Concurrent and Distributed Programming Patterns

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Overview

• A motivating application in AstroInformatics
• Programming techniques and patterns
  – farmer-worker computations,
  – iterative computations,
  – peer-to-peer agent networks,
  – soft real-time: priorities, delays
  – causal connections: named tokens, $\text{waitfor}$ property
• Distributed runtime architecture (World-Wide Computer)
  – architecture and implementation
  – distributed garbage collection
• Autonomic computing (Internet Operating System)
  – autonomous migration
  – split and merge
• Distributed systems visualization (OverView)
Milky Way Structure and Evolution

• Principal Investigators:
  H. Newberg (RPI Astronomy),
  M. Magdon-Ismail, B. Szymanski, C. Varela (RPI CS)

• Students:
  M. Newby, M. Arsenault, C. Rice, N. Cole (RPI Astronomy),
  T. Desell, J. Doran (RPI CS)

• Problem Statement:
  What is the structure, origin, and evolution of the Milky Way galaxy?
  How to analyze data from 10,000 square degrees of the north galactic cap collected in five optical filters over five years by the Sloan Digital Sky Survey?

• Applications/Implications:
  Astrophysics: origins and evolution of our galaxy; dark matter distribution.

• Approach:
  Experimental data analysis and simulation
  To use photometric and spectroscopic data for millions of stars to separate and describe components of the Milky Way

• Software:
  MilkyWay@Home BOINC project.
  Generic Maximum Likelihood Evaluation (GMLE) framework.
  N-body Simulations (using CPUs & GPUs.)

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How Do Galaxies Form?

Ben Moore, Inst. Of Theo. Phys., Zurich

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Tidal Streams

- Smaller galaxy gets tidally disrupted by larger galaxy
- Good tracer of galactic potential/dark matter
- Sagittarius Dwarf Galaxy currently being disrupted
- Three other known streams thought to be associated with dwarf galaxies

Kathryn V. Johnston, Wesleyan Univ.
Sloan Digital Sky Survey Data

- **SDSS**
  - ~ 9,600 sq. deg.
  - ~ 287,000,000 objects
  - ~ 10.0 TB (images)

- **SEGUE**
  - ~ 1,200 sq. deg.
  - ~ 57,000,000 objects

- **GAIA (2010-2012)**
  - Over one billion estimated stars

http://www.sdss.org
Maximum Likelihood Evaluation on RPI Grid and BlueGene/L Supercomputer

Computation Time (seconds)

- 2 Minute Evaluation
  - MLE requires 10,000+ Evaluations
  - 15+ Day Runtime

- ~100x Speedup
  - 1.5 Day Runtime

- ~230x Speedup
  - <1 Day Runtime

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February 2010: 1.1Petaflops

April 2010: 1.6Petaflops
Programming Patterns
Farmer Worker Computations

- Most common “Massively Parallel” type of computation
- Workers repeatedly request tasks or jobs from farmer and process them
Farmer Worker Computations

Farmer

Worker 1

Worker n

get

get

get

get

rec

rec

rec

rec

process

process

process

process

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Iterative Computations

• Common pattern for partial differential equations, scientific computing and distributed simulation
• Workers connected to neighbors
• Data location dependent
• Workers process an iteration with results from neighbors, then send results to neighbors
• Performance bounded by slowest worker
Iterative Farmer/Worker

Farmer → Worker 1 → Worker n

process → process → process

...
Iterative P2P

Worker 1  Worker 2  Worker 3  Worker 4

comm.  process  comm.  process  comm.  process  comm.  process

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Case Study: Heat Diffusion Problem

• A problem that models heat transfer in a solid

• A two-dimensional mesh is used to represent the problem data space

• An Iterative Application

• Highly synchronized
Parallel Decomposition of the Heat Problem

Legend:
- Ghost Cells
- Data Cells
- Boundary Cells
- Ghost Cell Exchange
- 4-pt update stencil
Peer-to-Peer Computations
Peer-to-peer systems (1)

