

A Language and Architecture for Distributed Computing over the Internet

Carlos A. Varela

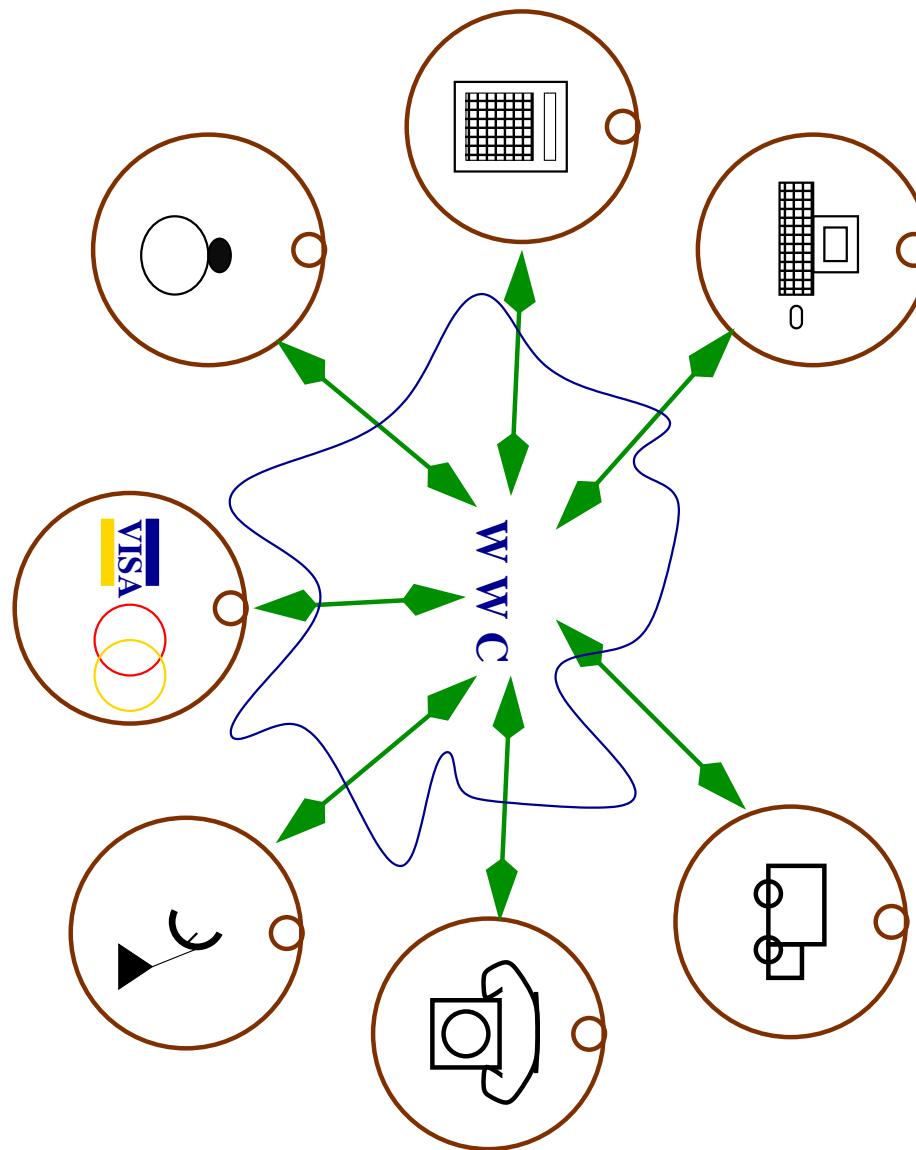
Worldwide Computing Lab

Department of Computer Science
Rensselaer Polytechnic Institute

<http://www.cs.rpi.edu/wcc/>

July 2002

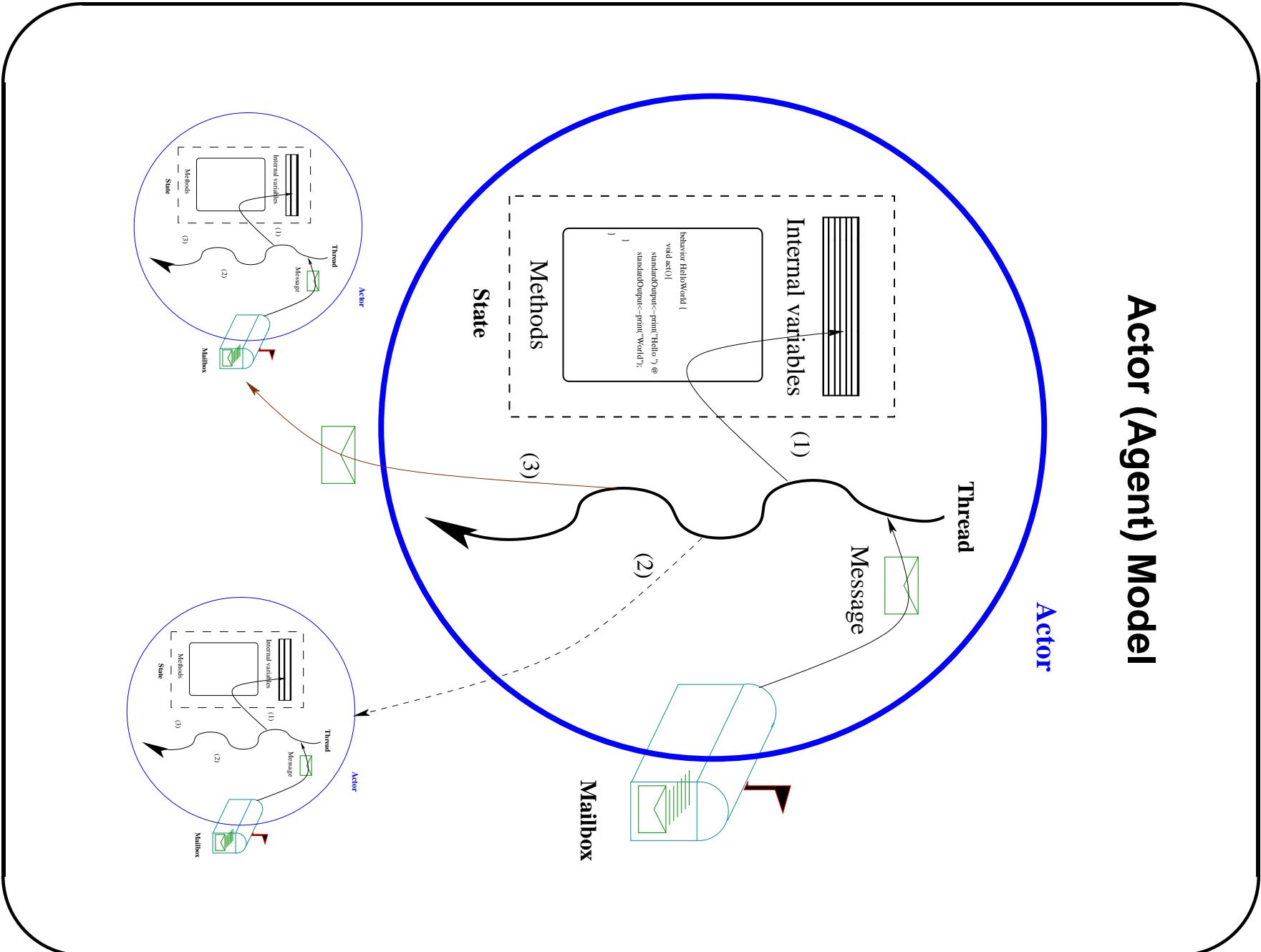
Worldwide Computing



- only primitive synchronization mechanisms
- non-universal naming
- synchronous communication
- shared memory
- passive objects

L i m i t a t i o n s o f j a v a f o r W o r l d w i d e C o m p u t i n g

Actor (Agent) Model



- inherent concurrency – threads are encapsulated in objects
- non-blocking, asynchronous communication

Coordination of actors in worldwide open systems remains difficult.

- universal naming
- distributed memory

Actor mobility is simpler than object mobility:

Actor (Agent) Model (continued)

Worldwide Computing Architecture

The **World-Wide Computer** consists of concurrent, distributed, and mobile Universal Actors.

- Universal **naming** strategy
- Run-time support: **Theatres**

