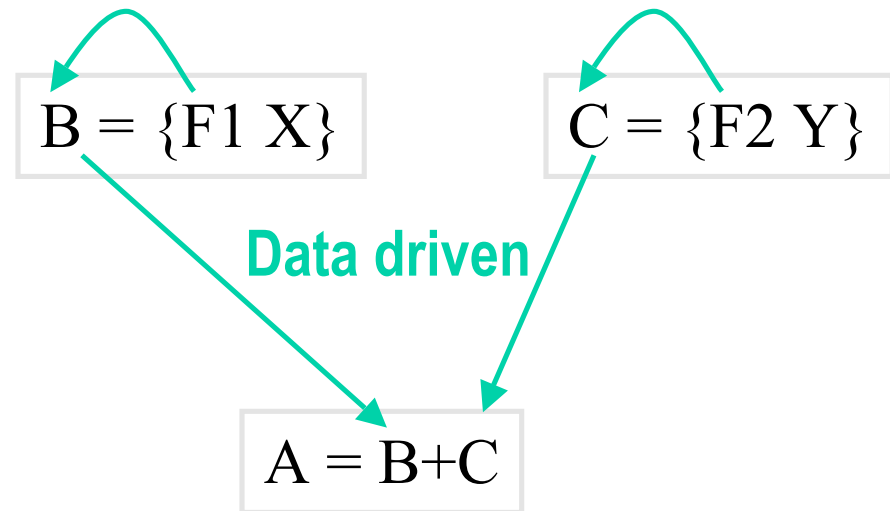
Example

- In lazy execution, an operation suspends until its result is needed
- The suspended operation is triggered when another operation needs the value for its arguments
- In general multiple suspended operations could start concurrently



Example II

- In data-driven execution, an operation suspends until the values of its arguments results are available
- In general the suspended computation could start concurrently



Using Lazy Streams

```
fun {Sum Xs A Limit}  
  if Limit>0 then  
    case Xs of X|Xr then  
      {Sum Xr A+X Limit-1}  
    end  
  else A end  
end
```

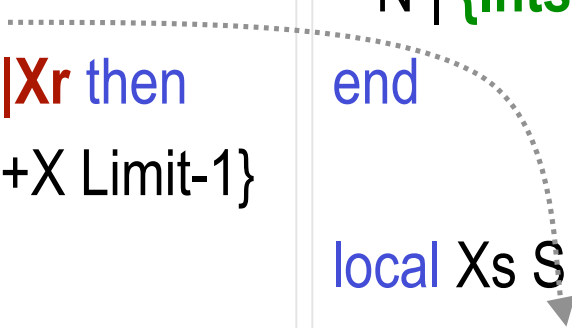
```
local Xs S in  
  Xs={Ints 0}  
  S={Sum Xs 0 1500}  
  {Browse S}  
end
```