- Network transparency works well for a small number of nodes; what do we do when the number of nodes becomes very large?
  - This is what is happening now
- We need a **scalable way to handle large numbers of nodes**
- Peer-to-peer systems provide one solution
  - A distributed system that connects resources located at the edges of the Internet
  - Resources: storage, computation power, information, etc.
  - Peer software: all nodes are functionally equivalent
- Dynamic
  - Peers join and leave frequently
  - Failures are unavoidable
Peer-to-peer systems (2)

• Unstructured systems
  – Napster (first generation): still had centralized directory
  – Gnutella, Kazaa, … (second generation): neighbor graph, fully decentralized but no guarantees, often uses superpeer structure

• Structured overlay networks (third generation)
  – Using non-random topologies
  – Strong guarantees on routing and message delivery
  – Testing on realistically harsh environments (e.g., PlanetLab)
  – DHT (Distributed Hash Table) provides lookup functionality
  – Many examples: Chord, CAN, Pastry, Tapestry, P-Grid, DKS, Viceroy, Tango, Koorde, etc.
Examples of P2P networks

- Hybrid (client/server)
  - Napster

- Unstructured P2P
  - Gnutella

- Structured P2P
  - Exponential network
  - DHT (Distributed Hash Table), e.g., Chord

\[
R = N - 1 \quad \text{(hub)}
\]
\[
R = 1 \quad \text{(others)}
\]
\[
H = 1
\]

\[
R = ? \quad \text{(variable)}
\]
\[
H = 1 \ldots 7
\]
\[
\text{(but no guarantee)}
\]

\[
R = \log N
\]
\[
H = \log N
\]
\[
\text{(with guarantee)}
\]
Properties of structured overlay networks

• Scalable
  – Works for any number of nodes

• Self organizing
  – Routing tables updated with node joins/leaves
  – Routing tables updated with node failures

• Provides guarantees
  – If operated inside of failure model, then communication is guaranteed with an upper bound on number of hops
  – Broadcast can be done with a minimum number of messages

• Provides basic services
  – Name-based communication (point-to-point and group)
  – DHT (Distributed Hash Table): efficient storage and retrieval of (key,value) pairs
Self organization

- Maintaining the routing tables
  - Correction-on-use (lazy approach)
  - Periodic correction (eager approach)
  - Guided by assumptions on traffic

- Cost
  - Depends on structure
  - A typical algorithm, DKS (distributed k-ary search), achieves logarithmic cost for reconfiguration and for key resolution (lookup)

- Example of lookup for Chord, the first well-known structured overlay network
Chord: lookup illustrated

Given a key, find the value associated to the key
(here, the value is the IP address of the node that stores the key)

Assume node 0 searches for the value associated to key K with virtual identifier 7

<table>
<thead>
<tr>
<th>Interval</th>
<th>node to be contacted</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0,1)</td>
<td>0</td>
</tr>
<tr>
<td>[1,2)</td>
<td>6</td>
</tr>
<tr>
<td>[2,4)</td>
<td>6</td>
</tr>
<tr>
<td>[4,8)</td>
<td>6</td>
</tr>
<tr>
<td>[8,0)</td>
<td>12</td>
</tr>
</tbody>
</table>
Soft Real-Time
Message Properties

- SALSA provides message properties to control message sending behavior:
  - priority
    - To send messages with priority to an actor
  - delay
    - To delay sending a message to an actor for a given time
  - waitfor
    - To delay sending a message to an actor until a token is available
Priority Message Sending

• To (asynchronously) send a message with high priority:

\[
a \leftarrow \text{book(flight)}:\text{priority};
\]

*Message is placed at the beginning of the actor’s mail queue.*
Delayed Message Sending

• To (asynchronously) send a message after a given delay in milliseconds:

\[a \leftarrow \text{book(flight)}:\text{delay(1000)};\]

*Message is sent after one second has passed.*
Causal Connections
Synchronized Message Sending

- To (asynchronously) send a message after another message has been processed:

  ```
  token fundsOk = bank <- checkBalance();
  ...
  a <- book(flight):waitFor(fundsOk);
  ```

*Message is sent after token has been produced.*
Named Tokens

- Tokens can be named to enable more loosely-coupled synchronization

  - Example:

    ```
    token t1 = a1 <- m1();
    token t2 = a2 <- m2();
    token t3 = a3 <- m3(t1);
    token t4 = a4 <- m4(t2);
    a <- m(t1,t2,t3,t4);
    
    Sending m(...) to a will be delayed until messages m1()..m4() have been processed. m1() can proceed concurrently with m2().
    ```
Named Tokens (Multicast)

- Named tokens enable multicast:
  - Example:

    ```java
    token tl = a1 <- m1();
    
    for (int i = 0; i < a.length; i++) a[i] <- m(tl);
    
    Sends the result of m1 to each actor in array a.
    ```
Named Tokens (Loops)

- Named tokens allow for synchronized loops:
  
  - Example 1:
    ```java
    token t1 = initial;
    for (int i = 0; i < n; i++) {
        t1 = a <- m(t1);
    }
    
    Sends m to a n times, passing the result of the previous m as an argument.
    
    - Example 2 (using waitfor):
      ```java
      token t1 = null;
      for (int i = 0; i < a.length; i++) {
          t1 = a[i] <- m(i) : waitfor(t1);
      }
      
      Sends m(i) to actor a[i], message m(i) will wait for m(i-1) to be processed.
Join Blocks

- Join blocks allow for synchronization over multiple messages
- Join blocks return an array of objects (Object[]), containing the results of each message sent within the join block. The results are in the same order as how the messages they were generated by were sent.

  - Example:
    
    ```
    token t1 = a1 <- m1();
    join {
      for (int i = 0; i < a.length; i++) {
        a[i] <- m( t1 );
      }
    } @ process( token );
    
    Sends the message m with the result of m1 to each actor in array a. After all the messages m have been processed, their results are sent as the arguments to process.
    ```
Current Continuations

- Current continuations allow for first class access to a message’s continuation
- Current continuations facilitate writing recursive computations
  - Example:

```java
int fibonacci(int n) {
  if (n == 0) return 0;
  else if (n == 1 || n == 2) return 1;
  else {
    token a = fibonacci(n - 1);
    token b = fibonacci(n - 2);
    add(a, b) @ currentContinuation;
  }
}
```

Finds the nth fibonacci number. The result of add(a, b) is sent as the return value of fibonacci to the next message in the continuation.
Current Continuations (Loops)

• Current Continuations can also be used to perform recursive loops:
  – Example:
    ```
    void loop(int n) {
      if (n == 0) {
        m(n) @
        currentContinuation;
      } else {
        loop(n - 1) @
        m(n) @
        currentContinuation;
      }
    }
    ```
    Sends the messages m(0), m(1), m(2) ...m(n). m(i) is always processed after m(i-1).
Current Continuations (Delegation)

- Current Continuations can also be used to delegate tasks to other actors:
  - Example:

    ```java
    String getAnswer(Object question) {
        if (question instanceof Question1) {
            knowsQ1 <- getAnswer(question) @ currentContinuation;
        } else if (question instanceof Question2) {
            knowsQ2 <- getAnswer(question) @ currentContinuation;
        } else return "don't know!";
    }
    ```

    If the question is Question1 this will get the answer from actor knowsQ1 and pass this result as its token, if the question is Question2 this will get the answer from actor knowsQ2 and pass that result as its token, otherwise it will return “don't know!”.
Distributed run-time (WWC)
World-Wide Computer Architecture

- **SALSA application layer**
  - Programming language constructs for actor communication, migration, and coordination.

- **IOS middleware layer**
  - A Resource Profiling Component
    - Captures information about actor and network topologies and available resources
  - A Decision Component
    - Takes migration, split/merge, or replication decisions based on profiled information
  - A Protocol Component
    - Performs communication between nodes in the middleware system

- **WWC run-time layer**
  - Theaters provide runtime support for actor execution and access to local resources
  - Pluggable transport, naming, and messaging services
WWC Theaters

![Diagram of WWC Theaters]

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Scheduling

• The choice of which actor gets to execute next and for how long is done by a part of the system called the scheduler.

• An actor is *non-blocked* if it is processing a message or if its mailbox is not empty, otherwise the actor is *blocked*.

