



Detecting Conversing Groups of Chatters: A Model, Algorithm and Tests

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Motivation and Problem

- Internet chatrooms are open for exploitation by malicious users
 - Chatrooms are open forums which offer anonymity.
 - The real identity of participants are decoupled from their chatroom nicknames.
 - Multiple threads of communication can co-exist concurrently.
- Our goal is
 - to provide automated tools to study chatrooms
 - to discover who is chatting with whom?
 - Human monitoring is possible but not scalable.

Motivation and Problem (cont.)

■ Not a trivial task even for a well trained human eye

[20:19:40] <id1> what u been up to or down to?

[20:19:42] <id2> Hi!! Anyone from Boston around?

[20:20:00] <id3> amenazas d eataque

[20:20:04] <id3> sorry

[20:20:29] <id4> laying around sick all weekend...

[20:20:30] <id1> can anyone in here speak english??

[20:20:37] <id4> what about you ?

[20:20:40] <id1> whats wrong?

[20:20:46] <id5> not me

[20:20:48] <id4> sinus infection

[20:20:57] <id1> i had that last week

[20:20:58] <id6> Hmm, those seem rather infectious down there

[20:21:07] <id1> its this stupid weather

[20:21:32] <id4> yeah...darn weather

[20:21:43] <id1> i wont get good and over them until we get a good rain or a really hard freeze

Outline

- Related work
- Contributions
- The Model
- Algorithms
 - Cluster
 - Connect
 - Color & Merge
- Results
- Conclusion

IRC - Internet Relay Chat

- RFC 2810 ... 2813
- Interactive and public forum of communication for participants with diverse objectives.
- IRC is a multi-user, multi-channel and multi-server chat system which runs on a Network.
- It is a protocol for text based conferencing:
 - Provides people all over the world to talk to one another in real time.
 - Conversation or chat takes place either in private or on a public channel.

Related Work

- **Multi-users in open forums**
 - H.-C. Chen, M. Goldberg, M. Magdon-Ismael, “Identifying multi-users in open forums.” ISI 04.
- **An automated surveillance system**
 - S. A. Camtepe, M. S. Krishnamorthy, B. Yener, “A tool for Internet chatroom surveillance”, ISI 04.
- **PieSpy**
 - P. Mutton and J. Golbeck, “Visualization of Semantic Metadata and Ontologies”, IV03.
 - P. Mutton, “PieSpy Social Network Bot”, <http://www.jibble.org/piespy/>.
- **Chat Circle**
 - F. B. Viegas and J. S. Donath, "Chat Circles“, CHI 1999.
- **Social Network Analysis (SNA)**
 - V. Krebs, “An Introduction to Social Network Analysis”, <http://www.orgnet.com/sna.html>.

Contributions

- A model which does not use semantic information,
 - chatters are nodes in a graph,
 - collection of chatters is a hyperedge,
- Two efficient algorithms
 - Uses statistical information on the posts to create candidate hyperedges,
 - “Cleans” the hyperedges using:
 - Transitivity,
 - Graph coloring,
- Algorithms are rigorously tested using simulation on the model.

The Model - Assumptions

- We model a single chatroom which corresponds to a topic.
- Members form small groups and talk on one or more subtopics:
 - Subtopics are created at the beginning and never halts.
- A user participates in one subtopic only. A user:
 - arrives,
 - selects a subtopic to talk on,
 - stays in the same subtopic during his/her lifetime.
- At any time, only one user is selected to post in a subtopic
- Message interarrival times are random according to a given distribution.

The Model – Assumptions (cont.)

- User's arrival and departure times are selected uniformly at random. To make a user to post enough messages for analysis:
 - Simulation time is divided into “**n**” **regions**
 - Arrival times are selected uniformly at random from the first region,
 - Departure times are selected uniformly at random from the last region.
- At any time, messages coming from all subtopics are uniformly at random shuffled and output.

The Model - Parameters

- Simulation time and number of regions
- Number of users
- Number of subtopics
- Probability distribution and parameters (mean, variance, ...) for:
 - User to subtopic assignment
 - Message interarrival time
- Step size K (will be defined in the next slide)

The Model - Algorithm

- Single event queue
 - Message post events (post, user, subtopic, time)
 - User join events (join, user, subtopic, time)
 - User leave events (leave, user, subtopic, time)
- K-step posting probability for each subtopic
 - A list of size **K** named as “Probability History List”
 - A user who post recently is pushed to front
 - A user at the front has smallest probability of post next
 - A user at the end and users not in the list have the highest probability of post next

The Model - Algorithm (cont.)