How does it work?

```
fun {Sum Xs A Limit}
  if Limit>0 then
    case Xs of X|Xr then
      {Sum Xr A+X Limit-1}
    end
  else A end
end
```

```
fun lazy {Ints N}
  N | {Ints N+1}
end

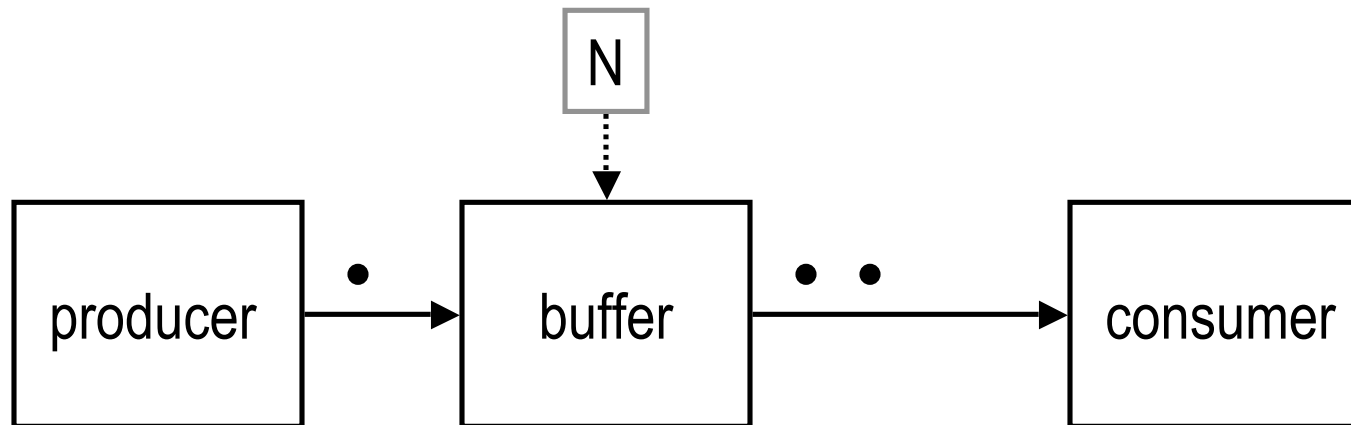
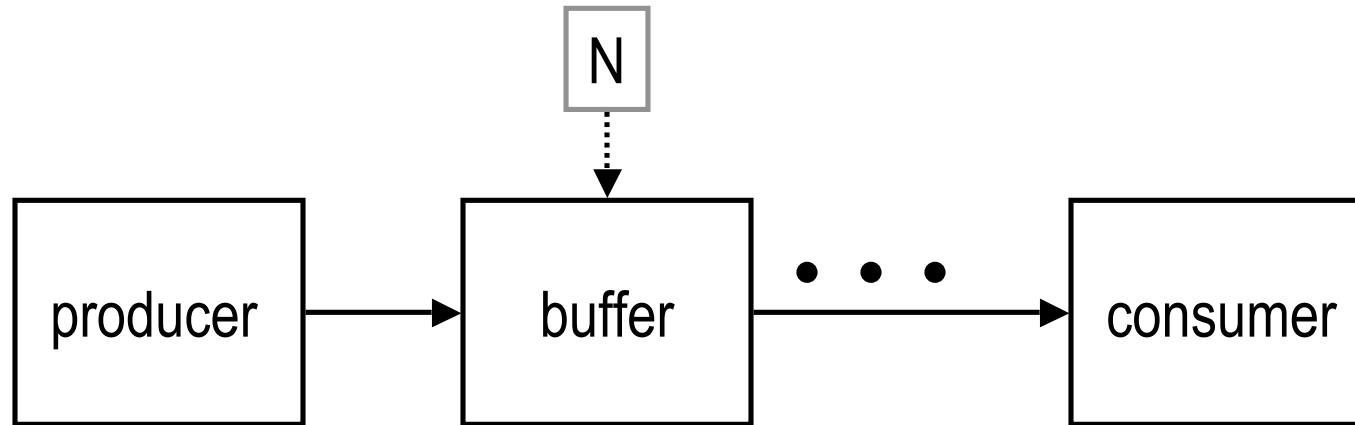
local Xs S in
  Xs = {Ints 0}
  S={Sum Xs 0 1500}
  {Browse S}
end
```



Improving throughput

- Use a lazy buffer
- It takes a lazy input stream `In` and an integer `N`, and returns a lazy output stream `Out`
- When it is first called, it first fills itself with `N` elements by asking the producer
- The buffer now has `N` elements filled
- Whenever the consumer asks for an element, the buffer in turn asks the producer for another element

The buffer example



The buffer

```
fun {Buffer1 In N}  
  End={List.drop In N}  
  
  fun lazy {Loop In End}  
    In.1|{Loop In.2 End.2}  
  end  
  
in  
  {Loop In End}  
end
```

Traversing the In stream,
forces the producer to emit
N elements

The buffer II

```
fun {Buffer2 In N}  
  End = thread  
    {List.drop In N}  
  end  
  fun lazy {Loop In End}  
    In.1|{Loop In.2 End.2}  
  end  
in  
  {Loop In End}  
end
```

