

## Quick Review:

### Community detection (CD)

→ Finding dense regions in a graph or network

→ Defining "dense"

\* cut edge / internal edges

\* modularity

\* conductance

### Approaches:

→ optimize "density" directly

→ agglomerative, divide heuristics

→ greedy algos

### Evaluation

→ relative to optimization metrics

→ relative to "ground truth"

Today: new approach

↳ Spectral clustering

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Spectral graph theory (mining)

Basically: study and applications of eigenvalues and eigenvectors of a graph's adjacency matrix or some modified adjacency matrix

PageRank:

- spectral algo. on transition probability matrix

$$- M = (D^{-1}A)^T \rightarrow Mx = \lambda x \rightarrow \lambda = 1$$
$$Mx = x$$

Eigenvector centrality

- similar to PageRank

$$- Ax = \lambda x$$

Other applications:

Other applications.

- Clustering
  - Partitioning
  - Coarsening
  - Visualization
- 

Graph Laplacian

$L$  = graph Laplacian

$L = D - A \rightarrow$  adj. matrix

↳ diagonal  
degrees

Normalized Laplacian

$$L_N = I - D^{-1/2} A D^{1/2}$$

↳ fixes largest eigenvalue to be 1  
(like transition prob. matrix)

Overall: maps discrete  $\rightarrow$  continuous space

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# Spectral clustering

spectral decomposition

$$L = U \Lambda U^T$$

$\Lambda$  = diagonal matrix of eigenvalues

$U$  = associated eigenvectors

↳ since  $L$  is symmetric, all of  $\Lambda$  are real, positive, and non-increasing

Fiedler vector/value

→ first non-null eigenvalue/vector of our Laplacian

Fiedler value is related to the connectivity of  $G$

Larger = more connected

smaller = less connected

$\mathcal{P}_i$ -clustering Algorithm

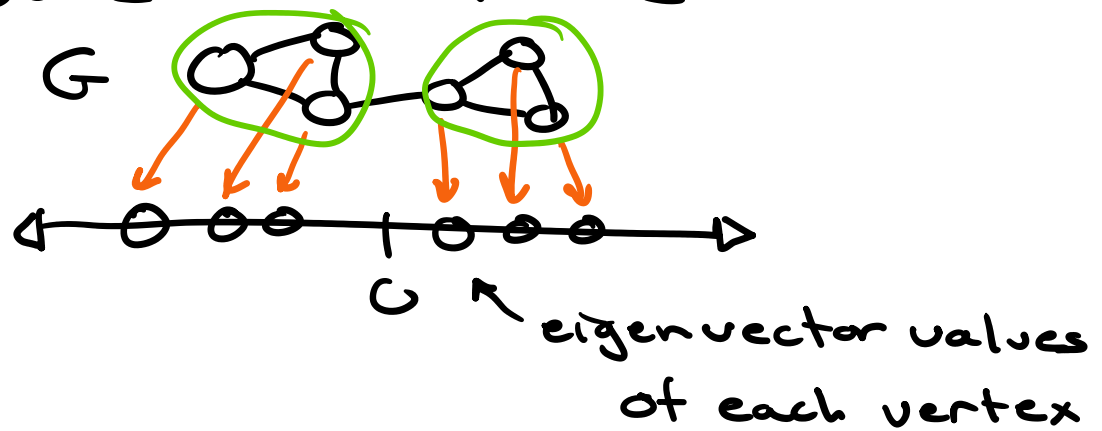
→ compute the Fiedler vector

→ we get values for each vertex

- we get values for each vertex
- if  $\text{value}[v] \geq 0 \rightarrow \text{cluster A}$
- $\text{value}[v] < 0 \rightarrow \text{cluster B}$

Intuitively: we mapping each vertex from discrete space to a 1D continuous space

Think of putting each vertex along some number line



What about some arbitrary  $k$  number of clusters?

