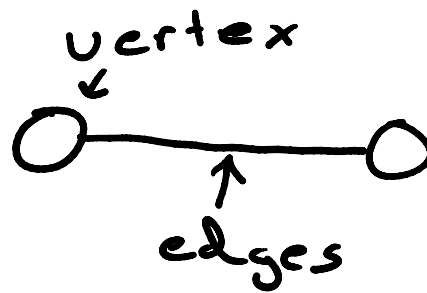
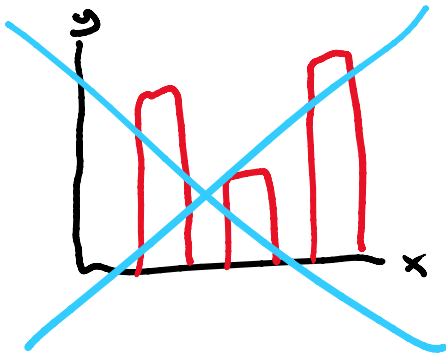


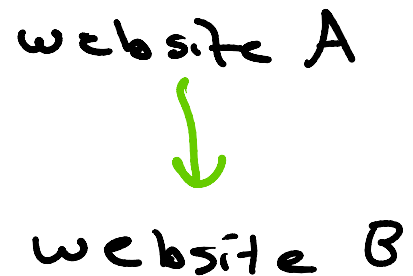
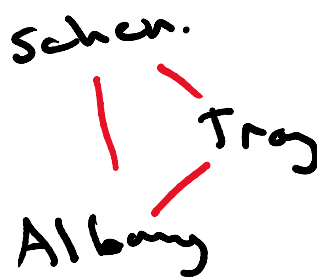
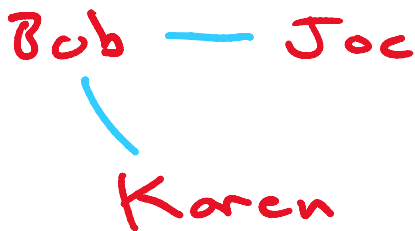
What is a graph?



Real-world graph

social networks

Road networks Info. nets.



Properties of real-world graphs

- Sparsity  $\rightarrow |E| \ll |V|^2$

- Skew  $\rightarrow$  # low degree verts  
 $\gg$

# high degree verts

- Hubs  $\rightarrow$  large degree or otherwise important vertices

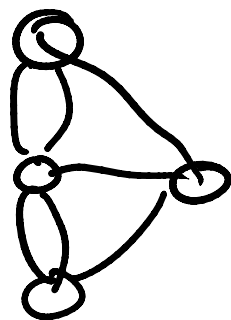
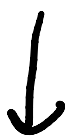
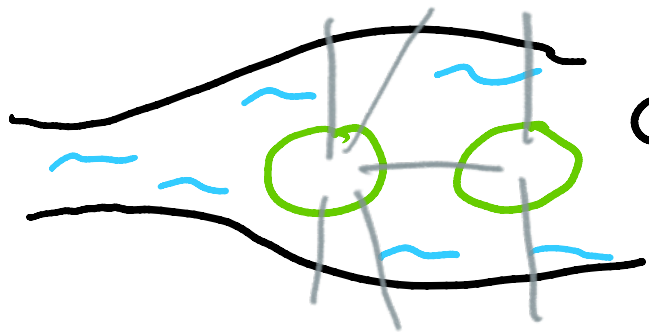
important vertices

- Irregularity  $\rightarrow$  not necessarily physically constrained
  - Small-world  $\rightarrow$  avg. shortest paths length is small
- "6 degrees of Kevin Bacon"

---

## History of Graph Theory

Königsberg



Euler: ? ? ?  
? ? ? ?

Q: Can I start at one point, traverse each bridge exactly once, and return to my start?

A: Invent graph theory  
"O.G. of G.T."

Ap2: Define properties of a graph for an

of a graph for an  
Euler Tour

---

## Basic Definitions

Graph  $G$ : tuple of vertices  $V(G)$   
and edges  $E(G)$

$$G = \{V(G), E(G)\}$$

for each  $e \in E(G)$ ,  $e \rightarrow (u, v)$   
 $u, v \in V(G)$

Terminology:

$u, v$  are endpoints of  $e$

$e$  joins  $u$  and  $v$

$u, v$  are incident with  $e$

$e$  is incident on  $u, v$

$u, v$  are adjacent to each other

$u, v$  are neighbors

$$u \in N(v), v \in N(u)$$

↑  
neighborhood  
of  $u$

The degree of  $v$  is its number of incident edges

$$|N(v)| = d(v) \leftarrow \text{degree of } v$$

(only in simple graphs)

