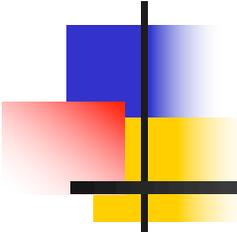


Announcements

- Accommodations and zone assignment out tomorrow
- Rules same as for Exam 1 with some tweaks; announcement tomorrow on Submittity
- Makeup date (tentative): Wed Nov.12, 5-8pm

- Homework
 - Write your contracts
 - Adhere to syntax and functionality restrictions



Exam 2 Topics

Topics

- Scheme (Lectures 12 and 13, plus chapters)
 - S-expression syntax
 - Lists and recursion
 - Shallow and deep recursion
 - Equality *eq? equal?*
 - Higher-order functions
 - **map, foldl, and foldr** *foldl: id (e1) e2 .. en)*
id1 e2
 - Programming with **map, foldl, and foldr** *foldr: (e1 e2 .. en-1 (en) id)*
(en-1 id1)
...
id → *total result of fold*
 - Tail recursion

Topics

- Scheme (Lecture 14, plus chapters)
 - Binding with **let**, **let***, **letrec**
 - Scoping in Scheme
 - Closures and closure bindings

Topics

- Scoping, revisited (Lecture 14, plus chapters)
 - Static scoping
 - Reference environment
 - Functions as third-class values vs.
 - Functions as first-class values
 - Dynamic scoping
 - With shallow binding
 - With deep binding

Topics

- Lambda calculus (Lectures 15, 16, and 17)
 - Syntax and semantics
 - Free and bound variables
 - Substitution
 - Rules of the Lambda calculus: Alpha-conversion and Beta-reduction
 - Normal forms *NF, HNF, WHNF*
 - Reduction strategies: Normal order and Applicative order
 - Fixed-point combinator and recursion

Quiz 5

Question 1. (2pts) Scheme's scoping discipline is

Select one:

- static scoping
- dynamic scoping

Question 2. (2pts) Scheme's typing discipline is

Select one:

- static typing
- dynamic typing

Quiz 5

Questions 3 and 4 refer to the following Scheme function:

```
(define (fun a b)
  (cond ((= a b) a)
        ((> a b) (fun (- a b) b))
        (else (fun a (- b a)))))
```

Question 3. (2pts) What does fun compute? Note: you may assume that the arguments are positive integers.

GCD

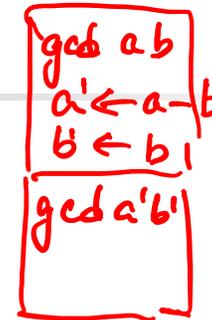
Question 4. (2pts) fun is tail-recursive.

Select one:

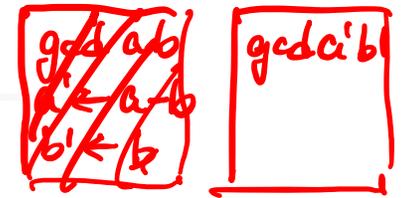
true

false

normal:



tail optimization:



Quiz 5

*(define (atom? x)
 (not (pair? x)))*

Question 5. (2pts) Function `atomcount`, defined below, attempts to count the number of atoms nested in a list. Recall predicate `atom?` that we wrote in class — it returns true if given an object that is *not a pair* (i.e., an object we cannot take the `car` or the `cdr` of).

```
(define (atomcount lis)
  (cond ((atom? lis) 1)
        ((null? lis) 0)
        (else (+ (atomcount (car lis)) (atomcount (cdr lis))))))
```