- Load parameters
- For each user
 - select an arrival time, generate an arrival event for the user
 - select a departure time, generate a departure event for the user
 - select a subtopic, generate join event
- For each timestep
 - For each events of current time
 - If post event
 - insert the message to message buffer
 - create new post event
 - select next user to send according to K-step probability
 - select time for next post (message interarrival time)
 - update K-step probability (probability history list)
 - If join event
 - add user to subtopic
 - If first user in the subtopic,
 - create a post event
 - update K-step probability (probability history list)
 - If leave event
 - remove user from subtopic
 - Shuffle the message buffer
 - Log the messages to a file

The Model - Output

- Sample chat log

TIME 6 USER 20

TIME 7 USER 15

TIME 9 USER 61

TIME 12 USER 41

TIME 12 USER 24

.....

- User to subtopic assignments

Subtopic	Members
1	15
2	20,41
3	61
4	24

Algorithms

- Initial processing of message logs
 - Consider every consecutive messages
 - Generate list of node pairs and the corresponding interarrival times

Sample Log	Node-pair, Interarrival list
TIME 6 USER 20	users (20,15) int. time 1
TIME 7 USER 15	
TIME 9 USER 61	users (15,61) int. time 2
TIME 12 USER 41	users (61,41) int. time 3
TIME 12 USER 24	users (41,24) int. time 0

Algorithms – Kmeans

- Simple Clustering (K-Means)
 - K-means clustering algorithm is applied on Interarrival list
 - Generates two clusters
 - Red: pairs which has small interarrival times are put into this cluster
 - Blue: pairs which has large interarrival times are put into this cluster

Algorithms – Kmeans (cont.)

- Simple Clustering (K-Means)
 - K-means clustering on Interarrival list
 - Generates two clusters:
 - Red: pairs which has small interarrival times
 - Blue: pairs which has large interarrival times
 - Declares:
 - Red pairs as not engaged in conversation
 - Blue pairs as engaged in conversation
 - Idea: interarrival time between messages of two users, who exchanges messages over a subtopic, can not be smaller then a threshold:
 - It takes time for user to read, interpret , prepare answer
 - Network and servers introduce additional latency

Algorithms – Kmeans (cont.)

- Issues:
 - Incomplete, it does not identify members of sub topics (conversing groups)
 - May include contradictory information
 - For group of three users a,b,c
 - (a,b) and (a,c) are blue, (b,c) is red
 - Are (a,b,c) in the same subgroup???
- Algorithms ***Connect*** and ***Color_and_Merge***
 - Reconcile possible contradictions
 - Produce complete output

Algorithms – Connect

- Takes blue and red clusters
- Trusts blue cluster
- Considers blue cluster as the edge set of a graph **B**
 - Finds connected components in **B**
 - breath-first search on **B**
 - Consider previous example
 - For group of three users a,b,c
 - (a,b) and (a,c) are blue, (b,c) is red
 - Connect concludes that (a,b,c) are in the same subgroup.

Algorithms – Color

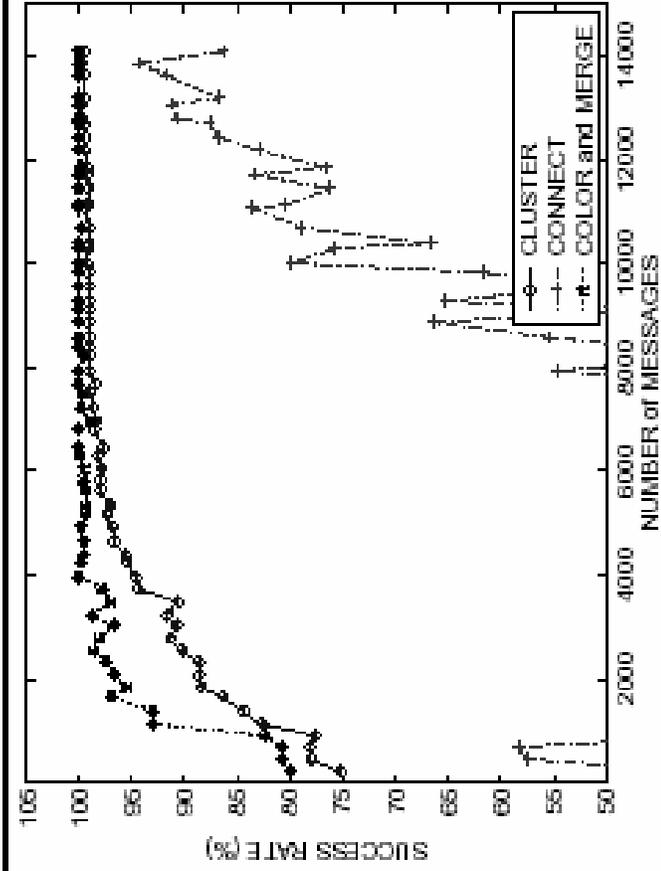
- Takes blue and red clusters
- Trusts red cluster more than blue cluster
- Considers red cluster as the edge set of a graph **R**
 - Applies vertex coloring
 - Uses heuristic Greedy to find an approximate solution
 - Generates color classes

