

CSCI-1200 Data Structures — Spring 2023

Lecture 27 — Concurrency & Asynchronous Computing

Final Exam General Information

- The final exam will be held:

Thursday May 4th from 3-6pm.

A makeup exam will only be offered if required by the RPI rules regarding final exam conflicts *-OR-* if a written excuse from the Dean of Students office is provided. Contact ds_instructors@cs.rpi.edu immediately if you have a conflict.

- Coverage: Lectures 1-27, Labs 1-14, and HW 1-10.
 - For the final for your notes you are allowed two physical pieces of 8.5x11” paper, that’s four “sides”. We will **not** collect these electronically and we will **not** pre-print them, you will have to bring these notes pages yourself if you want them. We will check at the start of the exam that you do not have more than two pieces of paper for your notes!
- The best thing you can do to prepare for the final is practice. Try the review problems (posted on the course website) with pencil & paper first. Then practice programming (with a computer) the exercises and other exercises from lecture, lab, and the homework.
- You will have 3 hours to complete the final.
- The final will be cumulative, with 150 points.
- Course evaluation responses to date: $140 / 466 = 30.00\%$
- Those of you interested in becoming an undergraduate mentor for Data Structures, or another CSCI course: contact your graduate lab TA or undergraduate programming mentor and ask him/her to recommend you for the position. Also fill out the online application (will be emailed shared by Shianne Hulbert & David Goldschmidt).

27.1 Today’s Class

- Computing with multiple threads/processes and one or more processors
- Shared resources & mutexes/locks
- Deadlock: the Dining Philosopher’s Problem

27.2 The Role of Time in Evaluation

- Sometimes the order of evaluation does matter, and sometimes it doesn’t.
 - The behavior of objects with *state* depends on sequence of events that have occurred.
 - *Referential transparency*: when equivalent expressions can be substituted for one another without changing the value of the expression. For example, a complex expression can be replaced with its result *if* repeated evaluations always yield the same result, independent of context.
- What happens when objects don’t change one at a time but rather act concurrently?
 - We may be able to take advantage of this by letting threads/processes run at the same time (a.k.a., in parallel).
 - However, we will need to think carefully about the interactions and shared resources.

27.3 Concurrency Example: Joint Bank Account

- Consider the following bank account implementation:

```
class Account {
public:
    Account(int amount) : balance(amount) {}
    void deposit(int amount) {
        int tmp = balance;           // A
        tmp += amount;               // B
        balance = tmp;               // C
    }
    void withdraw(int amount) {
        int tmp = balance;           // D
        if (amount > tmp)
            cout << "Error: Insufficient Funds!" << endl; // E1
        else {
            tmp -= amount;           // E2
        }
        balance = tmp;               // F
    }
private:
    int balance;
};
```

- We create a joint account that will be used by two people (threads/processes):

```
Account account(100);
```

- Now, enumerate all of the possible interleavings of the sub-expressions (A-F) if the following two function calls were to happen concurrently. What are the different outcomes?

```
account.deposit(50);
account.withdraw(125);
```

- What if instead the actions were:

```
account.deposit(50);
account.withdraw(75);
```

27.4 Correct/Acceptable Behavior of Concurrent Programs

- No two operations that change any shared state variables may occur at the same time.
 - Certain low-level operations are guaranteed to execute *atomic*-ly (from start to finish without interruption), but this varies based on the hardware and operating system. We need to know which operations are *atomic* on our hardware.
 - In the bank account example we *cannot* assume that the `deposit` and `withdraw` functions are atomic.
- The concurrent system should produce a result of the threads/processes running sequentially *in some order*.
 - We do not require that the threads/processes run sequentially, only that they produce results as if they had run sequentially.
 - *Note:* There may be more than one correct result!
- **Exercise:** What are the acceptable outcomes for the bank account example?

27.5 Serialization via a Mutex

- We can *serialize* the important interactions using a primitive, atomic synchronization method called a *mutex*.
- Once one thread has acquired the mutex (locking the resource), no other thread can acquire the mutex until it has been released.

- In the example below we use the STL mutex object (`#include <mutex>`). If the mutex is unavailable, the call to the mutex member function `lock()` *blocks* (the thread pauses at that line of code until the mutex is available).

```
class Chalkboard {
public:
    Chalkboard() { }
    void write(Drawing d) {
        board.lock();
        drawing = d;
        board.unlock();
    }
    Drawing read() {
        board.lock();
        Drawing answer = drawing;
        board.unlock();
        return answer;
    }
private:
    Drawing drawing;
    std::mutex board;
};
```

- What does the mutex do in this code?

27.6 The Professor & Student Classes

- Here are two simple classes that can communicate through a shared Chalkboard object:

```
class Professor {
public:
    Professor(Chalkboard *c) { chalkboard = c; }
    virtual void Lecture(const std::string &notes) {
        chalkboard->write(notes);
    }
protected:
    Chalkboard* chalkboard;
};
```

```
class Student {
public:
    Student(Chalkboard *c) { chalkboard = c; }
    void TakeNotes() {
        Drawing d = chalkboard->read();
        notebook.push_back(d);
    }
private:
    Chalkboard* chalkboard;
    std::vector<Drawing> notebook;
};
```

27.7 Launching Concurrent Threads

- So how exactly do we get multiple streams of computation happening simultaneously? There are many choices (may depend on your programming language, operating system, compiler, etc.).
- We'll use the STL `thread` library (`#include <thread>`). The new thread begins execution in the provided function (`student_thread`, in this example). We pass the necessary shared data from the main thread to the secondary thread to facilitate communication.

```

#define num_notes 10

void student_thread(Chalkboard *chalkboard) {
    Student student(chalkboard);
    for (int i = 0; i < num_notes; i++) {
        student.TakeNotes();
    }
}

int main() {
    Chalkboard chalkboard;
    Professor prof(&chalkboard);
    std::thread student(student_thread, &chalkboard);
    for (int i = 0; i < num_notes; i++) {
        prof.Lecture("blah blah");
    }
    student.join();
}

```

- The `join` command pauses to wait for the secondary thread to finish computation before continuing with the program (or exiting in this example).
- What can still go wrong? How can we fix it?