- we can consider clustering each vertex's eigenvector values for some  $k$  in some  $j$ -dimensional space
- i.e., via  $k$ -means

↳ i.e., via  $k$ -means

→ OR: recursive bisection

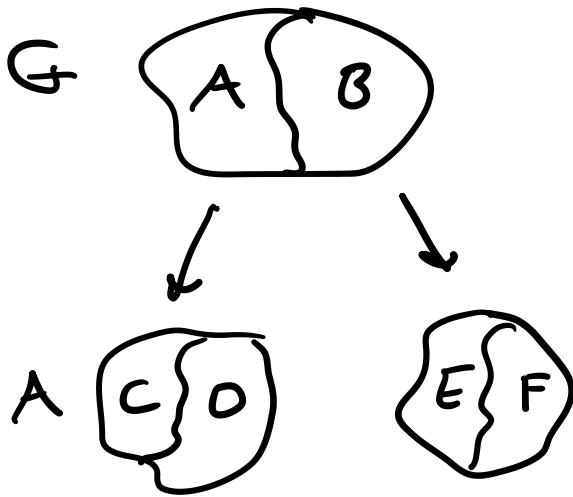
↓  
we get  $G = A + B$

via Fiedler cut

↳ can then get  $A = C + D$

$B = E + F$

using same approach  
on  $A, B$



Downside: we are  
constrained by  
power-of-2s

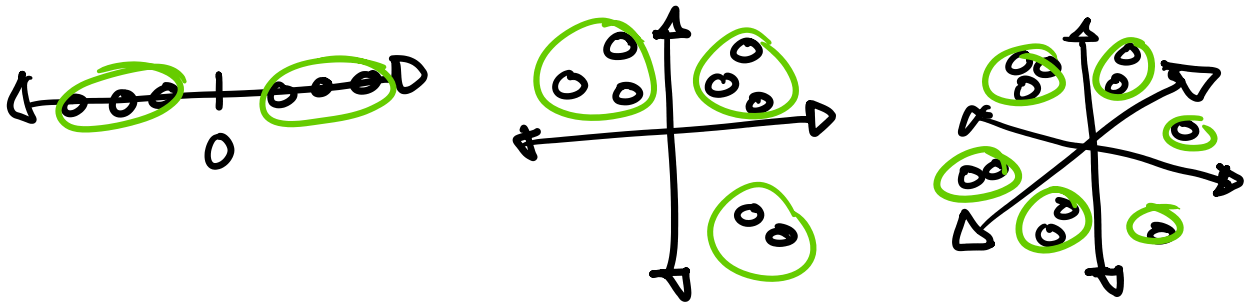
Using  $k$ -means to find  $k$ -clusters

→ Often:  $k$  eigenvectors of the  
largest  $k$  eigenvalues

→ Alternatively: we can use  
 $(\log_2 k) + 1$  with geometric  
bisection

↳ ... with 1D space

↳ we saw with 1D space  
for  $k=1$



Using  $j$ -eigen vectors to cluster  
into  $k$  clusters

- compute the  $j$  eigen vectors
- each vertex gets  $j$  eigen vector values
- this gives coordinates for each vertex in  $j$ -dimensional space
- we use those coordinates to do  $k$ -means clusters
- output of  $k$ -means is our cluster assignments

$k$ -means:

Unsupervised clustering algorithm  
to find  $k$  clusters in same

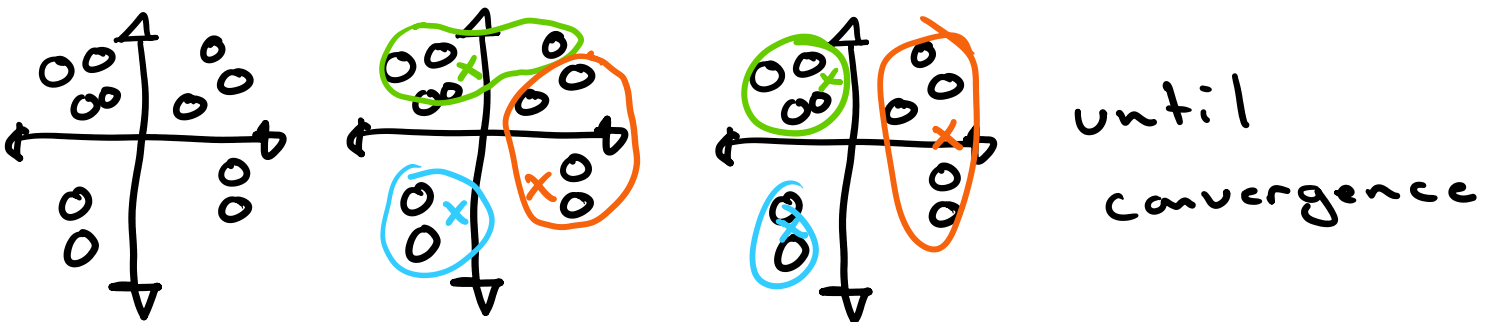
to find  $k$  clusters in some  $j$ -dimensional space

→ we initialize  $k$  points with  $j$  associated coordinates

→ each vertex joins the cluster defined by the smallest distance to one of those  $k$  points