Traversing the In stream,
forces the producer to emit
N elements **and at the same
time serves the consumer**

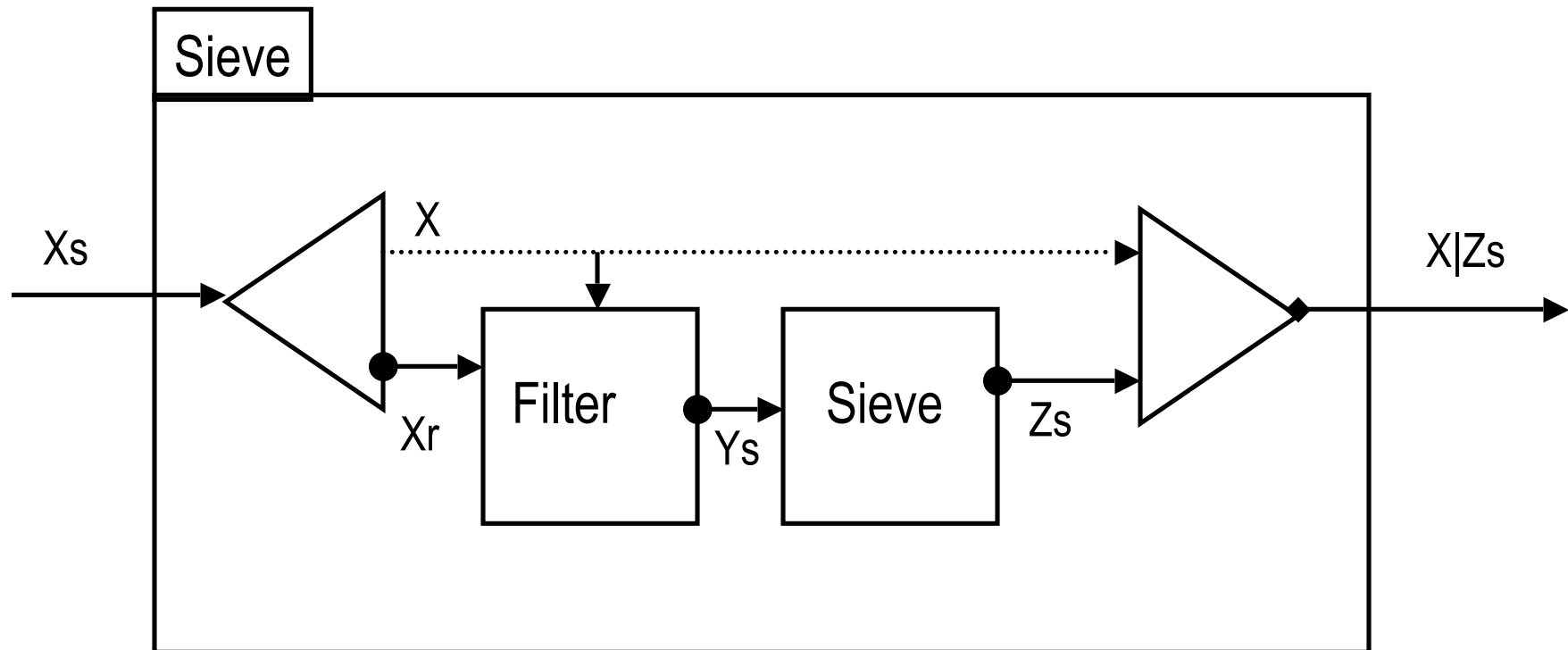
The buffer III

```
fun {Buffer3 In N}  
  End = thread  
    {List.drop In N}  
  end  
  fun lazy {Loop In End}  
    E2 = thread End.2 end  
    In.1|{Loop In.2 E2}  
  end  
in  
  {Loop In End}  
end
```

Traverse the In stream, forces the producer to emit N elements and at the same time serves the consumer, and requests the next element ahead