• A scheduler is fair if it does not starve a non-blocked actor, i.e. all non-blocked actors eventually execute.

• Fair scheduling makes it easier to reason about programs and program composition.
  – Otherwise some correct program (in isolation) may never get processing time when composed with other programs.
Remote Message Sending Protocol

• Messages between remote actors are sent using the Remote Message Sending Protocol (RMSP).
• RMSP is implemented using Java object serialization.
• RMSP protocol is used for both message sending and actor migration.
• When an actor migrates, its locator (UAL) changes but its name (UAN) does not.
Universal Actor Naming Protocol
Universal Actor Naming Protocol

- UANP includes messages for:
  - Binding actors to UAN, UAL pairs
  - Finding the locator of a universal actor given its UAN
  - Updating the locator of a universal actor as it migrates
  - Removing a universal actor entry from the naming service

- SALSA programmers need not use UANP directly in programs. UANP messages are transparently sent by WWC run-time system.
**UANP Implementations**

- Default naming service implementation stores UAN to UAL mapping in name servers as defined in UANs.
  - Name server failures may induce universal actor unreachability.

- Distributed (Chord-based) implementation uses consistent hashing and a ring of connected servers for fault-tolerance. For more information, see:


Actor Garbage Collection

- Implemented since SALSA 1.0 using pseudo-root approach.
- Includes distributed cyclic garbage collection.
- For more details, please see:


Challenge 1: Actor GC vs. Object GC

Actor Reference Graph

Passive Object Reference Graph
Challenge 2: Non-blocking communication

- Following references to mark live actors is not safe!

An example of mutation and asynchronous delivery of message
Challenge 2: Non-blocking communication

- Following references to mark live actors is not safe!
- What can we do?
  - We can protect the reference from deletion and mark the sender live until the sender knows the message has arrived.

![Diagram showing protected reference, marked live actor, blocked actor, message, and reference.]
Challenge 2: Non-blocking communication (continued)

- How can we guarantee the safety of an actor referenced by a message?
- The solution is to **protect the reference from deletion** and **mark the sender live** until the sender knows the message has arrived.
Challenge 3: Distribution and Mobility

- What if an actor is remotely referenced?
  - We can maintain an inverse reference list (only visible to the garbage collector) to indicate whether an actor is referenced.
  - The inverse reference registration must be based on non-blocking and non-First-In-First-Out communication!
  - Three operations change inverse references: actor creation, reference passing, and reference deletion.

![Actor Creation Diagram]
Challenge 3: Distribution and Mobility (continued)

- What if an actor is remotely referenced?
  - We can maintain an inverse reference list (only visible to the garbage collector) to indicate whether an actor is referenced.
  - The inverse reference registration must be based on non-blocking and non-First-In-First-Out communication!
  - Three operations are affected: actor creation, reference passing, and reference deletion.
Challenge 3: Distribution and Mobility (continued)

- What if an actor is remotely referenced?
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  - Three operations are involved: actor creation, reference passing, and reference deletion.

![Diagram showing reference passing and inverse reference registration.]

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Reference Passing 53
What if an actor is remotely referenced?

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The Pseudo Root Approach

- Pseudo roots:
  - Treat unblocked actors, migrating actors, and roots as pseudo roots.
  - Map *in-transit messages and references* into *protected references* and *pseudo roots*
  - Use inverse reference list (only visible to garbage collectors) to identify remotely referenced actors
- Actors which are not reachable from any pseudo root are garbage.
Autonomic Computing (IOS)
Middleware for Autonomous Computing

• Middleware
  – A software layer between distributed applications and operating systems.
  – Alleviates application programmers from directly dealing with distribution issues
    • Heterogeneous hardware/O.S.s
    • Load balancing
    • Fault-tolerance
    • Security
    • Quality of service

• Internet Operating System (IOS)
  – A decentralized framework for adaptive, scalable execution
  – Modular architecture to evaluate different distribution and reconfiguration strategies


Middleware Architecture

Application Layer

Component

Component

Component

Component

Physical Layer

CPU

Network

Memory

Component Profiling API

Diagnostics

FLOPS

Diagnostics

Latency / Bandwidth

Reconfiguration Requests

Virtual Network Layer (IOS Agent)