- Remote communication protocol

- Migration support

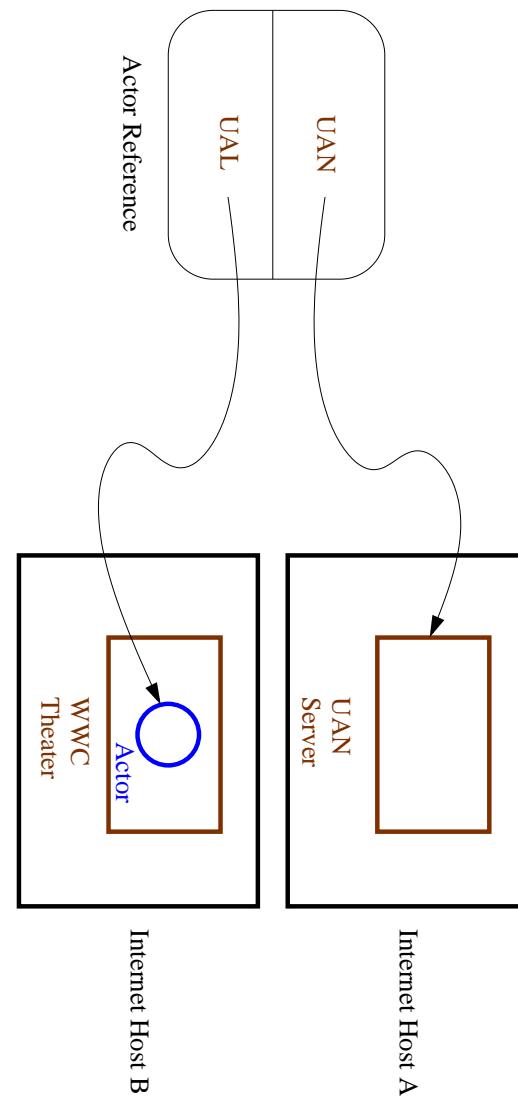
- Preliminary **performance** results

The main goals of **naming** in worldwide computing are to provide:

Naming in Worldwide Computing (Requirements)

- **both human and computer readability**
- **openness** by allowing unanticipated actor reference creation and protocols that provide access through names
- **transparent actor migration**
- **scalability** of name space management
- **platform independence** – names should appear coherent on all nodes independently of underlying architecture

Proposed Solution: Universal Actor Names and Locators



- Uniform Resource Identifiers (URI) syntax [Berners-Lee]
- UAN/UAL support **transparent** actor migration.

Sample UAN:

uan: //wwc.osl.cs.uiuc.edu:3030/Agha/Calendar

Sample Universal Actor Locators for this WWC actor:

```
rmsp://agha.cs.uiuc.edu:4040/Agents/Calendar  
rmsp://howard.cs.uiuc.edu:4040/AghaCalendar  
rmsp://agha.pda.com:1234/Calendar
```

An actor's location can be cached for faster future accesses.

