

What is Graph Mining?

Monday, January 8, 2024 8:57 AM

Plan for today:

- About me and the course ✓
- Go over syllabus ✓
- Go over website ✓
- Talk about project
- Talk about paper presentation
- Lecture
 - Define graphs
 - Real-world graph properties
 - Graph Mining and applications
 - Graph computation and processing
 - NetworkX and connected components

Project details

- Groups of 1-4

- Take a look at data repos

- Prior project ideas:

* Predicting chess match outcomes using a competition network

* Vertex centrality to identify population centers on a road network

* Stock market / coin prediction

* Recommender systems using

* Recommender systems using SNAP data

- Grad students: use your ongoing project
-

Paper Presentations

- One paper selected per student (from website)
 - 20-30 minute presentation
-

Graph Definitions

Graph \rightarrow tuple of

$U(G)$ = vertex set of graph G

$E(G)$ = edge set

$W_v(G)$ = vertex weights

$W_e(G)$ = edge weights

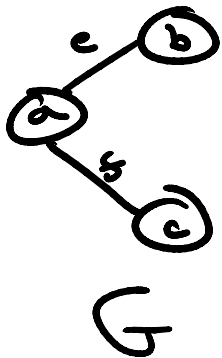
$D_v(G)$ = vertex metadata

$D_e(G)$ = edge metadata

$D_E(G) = \text{edge metadata}$

$$G = \{V, E, \omega_V, \omega_E, D_V, D_E\}$$

$G = \{V, E\} \rightarrow$ at a minimum
we need vertices and
edges to define graph



$$V(G) = \{a, b, c\}$$

$$E(G) = \{e = (a, b), f = (a, c)\}$$

For our GM context

vertices \Rightarrow represent discrete
objects, entities, etc.

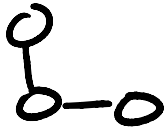
edges \Rightarrow represent connections
or interactions between
those objects, entities, etc.

weights \Rightarrow strength of connections
for edges or importance
or some other measure
for vertices.

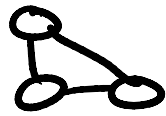
or some other measure
for vertices

metadata \Rightarrow generic catch-all
for other vertex/edge
labels or properties

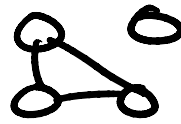
Temporal graphs \Rightarrow graphs that
evolve over time



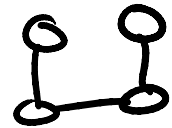
$G_{t=0}$



$G_{t=1}$



$G_{t=2}$



$G_{t=3}$

Real World Graphs

\hookrightarrow social networks, info networks,
biological networks, etc.

These graphs all have similar properties

Sparsity: $|E(G)| \ll |V(G)|^2$

Degree skew: # of low degree

$d(v) = \#$ of edges
attached to v

vertices $\gg \gg$ # of
high degree vertices

$d(v) = \#$ of edges
attached to v

vertices ...
high degree vertices




Hubs: high degree or otherwise
"important" vertices

Irregular: information, social, etc.
networks are often not
physically constrained

Small-world: average shortest path
lengths are small relative
to $|V(G)|$

"6 degrees of Kevin Bacon"

Note : Properties are
typical but not every
real-world graph has them

Graph Processing

"Vertex-centric" processing

- Each $v \in V(G)$ has same state $S(v)$
- We iteratively update $S(v)$
→ often using the states of v 's neighbors
- Many/most graph algorithms can be implemented in a vertex-centric way

Algorithmically:

Input: $G = \{V, E\}$

For all $v \in V(G)$:

$S(v) = \text{initialize}()$

For some # of iterations:

update Algo State()

For all $v \in V(G)$:

For all $u \in N(v)$:

← neighborhood of v

$S(v) = \text{update State}(S(v), S(u))$

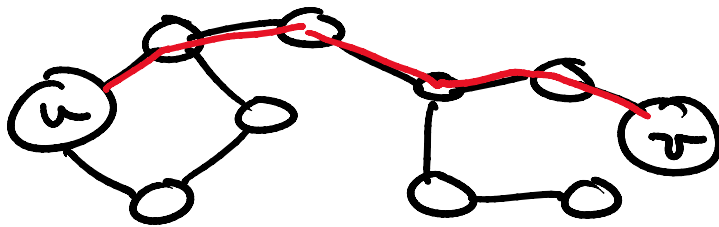
For all $v \in V(G)$:

$$S(v) = \text{finalize}()$$

Implement connectivity decomposition

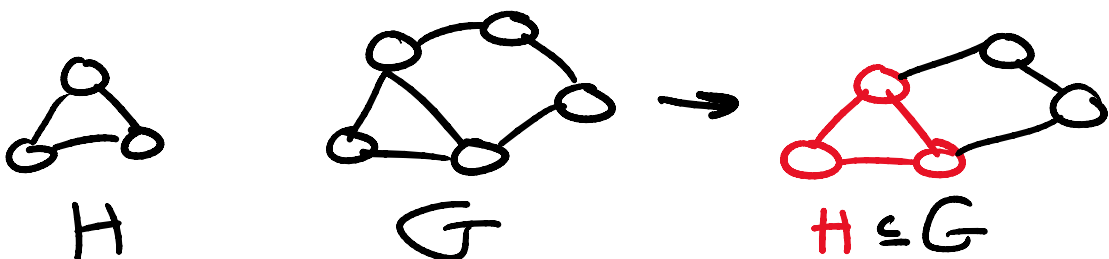


u, v are connected if there exists a path along edges from v to u



G is connected if there exists such a path for all u, v pairs

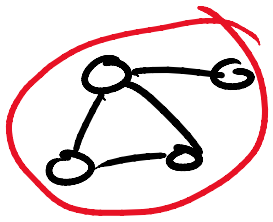
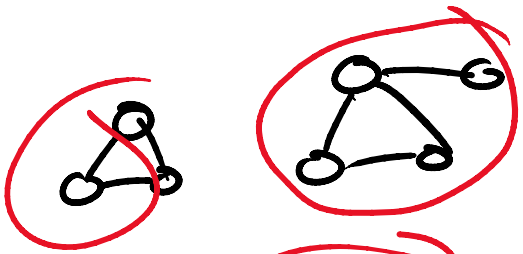
connectivity decomposition: find all maximal connected subgraphs within same G



H

G

$\bar{H} \subseteq \bar{G}$



H is a subgraph of G



3 connected maximal subgraphs

Easy way to solve using our framework:

Each vertex gets a unique label

do

all vertices assumes maximum label from its 1-hop neighborhood while some vertex updates a label

at the end \rightarrow each unique label corresponds to a unique connected component