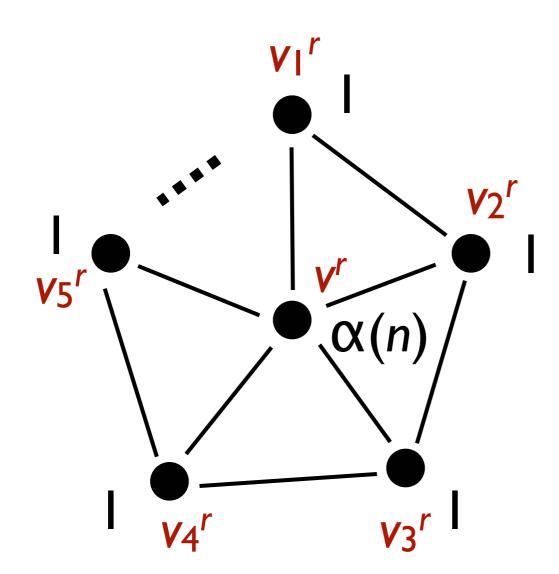
#### Piecewise Smooth Surface Reconstruction

Hugues Hoppe et al.

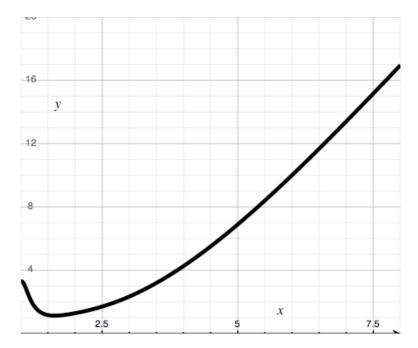
## Loop's Subdivision Scheme

- A mesh *M* is iteratively refined into smoother meshes *M*<sup>r</sup>
- Each vertex in M<sup>r+1</sup> is a weighted average of "nearby" vertices in M<sup>r</sup>
- Weights are given by masks

