

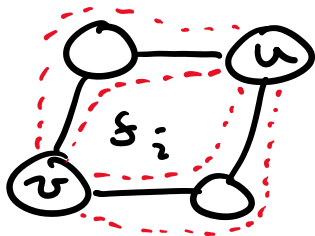
① Consider if there was a maximal planar graph which is not a triangulation

→ $\exists f_i$, a face with a length of 4 or greater

Consider non-adjacent $u, v \in V(f_i)$

→ we will always be able to add an edge (u, v) while remaining planar

→ simply draw (u, v) by following the boundary of f_i



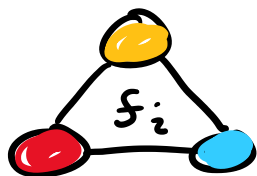
→ possible drawings of (u, v) , depending on if f_i is the outer face

Contradiction

⇒ we can always add such a

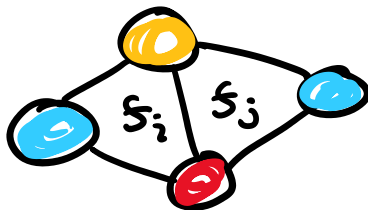
\Rightarrow as we can always add such a (u,v) to any $f_i: l(f_i) \geq 4$, a maximal planar G must be a triangulation \square

② Consider the vertex colorings of any face of G



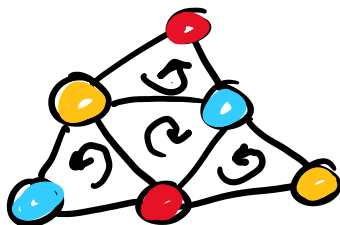
\rightarrow we note they must all have the same 3 colors

Consider two adjacent faces



\rightarrow we note there is a dependence via the shared edge

Consider an arbitrary number of adjacent faces



\rightarrow given some color order {e.g., $\text{yellow} \rightarrow \text{blue} \rightarrow \text{red}$ }, each face follows it clockwise or counter-clockwise

or counter-clockwise

→ No two adjacent faces have the same order, so any closed walk on the dual graph must be even, hence G^* is bipartite

Now consider $(G^*)^*$

→ we know there exists a way to draw $(G^*)^*$ s.t. $G \cong (G^*)^*$

→ we know that every face of G^* must be even

\Rightarrow this implies $G \cong (G^*)^*$ will have even degrees $\forall v \in V(G)$ and it is therefore Eulerian \square