Larger Example: The Sieve of Eratosthenes

- Produces prime numbers
- It takes a stream $2..N$, peels off 2 from the rest of the stream
- Delivers the rest to the next sieve



Lazy Sieve

```
fun lazy {Sieve Xs}
  X|Xr = Xs in
  X | {Sieve {LFilter
    Xr
    fun {$ Y} Y mod X \= 0 end
  }}
end

fun {Primes} {Sieve {Ints 2}} end
```

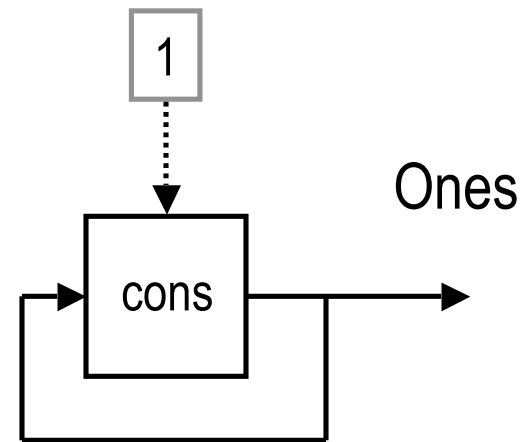
Lazy Filter

For the Sieve program we need a lazy filter

```
fun lazy {LFilter Xs F}  
  case Xs  
  of nil then nil  
  [] X|Xr then  
    if {F X} then X|{LFilter Xr F} else {LFilter Xr F} end  
  end  
end
```

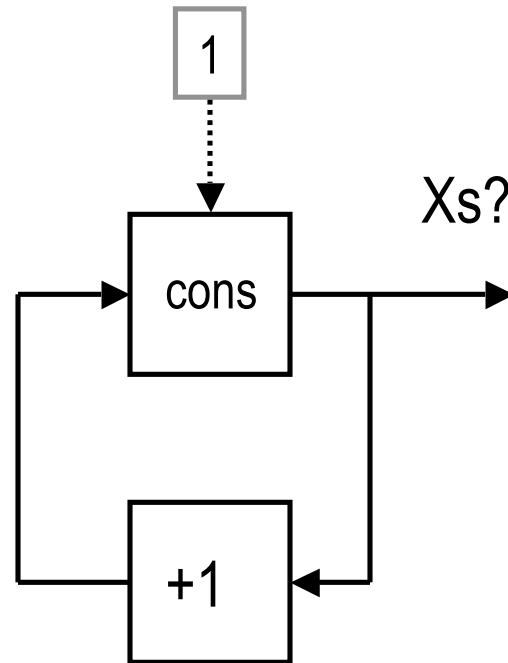
Define streams implicitly

- $\text{Ones} = 1 \mid \text{Ones}$
- Infinite stream of ones



Define streams implicitly

- $Xs = 1 | \{LMap\ Xs$
 `fun {$ X} X+1 end`
- What is Xs ?



The Hamming problem

- Generate the first N elements of stream of integers of the form: $2^a 3^b 5^c$ with $a, b, c \geq 0$ (in ascending order)