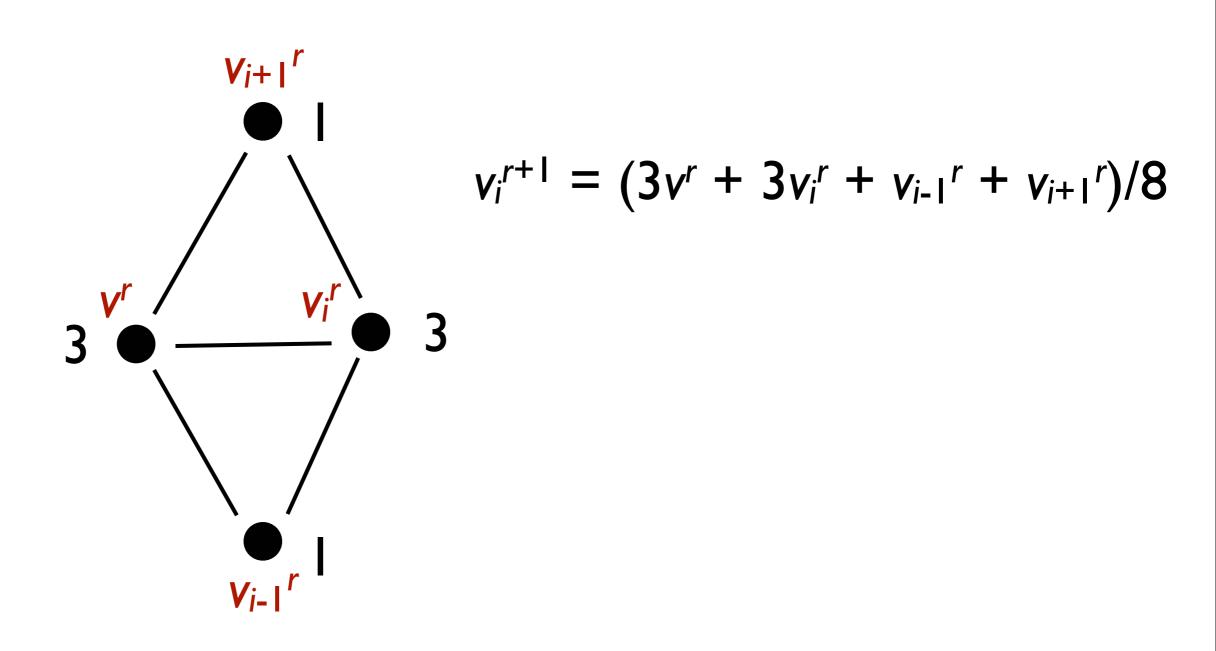
#### Subdivision masks



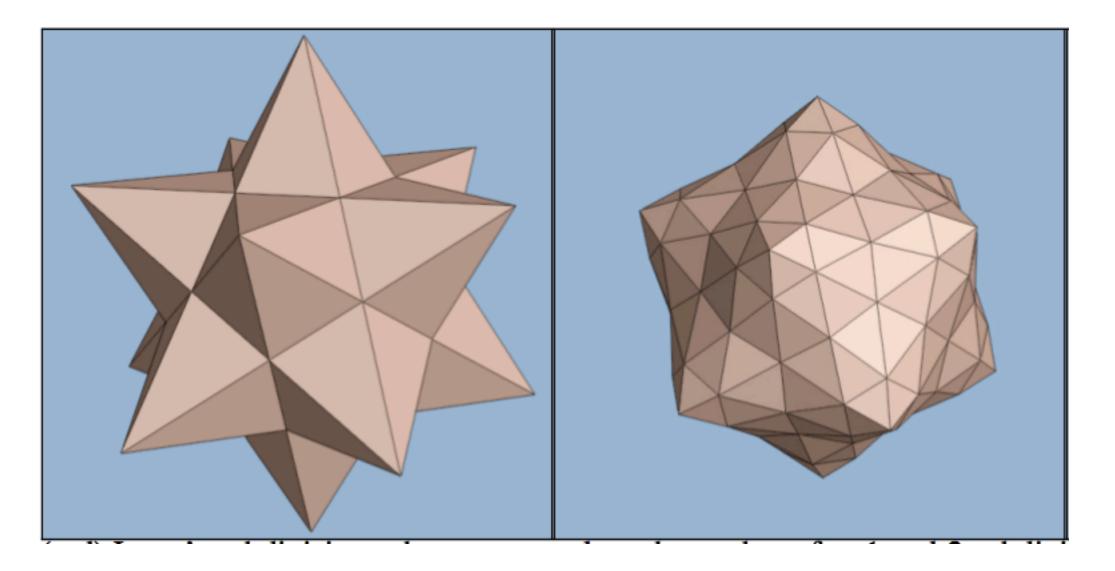
 $v^{r+1} = (\alpha(n) v^r + v_1^r + v_1^r)$  $\ldots + v_n^r)/(\alpha(n) + n)$ 



### Subdivision masks



### Subdivision masks



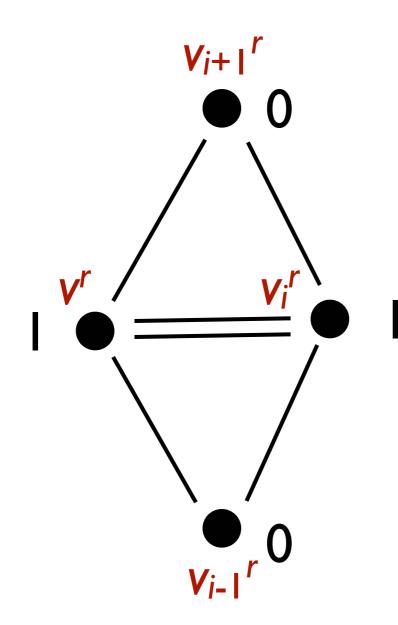
### Limits of the surface

- It turns out that we can calculate exactly the "eventual" position v<sup>∞</sup> of v using eigenanalysis (yay math!)
- We can also get exact normal vectors (e.g., for Phong shading)