Method	Parameters	Action
PUT	relative UAN, UAL	Creates a new entry in the database
GET	relative UAN	Returns the UAL entry in the database
DELETE	relative UAN	Deletes the entry in the database
UPDATE	relative UAN, UAL	Updates the UAL entry in the database

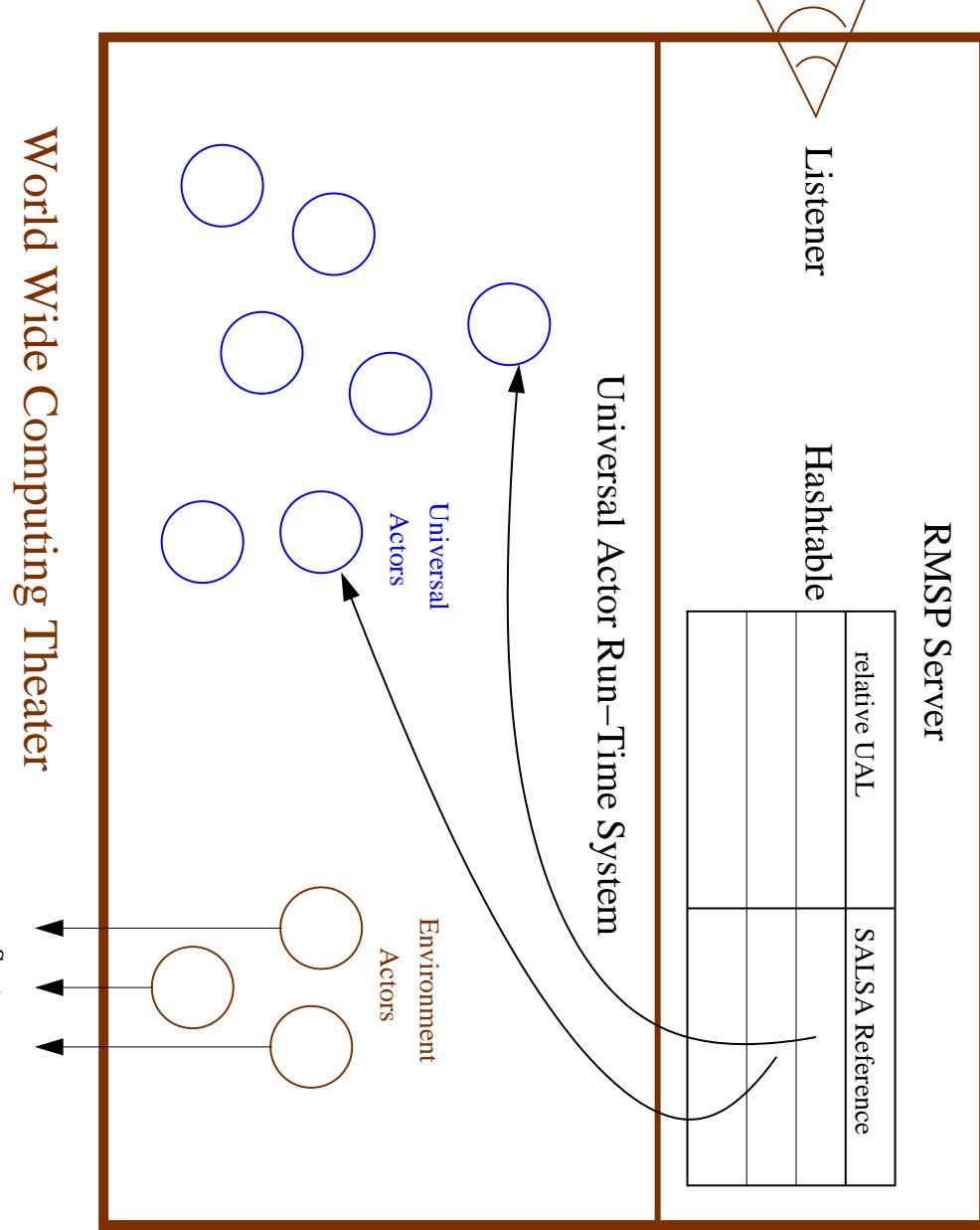
The **UANP** defines how to interact with the WWC Naming Service. Similarly to HTTP, UANP is text-based, and includes methods for the following actions:

Universal Actor Naming Protocol

World-Wide Computer Theaters

A WWC Theater provides runtime support to **Universal Actors**. A Theater contains:

- a **remote communication module** with a hashtable mapping relative UALs to actual SALSA actor references, and
- a **runtime system** for universal and environment actors.