Algorithms – Merge

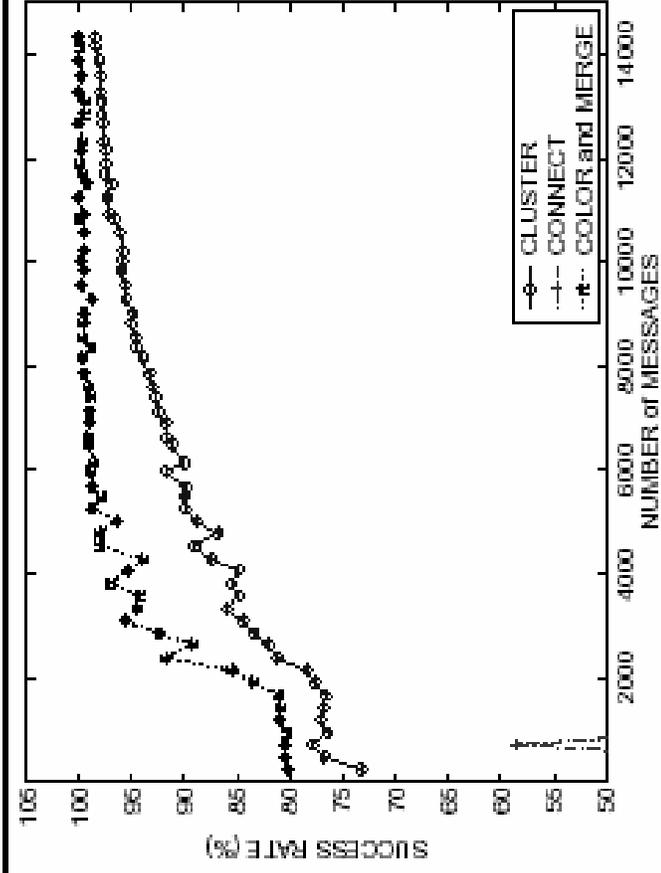
- Takes color classes generated by **color**
- For each pair of color classes C_1 and C_2
 - e_b = number of user pairs (x, y) where
 - (x, y) in blue cluster
 - $(x \text{ in } C_1 \text{ and } y \text{ in } C_2)$ or $(y \text{ in } C_1 \text{ and } x \text{ in } C_2)$
 - If $(e_b/|C_1| \cdot |C_2| \geq \text{threshold})$ merge C_1 and C_2
 - For our model, we found that threshold 0.7 gives good results
- Announce final color classes as subtopics.