Profiling Module

Component Monitor

CPU Monitor

Network Monitor

Local Profiling Information

Join/Leave Protocol

Communication Protocol

Replication Strategies

Decision Module

Split/Merge Strategies

Migration Strategies

Remote Profiling Information and Leave Requests

Inter-IOS Agent Communication

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IOS Architecture

• IOS middleware layer
  – A Resource Profiling Component
    • Captures information about actor and network topologies and available resources
  – A Decision Component
    • Takes migration, split/merge, or replication decisions based on profiled information
  – A Protocol Component
    • Performs communication with other agents in virtual network (e.g., peer-to-peer, cluster-to-cluster, centralized.)
A General Model for Weighted Resource-Sensitive Work-Stealing (WRS)

• Given:
  A set of resources, \( R = \{r_0 \ldots r_n\} \)
  A set of actors, \( A = \{a_0 \ldots a_n\} \)
  \( \omega \) is a weight, based on importance of the resource \( r \) to the performance of a set of actors \( A \)
  \[ 0 \leq \omega(r,A) \leq 1 \]
  \[ \sum_{r} \omega(r,A) = 1 \]

  \( \alpha(r,f) \) is the amount of resource \( r \) available at foreign node \( f \)
  \( \upsilon(r,l,A) \) is the amount of resource \( r \) used by actors \( A \) at local node \( l \)
  \( M(A,l,f) \) is the estimated cost of migration of actors \( A \) from \( l \) to \( f \)
  \( L(A) \) is the average life expectancy of the set of actors \( A \)

• The predicted increase in overall performance \( \Gamma \) gained by migrating \( A \) from \( l \) to \( f \), where \( \Gamma \leq 1 \):

  \[ \Delta(r,l,f,A) = (\alpha(r,f) - \upsilon(r,l,A)) / (\alpha(r,f) + \upsilon(r,l,A)) \]
  \[ \Gamma = \sum_{r} (\omega(r,A) * \Delta(r,l,f,A)) - M(A,l,f)/(10+\log L(A)) \]

• When work requested by \( f \), migrate actor(s) \( A \) with greatest predicted increase in overall performance, if positive.
Impact of Process/Actor Granularity

Experiments on a dual-processor node (SUN Blade 1000)
Component Malleability

• New type of reconfiguration:
  – Applications can dynamically change component granularity

• Malleability can provide many benefits for HPC applications:
  – Can more adequately reconfigure applications in response to a dynamically changing environment:
    • Can scale application in response to dynamically joining resources to improve performance.
    • Can provide soft fault-tolerance in response to dynamically leaving resources.
  – Can be used to find the ideal granularity for different architectures.
  – Easier programming of concurrent applications, as parallelism can be provided transparently.
Component Malleability

- Modifying application component granularity dynamically (at runtime) to improve scalability and performance.
- SALSA-based malleable actor implementation.
- MPI-based malleable process implementation.
- IOS decision module to trigger split and merge reconfiguration.
- For more details, please see:

Distributed Systems Visualization (OverView)
Distributed Systems Visualization

- Generic online Java-based distributed systems visualization tool
- Uses a declarative Entity Specification Language (ESL)
- Instruments byte-code to send events to visualization layer.
- For more details, please see:

Overview
Heat Example With 5 Initial Nodes
http://wcl.cs.rpi.edu/overview/
Final Remarks

• Thanks!

• Visit our web pages:
  – SALSA: http://wcl.cs.rpi.edu/salsa/
  – IOS: http://wcl.cs.rpi.edu/ios/
  – OverView: http://wcl.cs.rpi.edu/overview/
  – MilkyWay@Home: http://milkyway.cs.rpi.edu/

• Questions?
Exercises

82. Create a Producer-Consumer pattern in SALSA and play with message delays to ensure that the consumer actor mailbox does not create a memory problem.

83. Create an autonomous iterative application and run it within IOS so that the management of actor placement is triggered by the middleware.

84. Execute the Cell example with OverView visualizing actor migration.