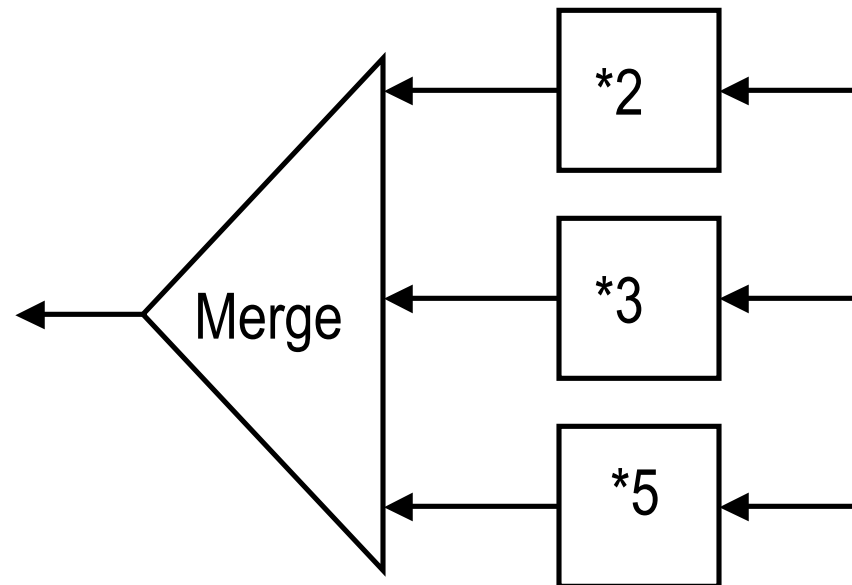
*2

*3

*5

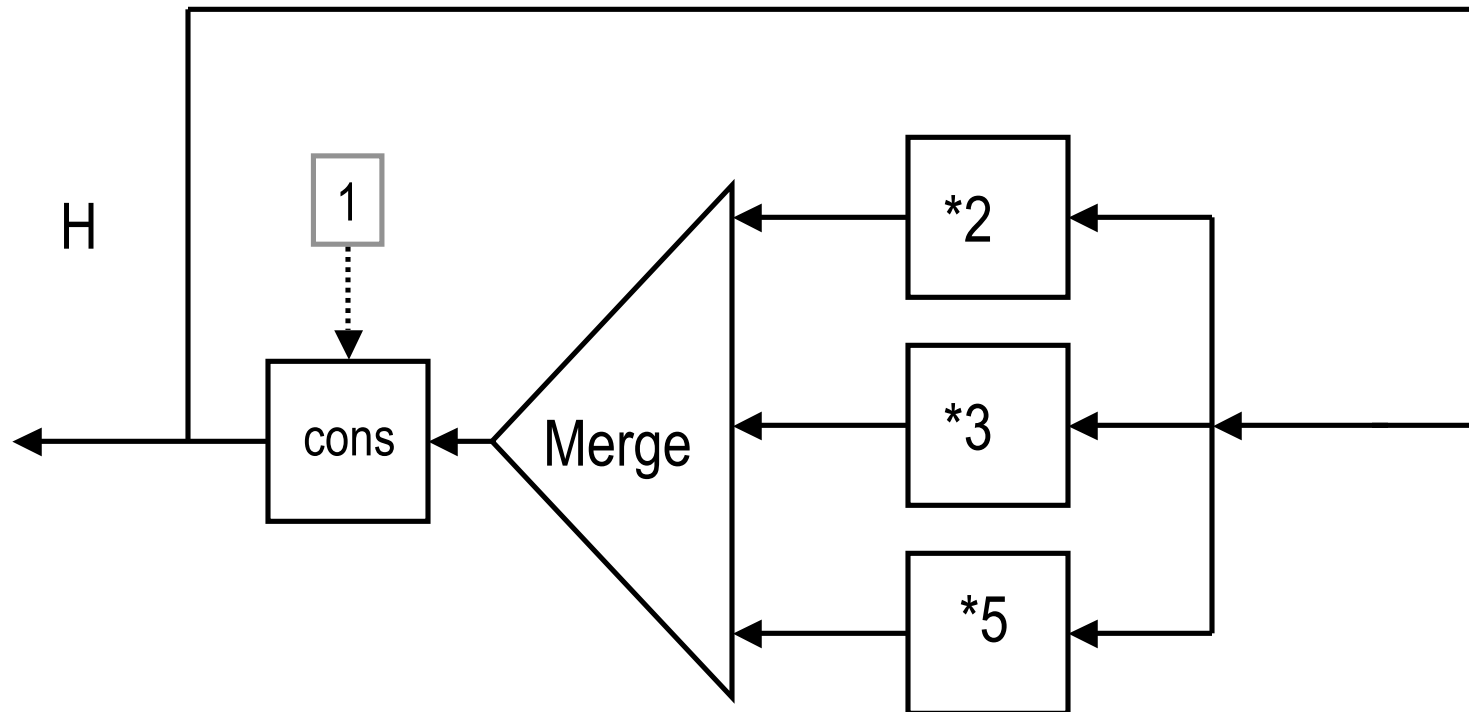
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Lazy File Reading

```
fun {ToList FO}
  fun lazy {LRead} L T in
    if {File.readBlock FO L T} then
      T = {LRead}
    else T = nil {File.close FO} end
  L
end
{LRead}
end
```

