

# Logic Programming (PLP 11)

Prolog: Arithmetic, Equalities, Operators, I/O,  
Natural Language Parsing

Carlos Varela

Rensselaer Polytechnic Institute

February 9, 2015

# Arithmetic Goals

$N > M$

$N < M$

$N = < M$

$N > = M$

- $N$  and  $M$  must be bound to numbers for these tests to *succeed* or *fail*.
- $X$  **is**  $1+2$  is used to *assign* numeric value of right-hand-side to variable in left-hand-side.

# Loop Revisited

```
natural(1).  
natural(N) :- natural(M), N is M+1.  
my_loop(N) :- N>0,  
              natural(I),  
              write(I), nl,  
              I=N,  
              !.  
  
my_loop(_).
```

Also called *generate-and-test*.

**= is not equal to == or ::=**

$X=Y$

$X\backslash=Y$

test whether  $X$  and  $Y$  **can be** or **cannot be** *unified*.

$X==Y$

$X\backslash==Y$

test whether  $X$  and  $Y$  are currently *co-bound*, i.e., have been bound to, or share the same value.

$X>::=Y$

$X=\backslash=Y$

test *arithmetic* equality and inequality.

# More equalities

$X=@=Y$

$X\backslash=@=Y$

test whether  $X$  and  $Y$  are *structurally identical*.

- $=@=$  is weaker than  $==$  but stronger than  $=$ .
- Examples:

$a=@=A$

**false**

$A=@=B$

**true**

$x(A, A) =@= x(B, C)$

**false**

$x(A, A) =@= x(B, B)$

**true**

$x(A, B) =@= x(C, D)$

**true**

# More on equalities

$$X==Y \Rightarrow X=@=Y \Rightarrow X=Y$$

but not the other way ( $\Leftarrow$ ).

- If two terms are currently **co-bound**, they are **structurally identical**, and therefore they can **unify**.
- Examples:

$a=@=A$	<b>false</b>
$A=@=B$	<b>true</b>
$x(A, A)=@=x(B, C)$	<b>false</b>
$x(A, A)=@=x(B, B)$	<b>true</b>
$x(A, B)=@=x(C, D)$	<b>true</b>

# Prolog Operators

```
:- op(P, T, O)
```

declares an operator symbol  $O$  with precedence  $P$  and type  $T$ .

- Example:

```
:- op(500, xfx, 'has_color')
```

```
a has_color red.
```

```
b has_color blue.
```

then:

```
?- b has_color C.
```

```
C = blue.
```

```
?- What has_color red.
```

```
What = a.
```

# Operator precedence/type

- Precedence **P** is an integer: the larger the number, the less the precedence (*ability to group*).
- Type **T** is one of:

<b>T</b>	<b>Position</b>	<b>Associativity</b>	<b>Examples</b>
<b>xfx</b>	Infix	Non-associative	is
<b>xfy</b>	Infix	Right-associative	, ;
<b>yfx</b>	Infix	Left-associative	+ - * /
<b>fx</b>	Prefix	Non-associative	?-
<b>fy</b>	Prefix	Right-associative	
<b>xf</b>	Postfix	Non-associative	
<b>yf</b>	Postfix	Left-associative	



# Testing types

**atom**(X)

tests whether X is an *atom*, e.g., 'foo', bar.

**integer**(X)

tests whether X is an *integer*; it does not test for complex terms, e.g., `integer(4/2)` fails.

**float**(X)

tests whether X is a *float*; it matches exact type.

**string**(X)

tests whether X is a *string*, enclosed in `` ... ``.

# Prolog Input

**seeing** (X)

succeeds if X is (or can be) bound to *current read port*.

X = user is keyboard (standard input.)

**see** (X)

*opens* port for input file bound to X, and makes it *current*.

**seen**

*closes* current port for input file, and makes user *current*.

**read** (X)

*reads* Prolog type expression from *current* port, storing value in X.

**end-of-file**

is returned by **read** at *<end-of-file>*.

