

Logic Programming and Prolog (conclusion)

Finish reading: Scott, Chapter 12

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Lecture Outline

- *Prolog programming*
 - *Parsing and language processing*
 - *Arithmetic*

- *Imperative control flow*
- *Negation by failure*
- *Generate and test paradigm*

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Negation by Failure: `not (X)` , `\+ (X)`

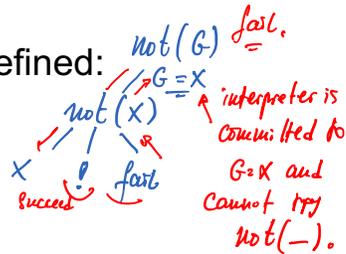
- `not (G)` succeeds when goal `G` cannot be made to succeed

- Called **negation by failure**, defined:

`not (X) :- X,!,fail.`

`not (_).`

→ ALWAYS TRUE.



- Not the same as negation in logic $\neg G$!
- In Prolog, we can assert that something is true, but we **cannot assert that something is false**

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Exercise

`takes(jane, his).`

`takes(jane, cs).`

`takes(ajit, art).`

`takes(ajit, cs).`

`classmates(X,Y) :- takes(X,Z),takes(Y,Z).`

`?- classmates(jane,Y).`

What are the bindings of `Y`?

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Exercise

- `p(X) :- q(X), not(r(X)).`
`r(X) :- w(X), not(s(X)).`
`q(a). q(b). q(c).`
`s(a). s(c).`
`w(a). w(b).`

- Evaluate:

- `?- p(a).`
- `?- p(b).`
- `?- p(c).`

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Generate and Test Paradigm

- Search in space
- Prolog rules to **generate** potential solutions
- Prolog rules to **test** potential solutions for desired properties
- Easy prototyping of search
`solve(P) :- generate(P), test(P).`

A Classical Example: n Queens

- Given an **n** by **n** chessboard, place each of **n** queens on the board so that no queen can attack another in one move
 - Queens can move either vertically,
 - horizontally, or
 - diagonally.
- A classical **generate and test** problem

n Queens

```
my_not(X):- X, !, fail.  %same as not
my_not(_).
in(H,[H|_]).           %same as member
in(H,[_|T]):- in(H,T).
```

```
nums(H,H,[H]).
nums(L,H,[L|R]):- L<H, N is L+1, nums(N,H,R).
```

```
queen_no(4).
```

n Queens (ii)

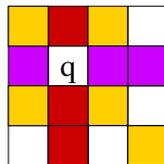
```
ranks(L):- queen_no(N), nums(1,N,L).
files(L):- queen_no(N), nums(1,N,L).
```

```
rank(R):- ranks(L), in(R,L).
```

```
file(F):- files(L), in(F,L).
```

n Queens (iii)

```
attacks((R,_),(R,_)).
attacks( (_,F), (_,F)). %a Prolog tuple
attacks( (R1,F1), (R2,F2)):-
    diagonal((R1,F1),(R2,F2)).
%%%can decompose a Prolog tuple by unification
(X,Y)=(1,2) results in X=1,Y=2; tuples have fixed
size and there is not head-tail type construct for
tuples
```



same rank
same file
same diagonal

What is safe placement
for next queen on board?

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n Queens (iv)

```
%%% Two squares are on the same diagonal if the slope of
the line between them is 1 or -1. Since / is used, real
number values for 1 and -1 are needed.
```

```
diagonal((X,Y),(X,Y)). %degenerate case
diagonal((X1,Y1),(X2,Y2)):-N is Y2-Y1,D is X2-X1,
    Q is N/D, Q is 1 . %diagonal needs bound
arguments!
diagonal((X1,Y1),(X2,Y2)):-N is Y2-Y1,D is X2-X1,
    Q is N/D, Q is -1 .
%%%because of use of "is", diagonal is NOT invertible.
```

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n Queens (v)

```
%%% This solution works by generating every list of
squares, such that the length of the list is the same
as the number of queens, and then checks every list
generated to see if it represents a valid placement of
queens to solve the N queens problem;
assume list length function
```

```
queens(P):- queen_no(N), length(P,N),
placement(P), ok_place(P).
```

“generate” code given first “test” code follows

n Queens (vi)

```
placement([]).
placement([(R,F)|P]):- placement(P), rank(R), file(F).
```

n Queens (vii)

```
ok_place([]).  
ok_place([(R,F)|P]):- no_attacks((R,F),P),ok_place(P).
```

```
no_attacks(_,[]).  
no_attacks((R,F),[(R2,F2)|P]):-  
    my_not(attacks((R,F),(R2,F2))), no_attacks((R,F),P).
```

The End