- This avoids reading the whole file in memory

List Comprehensions

- Abstraction provided in lazy functional languages that allows writing higher level set-like expressions
- In our context we produce lazy lists instead of sets
- The mathematical set expression
 - $\{x*y \mid 1 \leq x \leq 10, 1 \leq y \leq x\}$
- Equivalent List comprehension expression is
 - $[X*Y \mid X = 1..10 ; Y = 1..X]$
- Example:
 - $[1*1 \ 2*1 \ 2*2 \ 3*1 \ 3*2 \ 3*3 \ \dots \ 10*10]$

List Comprehensions

- The general form is
- $[f(x,y, \dots,z) \mid x \leftarrow \text{gen}(a_1,\dots,a_n) ; \text{guard}(x,\dots)$
 $y \leftarrow \text{gen}(x, a_1,\dots,a_n) ; \text{guard}(y,x,\dots)$

]
- No linguistic support in Mozart/Oz, but can be easily expressed

Example 1

- $z = [x\#x \mid x \leftarrow \text{from}(1,10)]$
- $Z = \{\text{LMap} \{\text{LFrom } 1 \ 10\} \text{ fun}\{\$ X\} X\#X \text{ end}\}$
- $z = [x\#y \mid x \leftarrow \text{from}(1,10), y \leftarrow \text{from}(1,x)]$
- $Z = \{\text{LFlatten}$
 $\{\text{LMap} \{\text{LFrom } 1 \ 10\}$
 $\text{fun}\{\$ X\} \{\text{LMap} \{\text{LFrom } 1 \ X\}$
 $\text{fun} \{\$ Y\} X\#Y \text{ end}$
 $\}$
 end
 $\}$
 $\}$

Example 2

- $z = [x\#y \mid x \leftarrow \text{from}(1,10), y \leftarrow \text{from}(1,x), x+y \leq 10]$
- $Z = \{\mathbf{LFilter}$
 $\{\mathbf{LFlatten}$
 $\{\mathbf{LMap} \{\mathbf{LFrom} 1 10\}$
 $\mathbf{fun} \{ \$ X \} \{\mathbf{LMap} \{\mathbf{LFrom} 1 X\}$
 $\mathbf{fun} \{ \$ Y \} X\#Y \mathbf{end}$
 $\}$
 \mathbf{end}
 $\}$
 $\}$
 $\mathbf{fun} \{ \$ X\#Y \} X+Y \leq 10 \mathbf{end} \}$ }

