

Joke: What football team used the power of graph theory to tour their way to the 1960 and 1961 AFL championships?

A: Houston Oilers

PageRank

- "A modern classic."

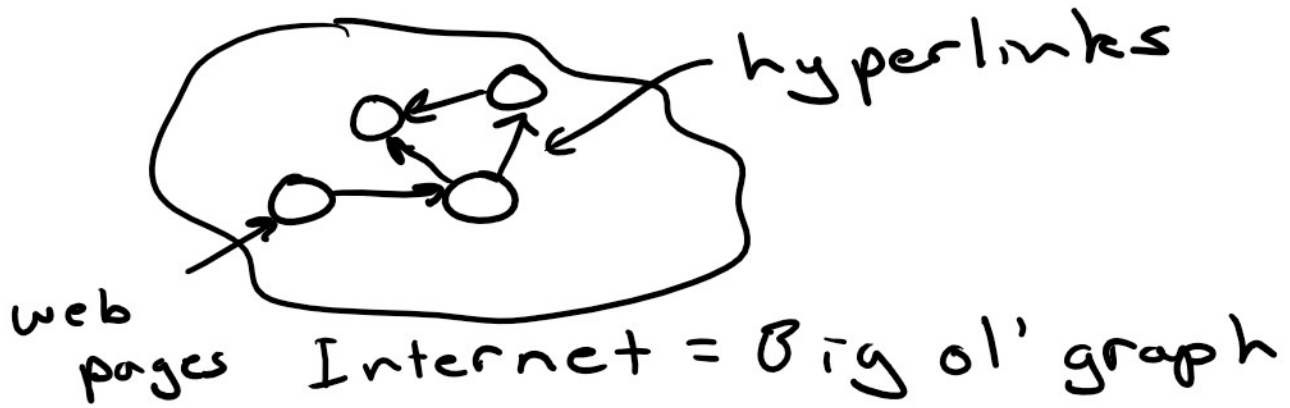
- Gittens

→ a centrality algorithm

↳ importance

BITD: a lot of internet,
but no good way
to find stuff

Google: we'll incorporate "trust" into our search



"Trust" → you're trusted if trustworthy site link to you

PageRank == trust

Random surfer model



Page Rank: probability that our random surfer is at vertex v at some point in time

Issue: what about $d^+(v) = 0$ ← sinks
 $d^-(v) = 0$

↳ randomly jump from a sink to any vertex with equal probability ← source

Graph Algorithmic Model

→ vertex centric computation over same iterations

Page Rank:

initialize $p(v) = \frac{1}{|V(G)|}$ for all $v \in V(G)$

we iterate $p(v) = \sum_{u \in N^-(v)} \frac{p(u)}{d^+(u)}$



To handle sinks / sources:

sinks: add edges to all $v \in V(G)$

In addition:

add some additional probability
of a random jump for each
step the surfer takes

(damping factor)

Linear Algebraic Model

consider our adj. matrix A

consider the diagonal degree
matrix D

- $D_{ii} = \text{out degree of vertex } i$

$$= \sum_j A_{ij}$$

- $D_{ij} = 0 \quad \forall i \neq j$

define transition probability

matrix M'

$M' = D^{-1}A$ defines transition

$M' = D^{-1}A$ ← defines transitional probabilities following out edges

$M = (D^{-1}A)^T$ ← in edge transition transition probabilities

Q: how can we calculate PageRanks

$$P_0(v) = \frac{1}{|V(G)|}$$

↑
PR on iteration zero

$$P_{i+1} = M P_i$$

eventually $P_\infty = M P_\infty$

recall for some matrix A

$$Ax = \lambda x$$

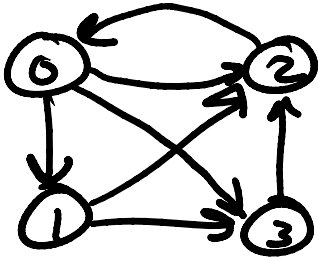
↑
eigenvector

↑
eigenvalues

↙ PageRanks are just the eigenvector of M that

eigenvector of M that corresponds to eigenvalue 1

Example calculation



$$A = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

$$D = \begin{bmatrix} 3 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$D^{-1} = \begin{bmatrix} 1/3 & \dots & \dots & 0 \\ \vdots & 1/2 & & \vdots \\ \vdots & & 1 & \\ 0 & \dots & & 1 \end{bmatrix}$$

$$M^T = \begin{bmatrix} 0 & 1/3 & 1/3 & 1/3 \\ 0 & 0 & 1/2 & 1/2 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$M = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1/2 & 0 & 0 & 0 \\ 1/3 & 1/2 & 0 & 1 \\ 1/3 & 1/2 & 0 & 0 \end{bmatrix}$$

$$P_0 = \begin{bmatrix} 1/4 \\ 1/4 \\ 1/4 \\ 1/4 \end{bmatrix}$$

$$P_1 = M P_0 = \begin{bmatrix} 1/4 \\ 1/2 \\ 11/24 \\ 5/24 \end{bmatrix}$$

Competition networks

vertices = competitors

edges = competitions

orient: add direction to
same undirected edge

For Page Rank:

- we can orient edges to victors

- PR "flows" to winners

↳ gives us a ranking of
competitors

Modifications:

- Can weight edges by the
margin of victory

Likely future special topic:
spectral clustering

spectral clustering

Graph Laplacian

$$L = D - A$$

It's spectrum contains information on how tightly coupled vertices are to each other

Application:

Graph Clustering

Graph Partitioning