The main goals of a remote communication protocol in worldwide computing are to provide:

- asynchronous, non-blocking communication

- an interface to the naming service for target actor location

- data and code mobility

Remote Communication in Worldwide Computing (Requirements)

- RMS^P defines how an actor sends a message to any other actor in the World-Wide Computer actor migration
- RMS^P is object-based, and includes support for message serialization, and transparent programming language support for sending messages to local and remote actors

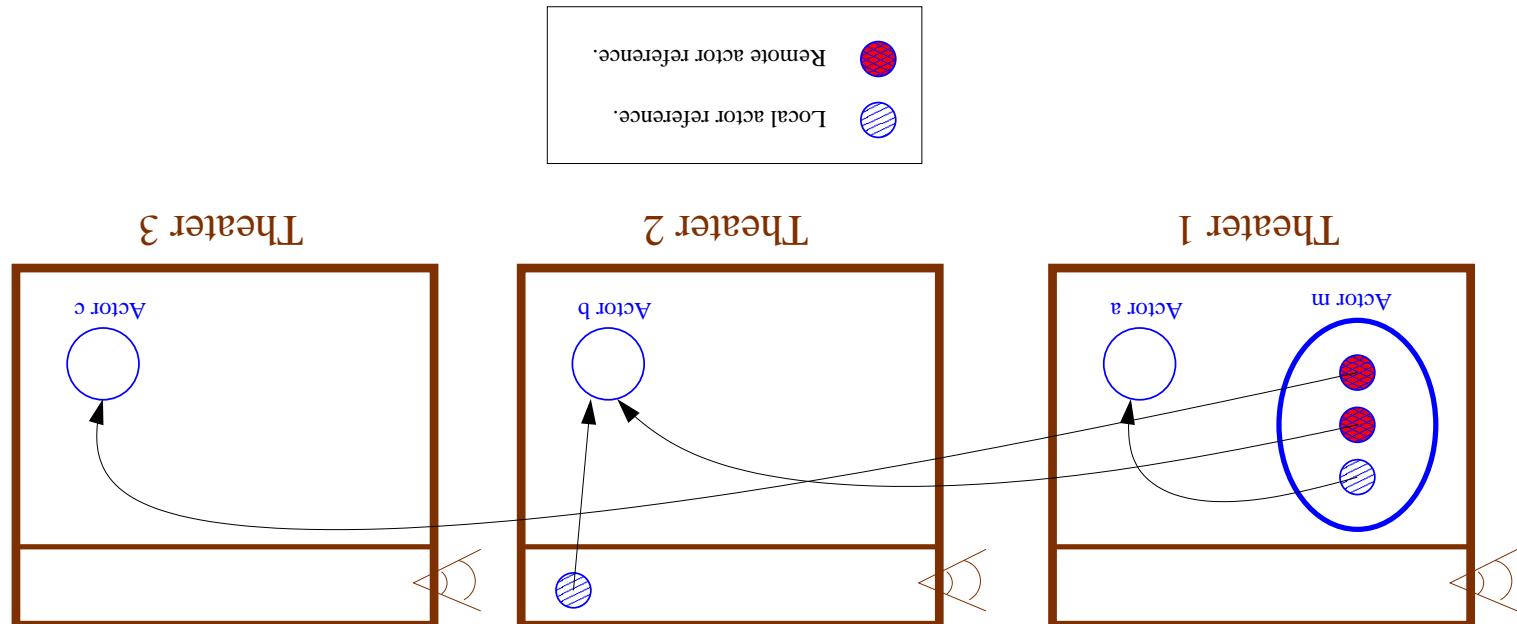
Proposed Solution: Remote Message Sending Protocol

Since actors do not have shared memory, actor migration is simpler and more efficient than object/thread migration.

- consistency protocols for shared memory.
- actor reference updating for more efficient local communication
- both fine-grained and coarse-grained mobility

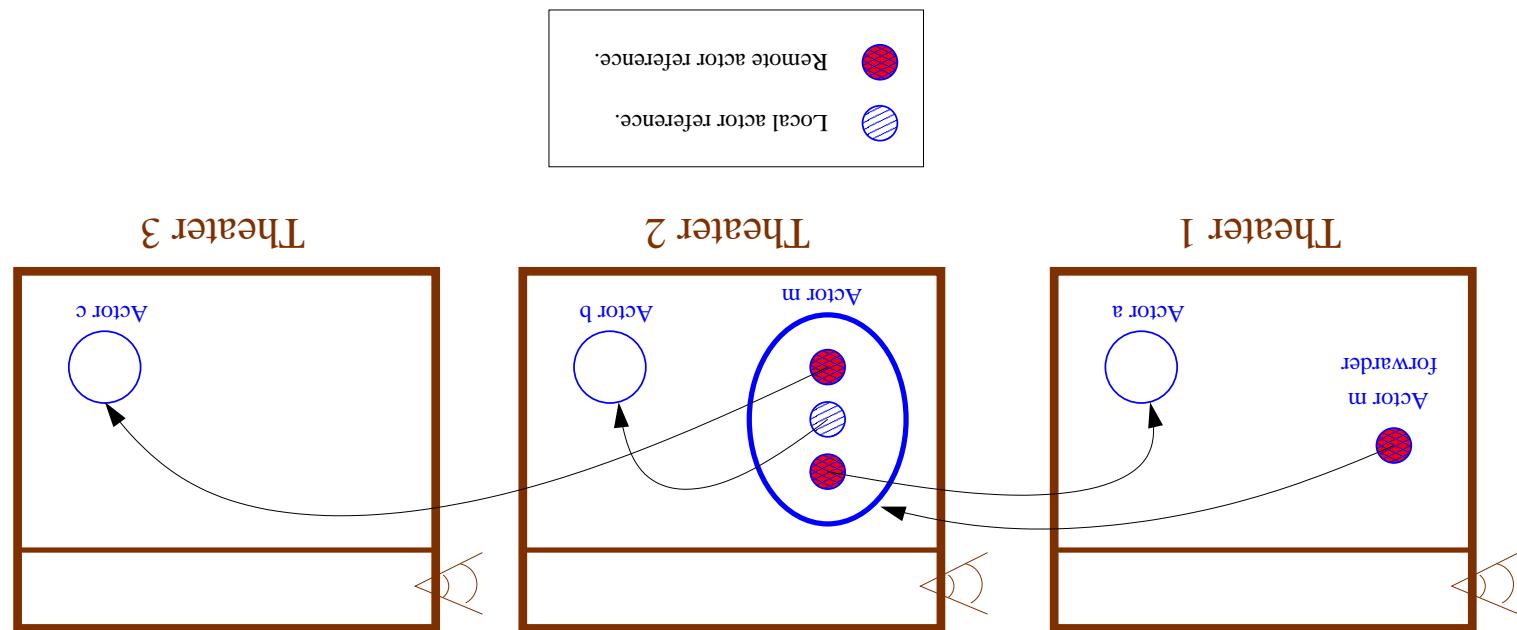
The main goals of migration in worldwide computing are to provide:

Migration in Worldwide Computing (Requirements)



Before migration of actor **m** from Theater 1 to Theater 2, its references to actors **b** and **c** are **remote**, while its reference to actor **a** is **local**.

Actor Migration



reference to actor **b** becomes **local**. Its reference to actor **c** remains unchanged.
After migration of actor **m**, its reference to actor **a** becomes **remote** and its

Actor Migration (continued)

Local actor creation time	386 μ s	Local message sending time	148 μ s	LAN message sending time	30-60ms	WAN message sending time	2-3 secs
LAN actor migration time	150-160ms (minimal actor)	LAN actor migration time	240-250ms (actor with 100kb of data)	LAN actor migration time	3-7secs (minimal actor)	WAN actor migration time	25-30secs (actor with 100kb of data)
WAN actor migration time	25-30secs (actor with 100kb of data)	WAN actor migration time	3-7secs (minimal actor)	WAN actor migration time	240-250ms (actor with 100kb of data)	LAN actor migration time	150-160ms (minimal actor)
WAN actor migration time	25-30secs (actor with 100kb of data)	WAN actor migration time	3-7secs (minimal actor)	WAN actor migration time	240-250ms (actor with 100kb of data)	LAN actor migration time	150-160ms (minimal actor)
WAN actor migration time	25-30secs (actor with 100kb of data)	WAN actor migration time	3-7secs (minimal actor)	WAN actor migration time	240-250ms (actor with 100kb of data)	LAN actor migration time	150-160ms (minimal actor)

Time Ranges (with SALSAs 0.3.2)

Machine Name	Location	OS-JVM	Processor
solarisr.co.jp	Tokyo, Japan	Solaris 2.6-JDK 1.1.6	Sparc 20
vulcan.ecoledoc.lip6.fr	Paris, France	Linux 2.2.5-JDK 1.2pre2	PII, 350MHz
yangtze.cs.uic.edu	Urbana, IL, USA	Solaris 2.5.1-JDK 1.1.6	Ultra 2

World-Wide Computer Testbed

Simple Actor Language, System and Architecture

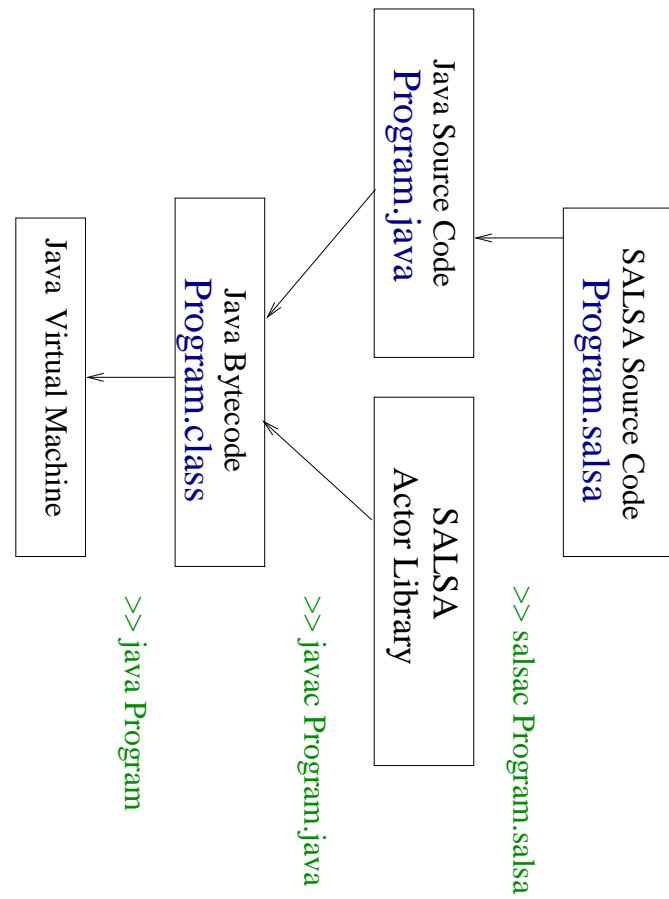
SALSA is a dialect of Java, with support for:

- Concurrent programming with **actors**.
- **Token-passing Continuations** to control concurrency by specifying customers for an actor message's return value.
- **Join Continuations** to provide a synchronization barrier for multiple actor computations into a single continuation.
- **Universal Naming** to bind and locate actors in the WWW using UANs and UALs.
- **RMS and Migration** to send messages to remote actors and to migrate actors across WWW Theaters.