Implementation of lazy execution

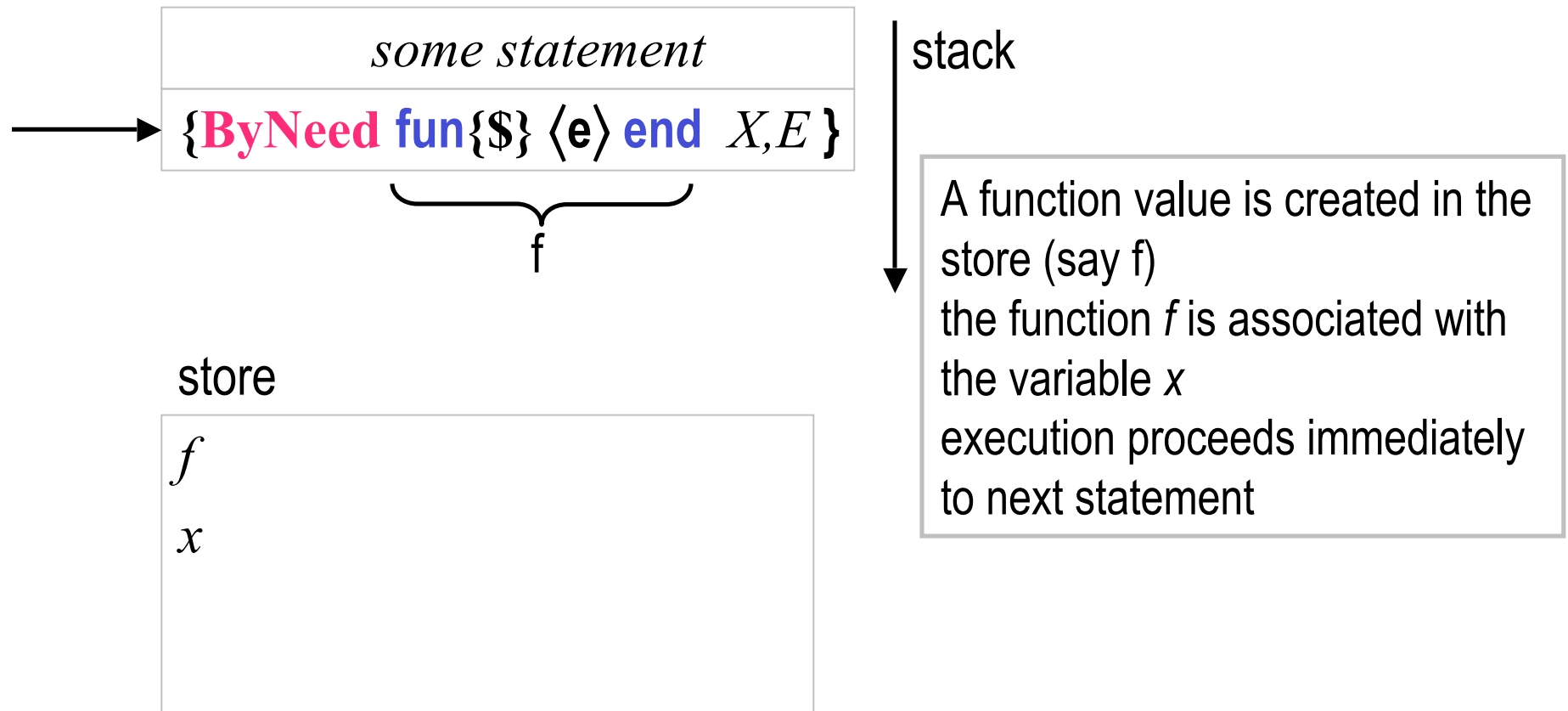
The following defines the syntax of a statement, $\langle s \rangle$ denotes a statement