### Creases

- This paper modifies Loop surfaces so you can mark edges as "creases"
- Only difference is that masks are different for points and edges on creases
  - Vertices on one side of a crease cannot affect vertices on the other side

## Crease edge mask



$$v_i^{r+1} = (v^r + v_i^r)/2$$

Crease vertices similarly "ignore" noncrease vertices

### The Problem

- Given a set V of vertices, find a mesh M which, when used as a Loop surface:
  - is concise (few control vertices)
  - has few crease edges
  - minimizes the "distance" from V to  $M^{\infty}$

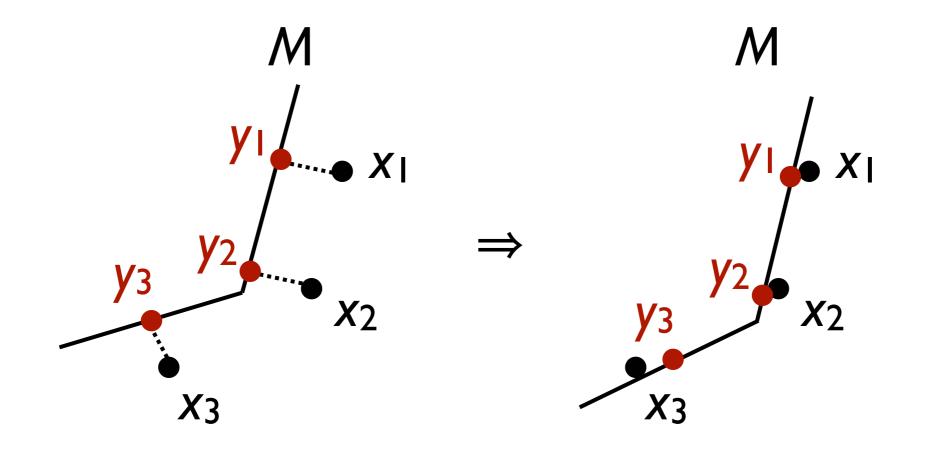
# Energy function

- $E(M, V) = E_{dist}(M, V) + c_{rep}m + c_{sharp}e$ 
  - *m* is number of vertices in *M*
  - e is number of crease edges in M

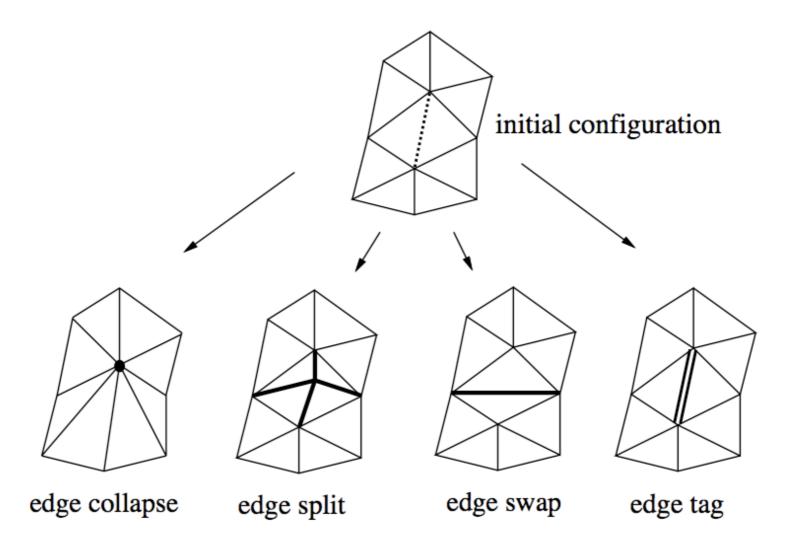
# Mesh Optimization

 If we keep the connectivity of the mesh constant and move the vertices, it turns out to be an iterative least-squares problem

### Mesh Optimization



# Mesh Optimization



## Pretty Pictures