# Prolog Output

**telling** (X)

succeeds if X is (or can be) bound to *current output port*.

X = user is screen (standard output.)

**tell** (X)

*opens* port for output file bound to X, and makes it *current*.

**told**

*closes* current output port, and reverses to screen output  
(makes user *current*.)

**write** (X)

*writes* Prolog expression bound to X into *current* output port.

**nl**

new line (line feed).

**tab** (N)

writes N spaces to current output port.

# I/O Example

```
browse(File) :-
    seeing(Old),          /* save for later */
    see(File),           /* open this file */
    repeat,
    read(Data),          /* read from File */
    process(Data),
    seen,                /* close File */
    see(Old),            /* prev read source */
    !.                  /* stop now */

process(end_of_file) :- !.
process(Data) :- write(Data), nl, fail.
```

# First-Class Terms Revisited

<code>call(P)</code>	Invoke predicate as a goal.
<code>assert(P)</code>	Adds predicate to database.
<code>retract(P)</code>	Removes predicate from database.
<code>functor(T, F, A)</code>	Succeeds if T is a <i>term</i> with <i>functor</i> F and <i>arity</i> A.
<code>findall(F, P, L)</code>	Returns a list L with all elements F satisfying predicate P
<code>clause(H, B)</code>	Succeeds if the clause <code>H :- B</code> can be found in the database.

# Natural Language Parsing

(Example from "Learn Prolog Now!" Online Tutorial)

```
word(article,a) .
```

```
word(article, every) .
```

```
word(noun, criminal) .
```

```
word(noun, 'big kahuna burger') .
```

```
word(verb, eats) .
```

```
word(verb, likes) .
```

```
sentence(Word1, Word2, Word3, Word4, Word5) :-
```

```
    word(article, Word1) ,
```

```
    word(noun, Word2) ,
```

```
    word(verb, Word3) ,
```

```
    word(article, Word4) ,
```

```
    word(noun, Word5) .
```

# Parsing natural language

- *Definite Clause Grammars (DCG)* are useful for natural language parsing.
- Prolog can load DCG rules and convert them automatically to Prolog parsing rules.

# DCG Syntax

-->

DCG *operator*, e.g.,

sentence-->subject, verb, object.

Each goal is assumed to refer to the *head* of a DCG rule.

**{prolog\_code}**

*Include* Prolog code in generated parser, e.g.,

subject-->modifier, noun, {write('subject') }.

**[terminal\_symbol]**

*Terminal* symbols of the grammar, e.g.,

noun-->[cat].



# Natural Language Parsing

(example rewritten using DCG)

sentence --> article, noun, verb, article, noun.

article --> [a] | [every].

noun --> [criminal] | ['big kahuna burger'].

verb --> [eats] | [likes].

# Natural Language Parsing (2)

sentence(V) --> subject, verb(V), subject.

subject --> article, noun.

article --> [a] | [every].

noun --> [criminal] | ['big kahuna burger'].

verb(eats) --> [eats].

verb(likes) --> [likes].

# Difference lists in Prolog

- A *difference list* is a pair of lists, each might have an unbound tail, with the invariant that one can get the second list by removing zero or more elements from the first list
- $X, X$                     % Represent the empty list
- $[], []$                     % idem
- $[a], [a]$                    % idem
- $[a,b,c|X], X$               % Represents  $[a,b,c]$
- $[a,b,c,d], [d]$             % idem

# Difference lists in Prolog (2)

- When the second list is unbound, an append operation with another difference list takes constant time

```
append_dl(S1,E1, S2,E2, S1,E2) :- E1 = S2.
```

- ?- append\_dl([1,2,3|X],X, [4,5|Y],Y, S,E).

Displays

```
X = [4, 5|_G193]
```

```
Y = _G193
```

```
S = [1, 2, 3, 4, 5|_G193]
```

```
E = _G193 ;
```

# Exercises

12. How would you translate DCG rules into Prolog rules?
13. PLP Exercise 11.8 (pg 571).
14. PLP Exercise 11.14 (pg 572).