$\langle s \rangle ::=$

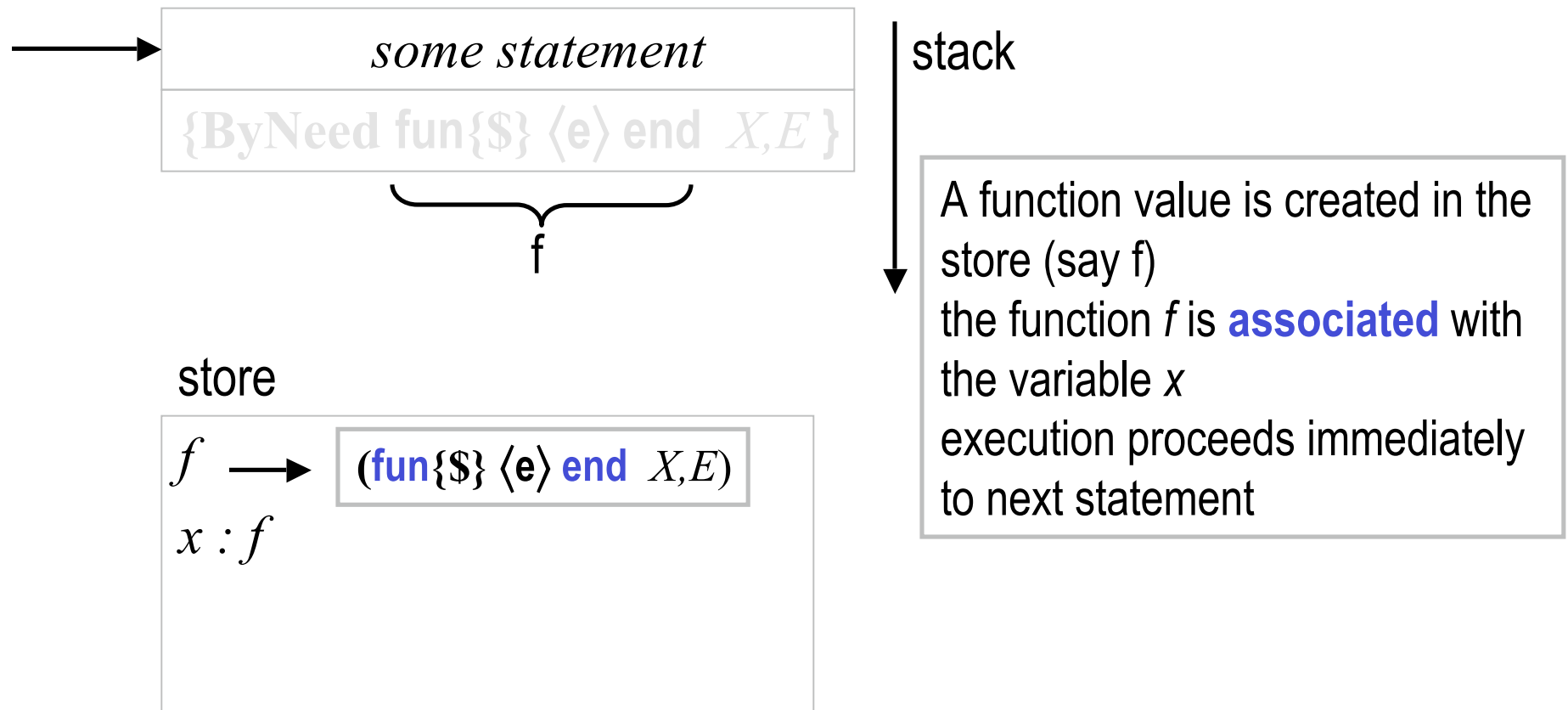
- skip** *empty statement*
- ...
- thread** $\langle s_1 \rangle$ **end** *thread creation*
- {ByNeed fun** $\{ \$ \}$ $\langle e \rangle$ **end** $\langle x \rangle$ *by need statement*

zero arity function variable

Implementation



Implementation



Accessing the ByNeed variable

- $X = \{\text{ByNeed } \text{fun}\{\$ \} \ 111*111 \ \text{end}\}$ (by thread T0)
- Access by some thread T1
 - if $X > 1000$ then $\{\text{Browse hello}\#X\}$ end

or

- $\{\text{Wait } X\}$
- Causes X to be bound to 12321 (i.e. $111*111$)

Implementation

Thread T1

1. X is needed
2. start a thread T2 to execute F (the function)
3. only T2 is allowed to bind X

Thread T2

1. Evaluate $Y = \{F\}$
2. Bind X the value Y
3. Terminate T2

4. Allow access on X

Lazy functions

```
fun lazy {Ints N}  
  N | {Ints N+1}  
end
```



```
fun {Ints N}  
  fun {F} N | {Ints N+1} end  
in {ByNeed F}  
end
```

Exercises

90. Write a lazy append list operation `LazyAppend`. Can you also write `LazyFoldL`? Why or why not?
91. Exercise VRH 4.11.10 (pg 341)
92. Exercise VRH 4.11.13 (pg 342)
93. Exercise VRH 4.11.17 (pg 342)