Language Implementation

- SALSA Actor Library
- Join Continuation Code Generation

SALSA Architecture



```
module helloWorld;
behavior HelloWorld {
    void act (String arguments[])
    {
        standardOutput <- print ("Hello " + arguments[0]);
        standardOutput <- print ("World!");
    }
}
```

SALSA Hello World Example

Actor Model Support

- All SALSA behaviors inherit from the `UniversalActor` class. To create a new actor instance:

```
HelloWorld helloworld = new HelloWorld();
```

- To send messages to acquaintance actors:

```
acquaintance <- message( arg1, arg2, ... );
```

For example:

```
standardOutput <- print( "Hello " );
hello();
```

- Sending a message returns immediately (it is **asynchronous**) with a void return value.
- Only an actor itself can **change its internal state** through assignments to its instance variables. **No shared memory** or static variables are allowed in SALSA.

Token-Passing Continuations

- SALSA messages are potential Java method invocations. Message passing is asynchronous.
- Continuations allow to specify a customer for a message's return value (called token):

```
acquaintance<-m1( args ) @  
customer<-m2( arg0, ... , token, ... , argn );
```

For example:

```
fractal<-computePixel( ) @  
screen<-draw(token);
```

- The return type of computePixel() needs to match the formal argument type of draw(argument). SALSA allows method overloading and will choose the most specific method according to the token run-time type.
- a<-m with no arguments, is syntactic sugar for a<-m(token). For example:


```
fractal<-computePixel( ) @ screen<-draw;
```
- It is possible to "chain" continuations. For example:


```
a1<-m1( ) @ a2<-m2 @ a3<-m3( token, 10 );
```

join statement.
argument is an array containing the return values of the messages inside the
authors have finished writing their chapters. The token passed as an
will only send the message review (token) to the editor when both

```
editor->review @ publicasher->print;  
join(authors[writeChapter(1), authors[writeChapter(2)]]) @
```

For example:

```
join(actorArray->m()) @ customer <- n;  
join(a1->m1(args), a2->m2(args), ...) @ customer <- n;
```

- A join statement allows to provide a synchronization barrier for multiple actor computations into a single continuation:

join Continuations

```
behavior Agent {  
    void printIntArray() { . . . }  
    void act(String[] args) {  
        Agent a = new Agent();  
        try {  
            a->bind("uan://Yangtze.cs.unique.edu:3030/agent",  
                    "rmp://Yangtze.cs.unique.edu:4040/agent");  
            catch (Exception e) {  
                standardOutput->println(e);  
            }  
        }  
    }  
}
```

Universal Actor Naming in SALSA

RMS and Actor Migration in SALSA

Getting a remote actor reference by name and sending a message:

```
Agent a = new Agent();  
a->getRemoteReference("uan://yangtze.cs.uiuc.edu/agent");  
a->printItinerary();  
a->getRemoteLocation("rmsp://yangtze.cs.uiuc.edu/agent");  
a->printItinerary();
```

Getting the reference by location:

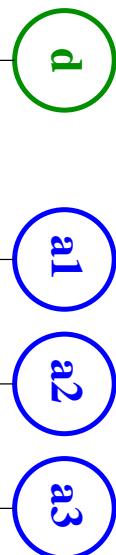
```
Agent a = new Agent();  
a->migrate("rmsp://vulcan.ecoledoc.lip6.fr/agent");  
a->getRemoteName("uan://yangtze.cs.uiuc.edu/agent")  
a->migrate("rmsp://vulcan.ecoledoc.lip6.fr/agent");
```

Migrating an agent to a remote WWC Theater:

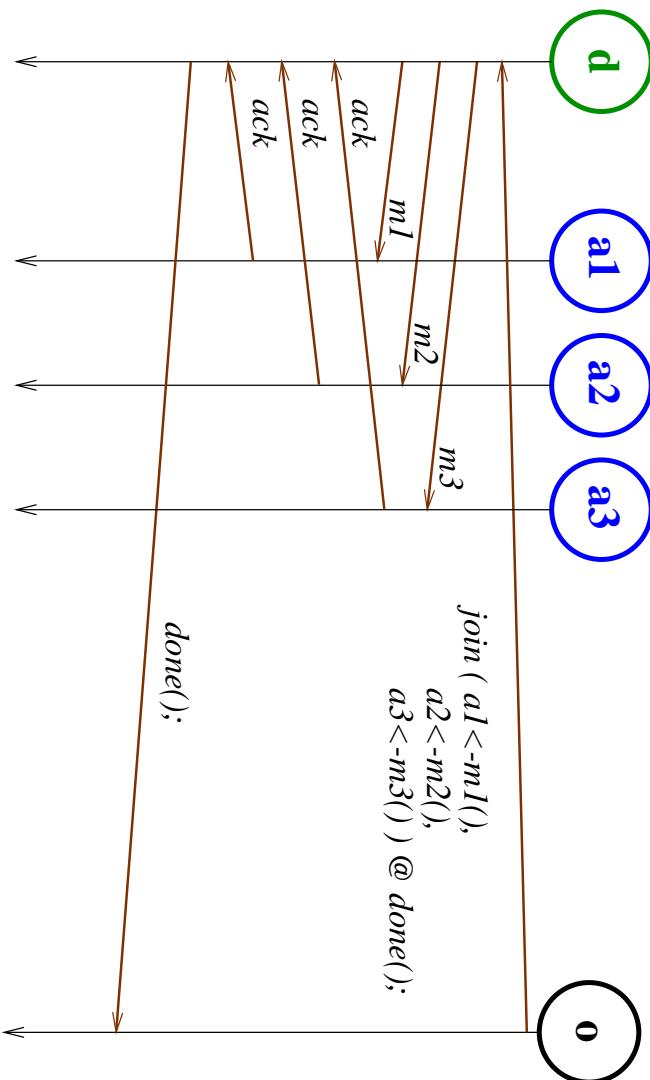
Example 1: Multicast Protocols

Multicast

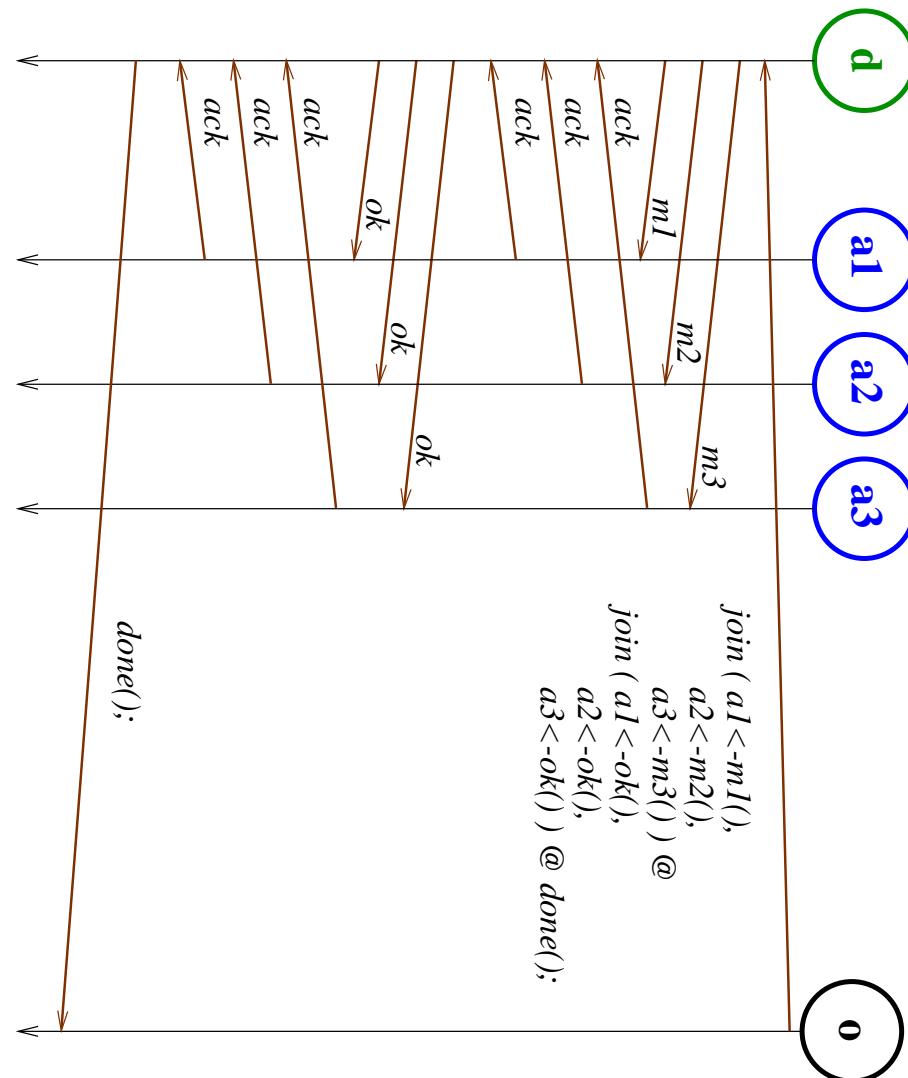
```
join ( al <-m1(),  
       a2 <-m2(),  
       a3 <-m3());
```