27.8 Condition Variables

- Here we've added a *condition variable*, `student_done`:

```

class Chalkboard {
public:
    Chalkboard() { student_done = true; }
    void write(Drawing d) {
        while (1) {
            board.lock();
            if (student_done) {
                drawing = d;
                student_done = false;
                board.unlock();
                return;
            }
            board.unlock();
        }
    }
    Drawing read() {
        while (1) {
            board.lock();
            if (!student_done) {
                Drawing answer = drawing;
                student_done = true;
                board.unlock();
                return answer;
            }
            board.unlock();
        }
    }
private:
    Drawing drawing;
    std::mutex board;
    bool student_done;
};

```

- *Note:* This implementation is actually quite inefficient due to “busy waiting”. A better solution is to use an operating system-supported *condition variable* that yields to other threads if the lock is not available and is signaled when the lock becomes available again. STL has a `condition_variable` type which allows you to wait for or notify other threads that it may be time to resume computation.

27.9 Exercise: Multiple Students and/or Multiple Professors

- Now consider that we have multiple students and/or multiple professors. How can you ensure that each student is able to copy a complete set of notes?

27.10 Multiple Locks & Deadlock

- For this last example, we add two public member variables of type `std::mutex` to the `Chalkboard` class, named `chalk` and `textbook`.
- And we derive two different types of lecturer from the base class `Professor`. The professors can lecture concurrently, but they must share the chalk and the book.

```
class CautiousLecturer : public Professor {
public:
    CautiousLecturer(Chalkboard *c) : Professor(c) {}
    void Lecture() {
        chalkboard->textbook.lock();
        Drawing d = FromBookDrawing();
        chalkboard->chalk.lock();
        Professor::Lecture(d);
        chalkboard->chalk.unlock();
        chalkboard->textbook.unlock();
    }
};
```

```
void checkDrawing(const Drawing &d) {}
```

```
class BrashLecturer : public Professor {
public:
    BrashLecturer(Chalkboard *c) : Professor(c) {}
    void Lecture() {
        chalkboard->chalk.lock();
        Drawing d = FromMemoryDrawing();
        Professor::Lecture(d);
        chalkboard->textbook.lock();
        checkDrawing(d);
        chalkboard->textbook.unlock();
        chalkboard->chalk.unlock();
    }
};
```

- What can go wrong? How can we fix it?
Why might philosophers discuss this problem over dinner?

27.11 Topics Covered

- Algorithm analysis: big O notation; best case, average case, or worst case; algorithm running time or additional memory usage
- STL classes: `string`, `vector`, `list`, `map`, & `set`, (we talked about but did not practice using `STL stack`, `queue`, `unordered_set`, `unordered_map`, & `priority_queue`)
- C++ Classes: constructors (default, copy, & custom argument), assignment operator, & destructor, classes with dynamically-allocated memory, operator overloading, inheritance, polymorphism
- Subscripting (random-access, pointer arithmetic) vs. iteration
- Recursion & problem solving techniques

- Memory: pointers & arrays, heap vs. stack, dynamic allocation & deallocation of memory, garbage collection, smart pointers
- Implementing data structures: resizable arrays (vectors), linked lists (singly-linked, doubly-linked, circularly-linked, dummy head/tail nodes), trees (for sets & maps), hash sets
- Binary Search Trees, tree traversal (in-order, pre-order, post-order, depth-first, & breadth-first), ropes
- Hash tables (hash functions, collision resolution), priority queues, heap as a vector
- Exceptions, concurrency & asynchronous computing

27.12 Course Summary

- Approach any problem by studying the requirements carefully, playing with hand-generated examples to understand them, and then looking for analogous problems that you already know how to solve.
- STL offers container classes and algorithms that simplify the programming process and raise your conceptual level of thinking in designing solutions to programming problems. Just think how much harder some of the homework problems would have been without generic container classes!
- When choosing between algorithms and between container classes (data structures) you should consider:
 - efficiency,
 - naturalness of use, and
 - ease of programming.
- Use classes with well-designed public and private member functions to encapsulate sections of code.
- Writing your own container classes and data structures usually requires building linked structures and managing memory through the big three:
 - copy constructor,
 - assignment operator, and
 - destructor.
- When testing and debugging:
 - Test one function and one class at a time,
 - Figure out what your program actually does, not what you wanted it to do,
 - Use small examples and boundary conditions when testing, and
 - Find and fix the first mistake in the flow of your program before considering other apparent mistakes.
- Above all, remember the excitement and satisfaction when your hard work and focused debugging is rewarded with a program that demonstrates your technical mastery and realizes your creative problem solving skills!