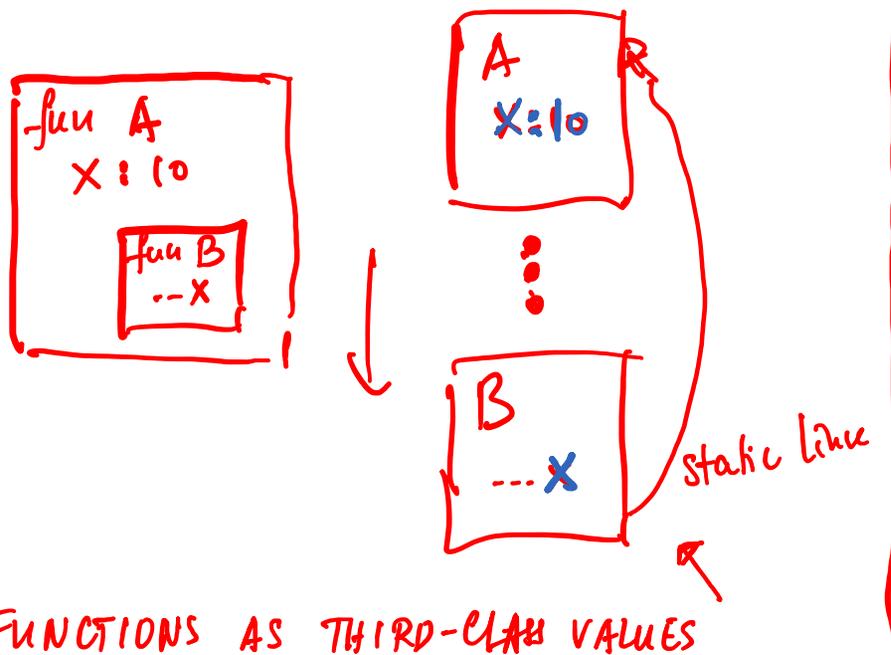
`(atomcount '(1 (2 (3))))` yields

6

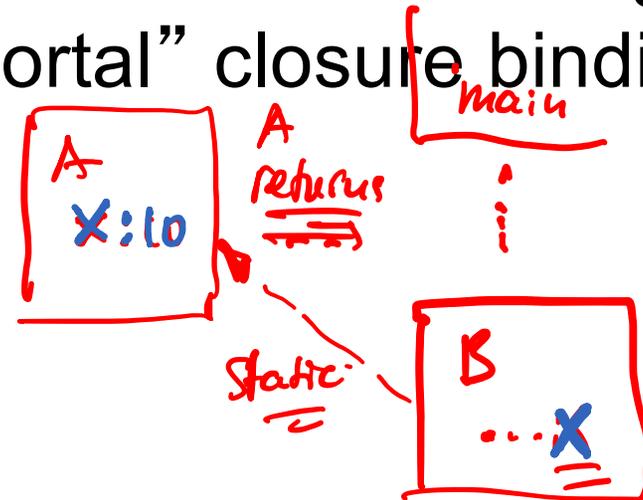
*(atom? 'a) → #t
(atom? '(a)) → #f
(atom? '()) → #t*

Scoping with First-Class Functions

- When functions are **first-class values**, static scoping gets more involved
 - Function value may outlive static referencing environment! Need “immortal” closure bindings



FUNCTIONS AS THIRD-CLASS VALUES



FUNCTIONS AS FIRST-CLASS VALUES.
A returns but B needs A.

Scoping with First-Class Functions

- Dynamic scoping gets more involved as well
 - Shallow binding vs. deep binding
- **Dynamic scoping with shallow binding**
 - Reference environment for function/routine is not created until the function is called
 - All non-local references are resolved using the most-recent-frame-on-stack rule --- just follow dynamic links

Scoping with First-Class Functions

- **Dynamic scoping with deep binding**
 - When a function/routine is passed as an argument, the code that passes the function/routine has a particular reference environment (the current one!) in mind. It passes this reference environment along with the function value (it passes a closure).

Example

$v : \text{integer} := 10$

people : database

print_routine (p : person)

if p.age > v

write_person(p)

other_routine (db : database, P : procedure)

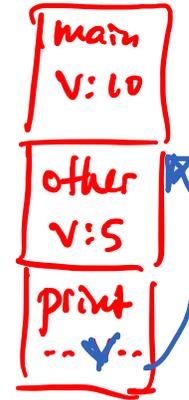
$v : \text{integer} := 5$

foreach record r in db

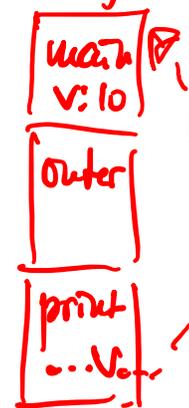
P(r)

other_routine(people, print_routine) /* call in main */

Shallow :



deep :



P is a closure } print-routine + main's env.

WITH DEEP BINDING:
main PASSES A CLOSURE

} print-routine + main's ENVIRON.

Quiz 6

Questions 1-3 below refer to the following Scheme function:

```
(define A
  (lambda ()
    (let* ((x 2)
           (C (lambda (P) (let ((x 4)) (P) )))
           (D (lambda () x))
           (B (lambda () (let ((x 3)) (C D))))))
      (B))))
```

with shallow: D is just $\lambda().x$
 with deep: $\left\{ \begin{array}{l} \lambda().x \\ \text{ref env of B} \end{array} \right.$

Question 1. (1pts) What gets printed when we call the function in the interpreter?

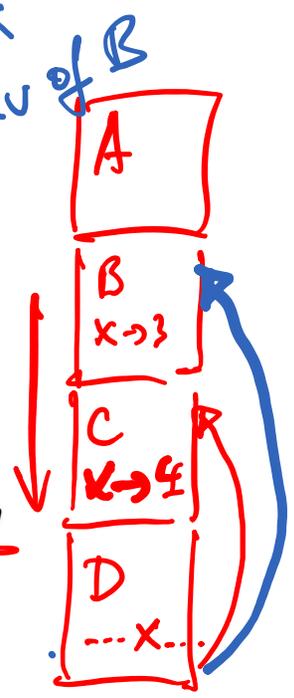
2 (A) D is "constant 2" function

Question 2. (2pts) What would get printed if Scheme used *dynamic scoping with shallow binding*?

4 D is just $\lambda().x$ function.
 x is resolved following dynamic link

Question 3. (2pts) What would get printed if Scheme used *dynamic scoping with deep binding*?

3 D is a closure $\left\{ \begin{array}{l} \lambda().x \\ + \\ \text{env. of B} \end{array} \right.$



Quiz 6

Question 4. (2pts) Willy Wazoo wrote a function f. What does it do?

```
(define (f lis)
  (foldl (lambda (x y) (if (and (number? x) (> x 5)) (cons x y) y)) lis '()))
```

partial result →
next element →

~~_____~~

Select one:

- selects all numerical elements in lis greater than 5
- selects all numerical elements in lis less or equal than 5
- neither; Willy's function has a bug

Quiz 6

Question 5. (1pt) Consider the lambda term $(\lambda v.v) ((\lambda v.v) (\lambda x.(\lambda x.x) z))$. There are this many reducible expressions in this term: 3

3

Question 6. (1pt) Is lambda term x $((\lambda y.y) z)$ in Weak Head Normal Form (WHNF)?

YES

Question 7. (1pt) Is term x $((\lambda y.y) z)$ in Normal Form (NF)?

No

Questions?

$(\text{cons } 'a \ 10)$ \rightarrow  $(a \ 10)$

$(\text{list } 'a \ 10)$ \rightarrow  $(a \ 10)$

\parallel
 $(\text{cons } 'a \ (\text{cons } 10 \ '()))$ \rightarrow  $(a \ 10 \ ())$

The End