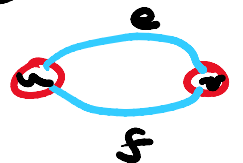
$$\text{order of } G = |V(G)| = n$$

$$\text{size of } G = |E(G)| = m$$

Multi-edges are edges in  $G$  that have the same endpoints

$$\text{e.g. } e = (u, v), f = (u, v)$$

$e, f$  are multi-edges



a self-loop in  $G$  is an edge that connects the same vertex to itself

$$\text{e.g. } e = (u, u)$$



a simple graph has no multi-edges or self-loops

a loopy graph has self-loops

... graph has multi-edges

a multi-graph has multi-edges

If  $|V|$  and  $|E| < \infty$

$\rightarrow G$  is finite

If  $|V|$  and  $|E| = 0$   $V = \emptyset$   
empty set

$\rightarrow G$  is null

If  $|V| = 1$  and  $|E| = 0$

$\rightarrow G$  is trivial

If  $|V| \geq 1$  and  $|E| = 0$

$\rightarrow G$  is empty

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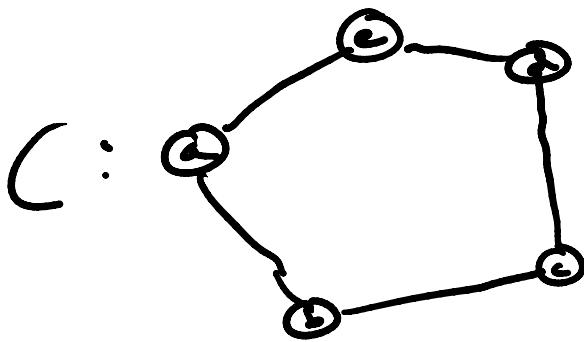
Path: a simple graph whose vertices  
can be listed s.t. any two vertices  
(such that)  
are adjacent iff they are  
(if and only if)  
consecutive in the list



P:  $\{a, b, c, d\}$

P:  $\{a, e, b, f, c, g, d\}$

Cycle: a simple graph with an equal number of vertices and edges whose vertices can be placed in a circle and two vertices are adjacent iff they appear consecutively along that circle



C:  $\{a, b, c, d, e, a\}$

Note: can also be called a closed path

Tree: a simple connected graph with no cycles

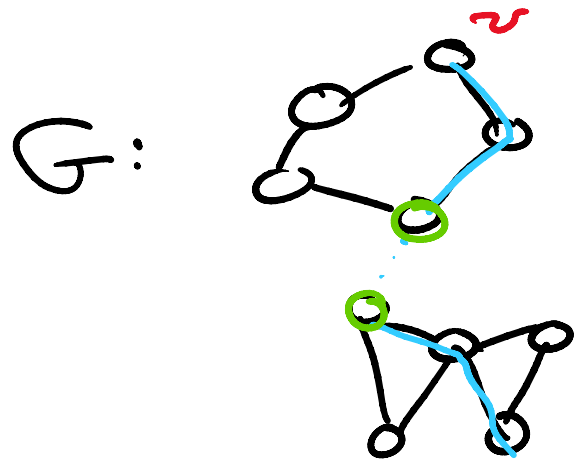
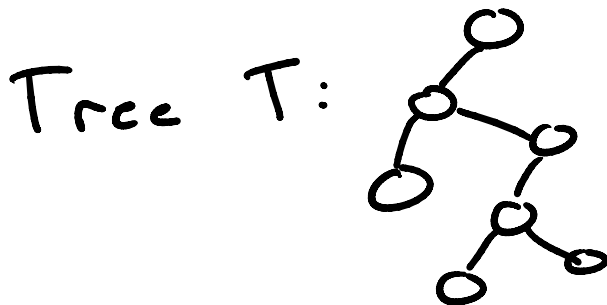
Connected graph G:  $\forall u, v \in V(G):$   
 $\leftarrow$  for all  $\leftarrow$  in

Connected graph  $G: \forall u, v \in V(G):$

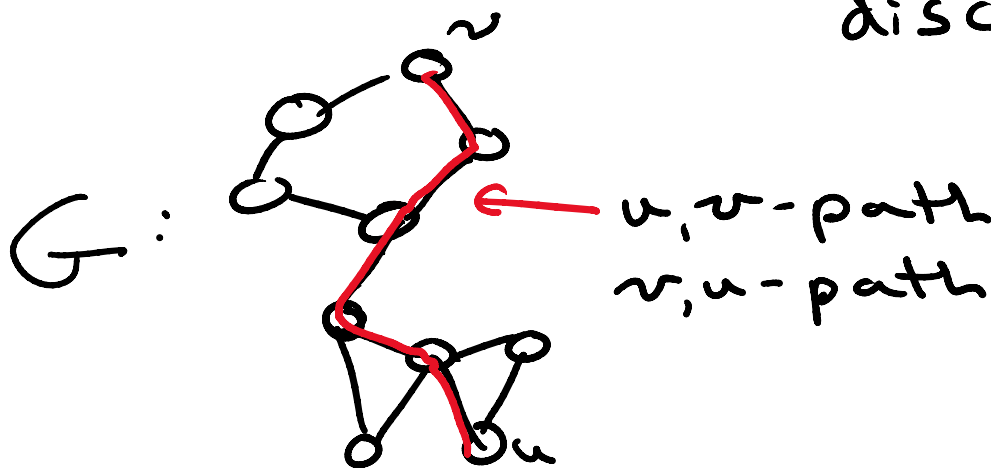
$\exists u, v$ -path

↖ there exists

$u, v$ -path: a path from  $u$  to  $v$



disconnected



$u, v$ -path  
 $v, u$ -path

connected