Tests

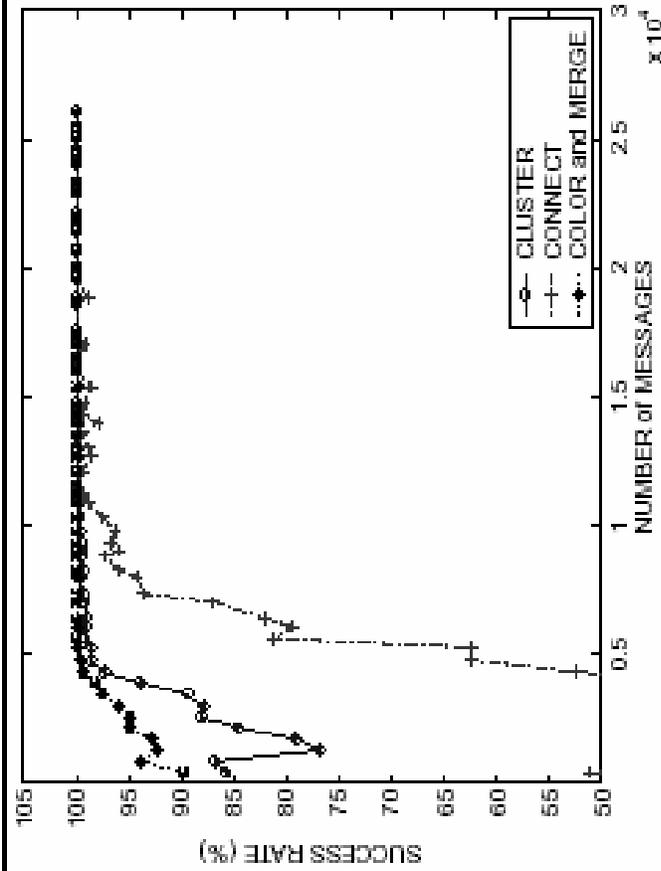
- Parameters of the model are tuned according to observations and statistical analysis over real chatroom logs.
- A user pair which is announced correctly as being in the same subtopic is accepted as a correct result
- Success rate = # correct results / all
- Following slide lists results for:
 - 5 topics, 50 users
 - 5 topics, 75 users
 - 10 topics, 50 users
 - 10 topics, 75 users



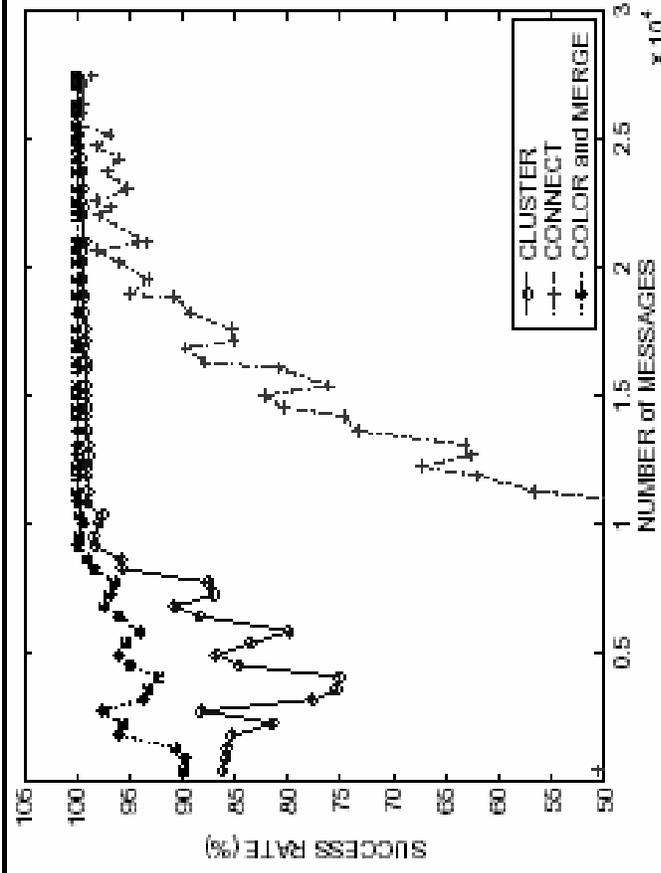
(a) 5 Topics and 50 Users



(b) 5 Topics and 75 Users



(c) 10 Topics and 50 Users



(a) 10 Topics and 75 Users

Results

- For sufficiently long log size, all algorithms converges to 100% success
- Critical factor is number messages per user.
 - As the number of users increases, larger logs are required
- Color_and_Merge algorithm provides the best result.
 - Converges to 100% success very quickly
- Connect is the most sensitive algorithm to log size
- As the log size decrease, connect fails faster
 - Why? A single false edge may connect two components yielding too much false results

Conclusion

- We presented a model for which we showed that it is possible to accurately determine the conversation
- Ideas can be generalized to more elaborate models
- Future work:
 - Enhance the model
 - Users may belong to multiple conversations
 - Users may switch between conversations
 - Apply algorithms to real chatroom logs