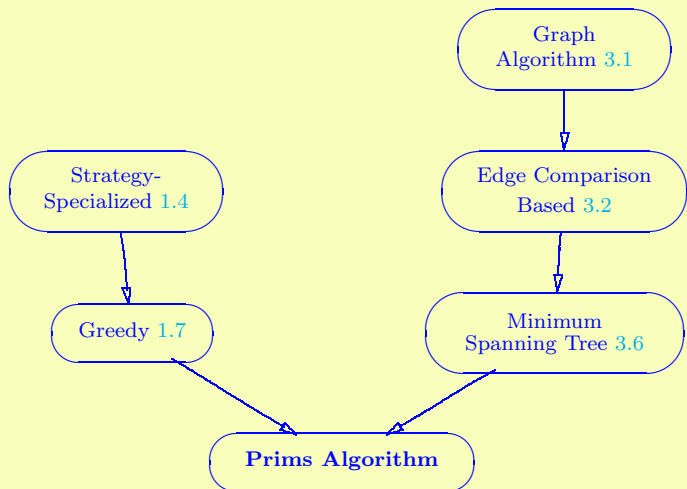


3.3. Prim's Minimum Spanning Tree

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Refinement of: Greedy (§1.7) Minimum Spanning Tree (§3.6). Prim's algorithm is a greedy algorithm for solving the Minimum spanning tree problem. It is closely related to Dijkstra's (§3.16) single source shortest path algorithm.

Prototype:

```
template <class VertexListGraph,  
         class DijkstraVisitor, class PredecessorMap,  
         class DistanceMap, class WeightMap,  
         class IndexMap>  
void prim_minimum_spanning_tree(  
    const VertexListGraph& g,  
    vertex_descriptor s, PredecessorMap pred,  
    DistanceMap distance, WeightMap weight,  
    IndexMap index_map)
```

Input: Directed or undirected graph G , a model of VERTEXLISTGRAPH, with edge set E and vertex set V .

Start node s .

WEIGHTMAP, which maps each $e \in E$ to real values.

Empty PREDECESSORMAP, which will map each $v_1 \in V$ to another $v_2 \in V$, where v_2 corresponds to the parent of v_1 in the minimum spanning tree.

INDEXMAP, which maps each $v \in V$ to an integer in the range $[0, \text{num_vertices}(g))$.

Empty DISTANCEMAP, which will map each $v \in V$ to its shortest path weight from source vertex s in the minimum spanning tree. [1]

Output: Modified PREDECESSORMAP, which maps each $v_1 \in V$ to another $v_2 \in V$, where v_2 corresponds to the parent of v_1 in the minimum spanning tree, if one exists.

A Modified DISTANCEMAP, which maps each $v \in V$ to its shortest path weight from source vertex s in the minimum spanning tree.

Effects: No effects on the input.

Asymptotic complexity: Let E = number edges in the graph and let V = number of vertices in the graph.

- Best case (sparse graph where $E \lesssim V$):
 $O(V \log_2 V)$
- Average case: $O(E \log_2 V)$
- Worst case (dense graph where $E \approx V^2$):
 $O(V^2 \log_2 V)$

| V | E | Assignments | Iterator | Integer | Compare |
|-----|------|-------------|----------|---------|---------|
| 100 | 100 | 414 | 514 | 400 | 102 |
| 100 | 200 | 2914 | 1800 | 1125 | 980 |
| 100 | 400 | 5150 | 3288 | 2287 | 1894 |
| 100 | 800 | 6946 | 5710 | 4301 | 2676 |
| 100 | 1600 | 9772 | 10510 | 8301 | 3928 |
| 100 | 3200 | 14648 | 20110 | 16301 | 6263 |
| 100 | 6400 | 23270 | 39310 | 32301 | 10716 |

| V | E | Assignments | Iterator | Integer | Compare |
|------|-----|-------------|----------|---------|---------|
| 100 | 100 | 414 | 514 | 400 | 102 |
| 200 | 100 | 964 | 1104 | 836 | 232 |
| 400 | 100 | 1614 | 2014 | 1600 | 402 |
| 800 | 100 | 3214 | 4014 | 3200 | 802 |
| 1600 | 100 | 6414 | 8014 | 6400 | 1602 |
| 3200 | 100 | 12848 | 16034 | 12808 | 3209 |
| 6400 | 100 | 25614 | 32014 | 25600 | 6402 |

Operation Counts in the Average Case:

Value comparisons: $0.27E \log_2 V + 0.08V \log_2 V$

Value assignments: $0.6E \log_2 V + 0.31V \log_2 V$

Iterator operations: $0.93E \log_2 V + 0.39V \log_2 V$

Integer operations: $0.76E \log_2 V + 0.31V \log_2 V$

Algorithm Animation: An animation for Prim's MST algorithm can be found both [here](#).