Acknowledged multicast

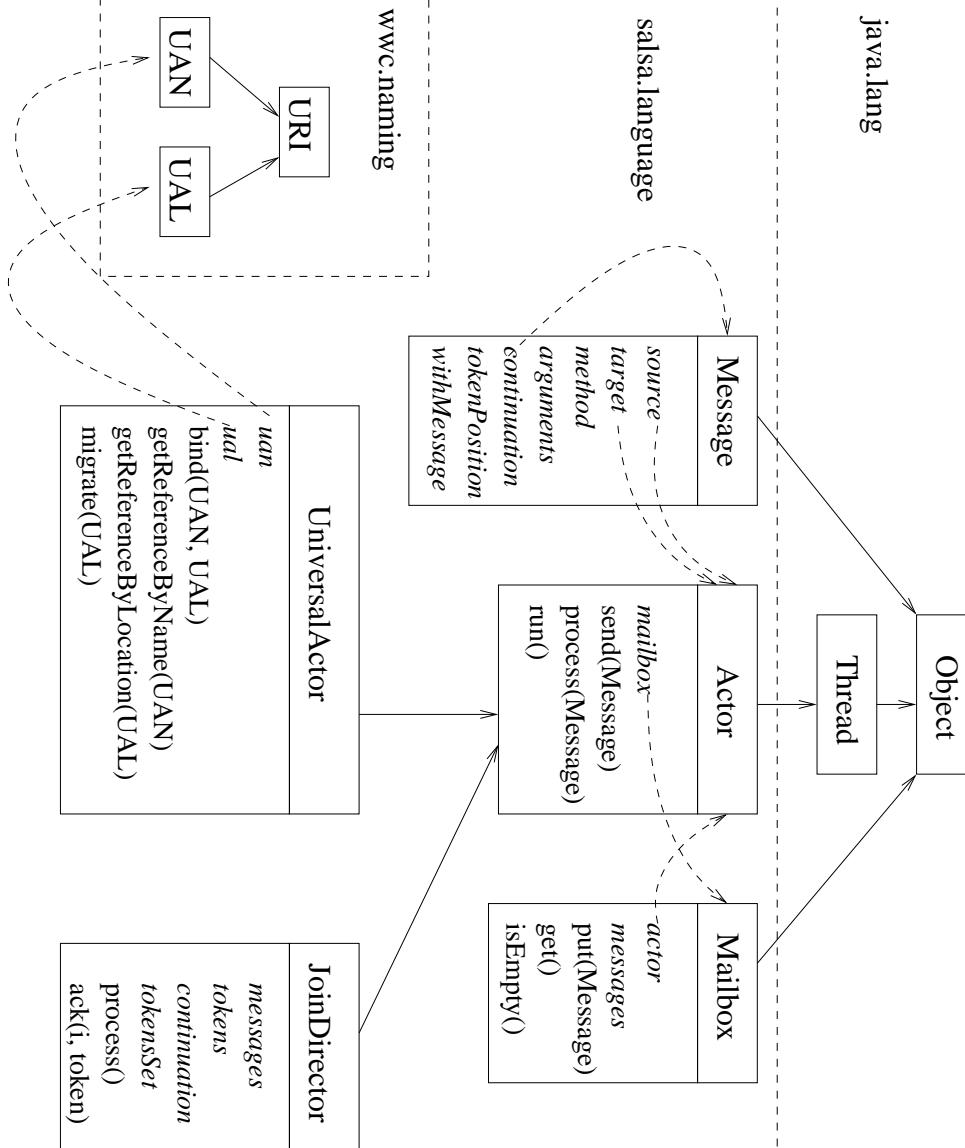


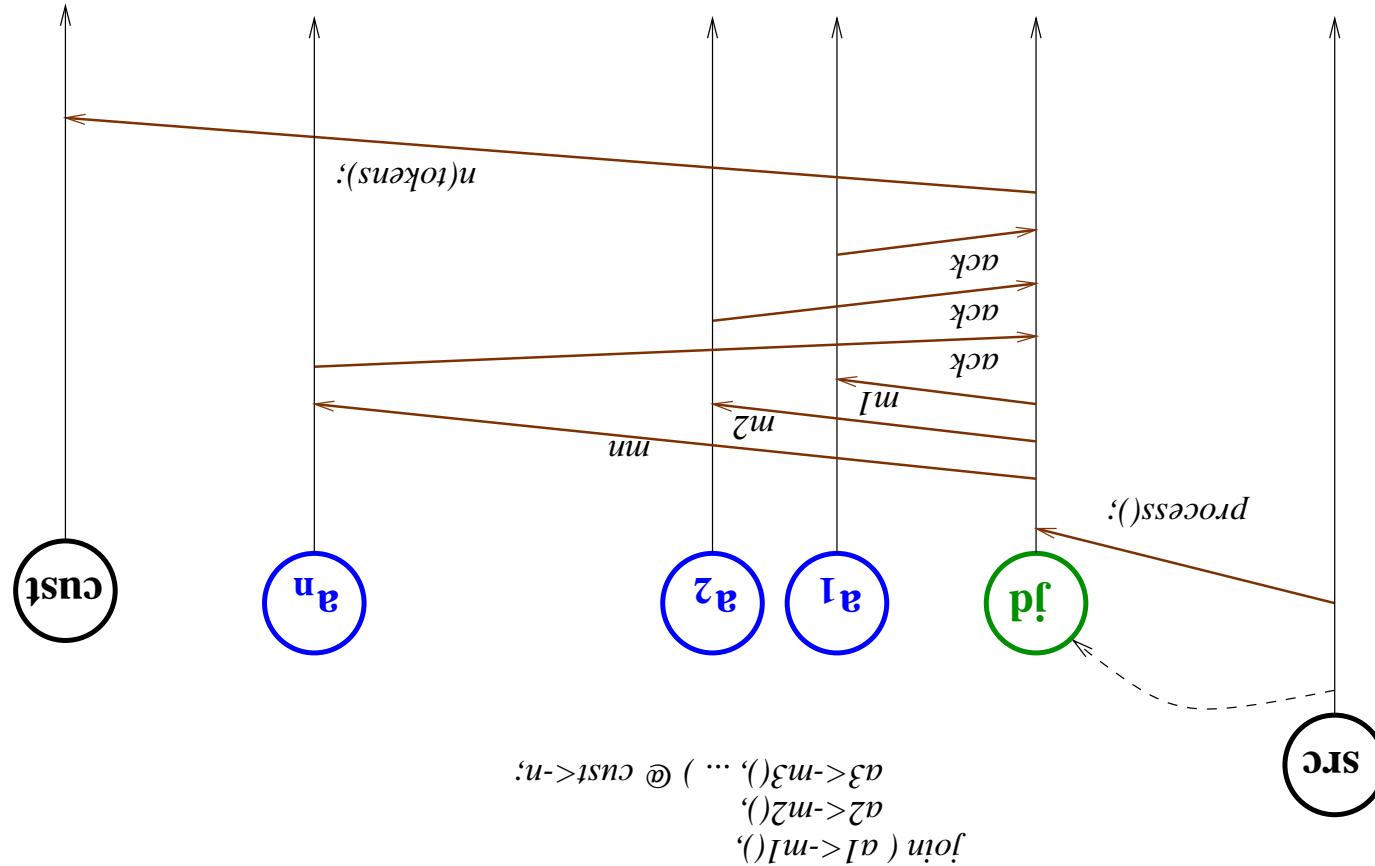
Group knowledge multicast (Fagin, Halpern, Moses and Vardi, 1995)



Lines of Code Comparison for Multicast Protocols

SALSA Actor Library





Join Continuation Code Generation