→ each of those  $k$  points updates their coordinates to the mean coordinate values of all member vertices in the cluster

→ iterate until done



$k=3$

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Spectral partitioning



I.e., balanced graph partitioning

→ Given  $G$ , we want  $k$  "equal" sized parts

→ Usually formulated as an optimization problem

\* minimize edge cut (inter-part edges)  
\* constrain part sizes

such that  $|S_i| \approx |S_j| \forall i, j$

$$|S_i| \leq \frac{|V(G)|}{k} \epsilon \leftarrow \begin{array}{l} \text{imbalance} \\ \text{tolerance} \end{array}$$

Note: many possible variants to this optimization problem

$$\epsilon = 1.01$$

↳ 1% imbalance

Why: often used in scientific computations and parallel processing

Meshes == graphs, which are used in 99.999...% of scientific problems

Edge cut = amount of communication

Imbalance = workload imbalance

## Spectral partitioning

→ approach it like clustering, but  
constrain cluster sizes

→ Recall = "good" clusters have high  
density and therefore  
a low edge cut

↳ Intuition: we can optimize cut  
by finding good clusters

Q: How can we constrain cluster sizes?

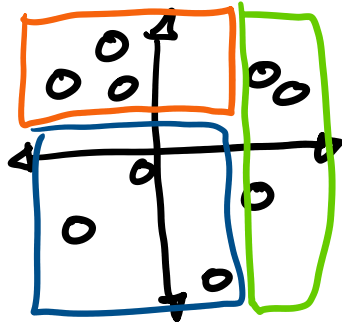
A: For bipartitioning, select the median  
value for the cluster assignments

↳ same for power-2 partitioning  
via recursive bisection

For  $k$ -way partitioning: we need to  
enforce balanced  $k$ -means

OR: a geometric-based approach

UK: a geometric-based approach



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## Visualization

Generally: drawing a graph in  
2D or 3D or some more  
complex multi-level space

## Approaches:

Spring-model  $\rightarrow$  model edges as  
springs and vertices as weights,  
and solve for some stable  
solution in 2D/3D space

Spectral  $\rightarrow$  basically what we've  
been talking about, using the  
eigenvector values to map vertices  
to Euclidean <sup>(sp.)</sup> space

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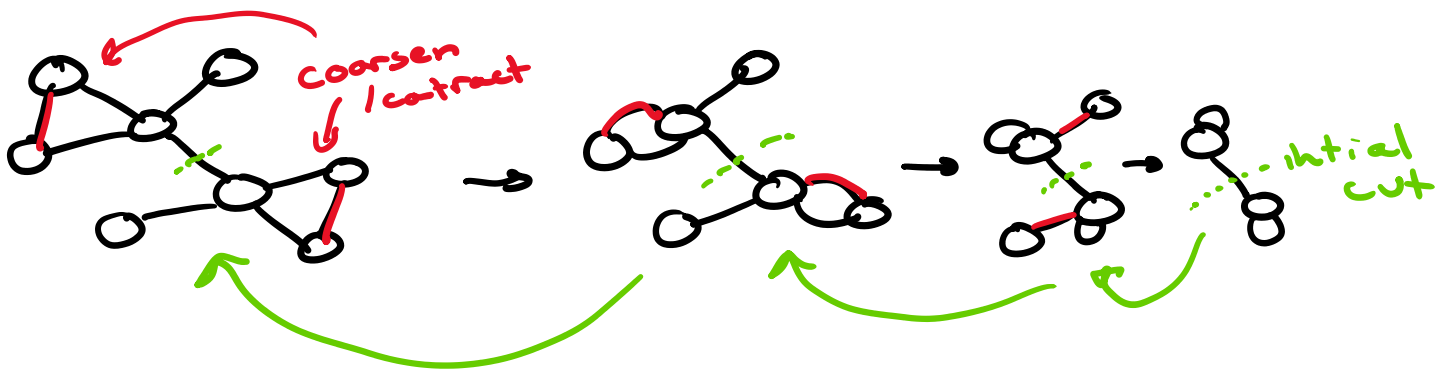
## Graph Coarsening

# Graph Coarsening

↳ Representing same graph as a reduced order model

Use: solving complex graph problems on a smaller graph, and then mapping that solution onto the larger original graph

(often done with partitioning)



Approches for Coarsening:

- Label propagation
- max-weight/cardinality matching
- Spectral coarsening: we select the vertices with the closest eigenvector values to contract

eigenvector values to contract