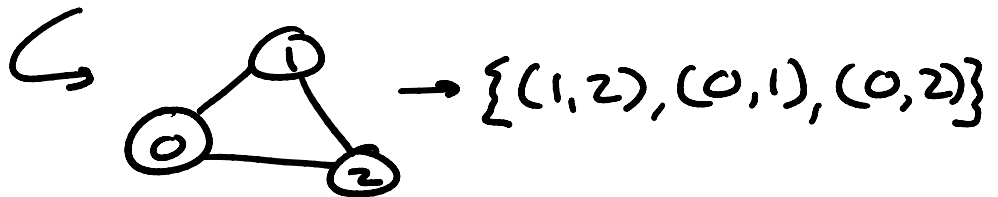


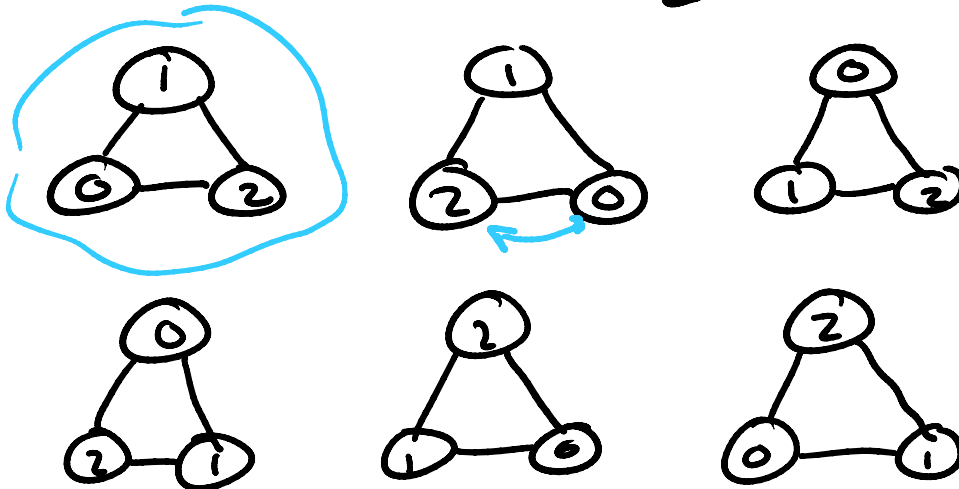
'Graph Theory is a terminological jungle, in which any newcomer may plant a tree.'  
 - John Barnes

# Automorphism

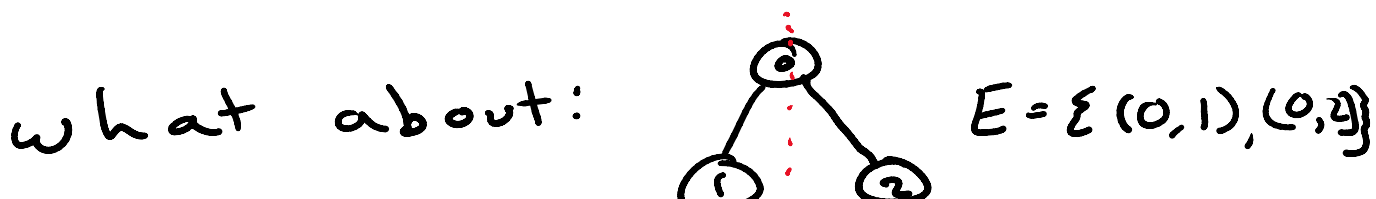
→ An isomorphic mapping of same  $G$  to itself  
 (s.t. the edge list is preserved)



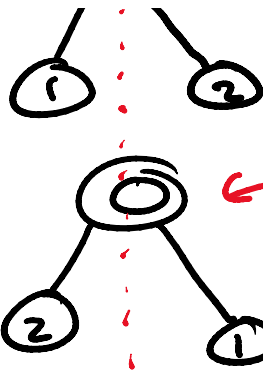
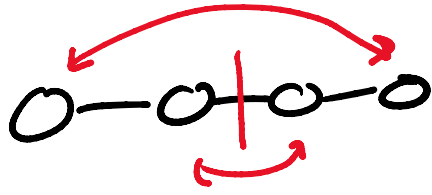
$K_n$ : we can map any vertex to any other vertex



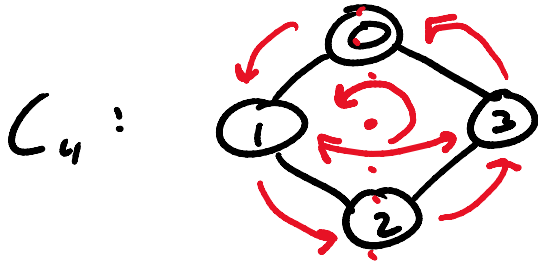
Note:  $n!$  total automorphisms



What about.



$$E = \{(0,1), (0,2)\}$$



$$E = \{(0,1), (1,2), (2,3), (0,3)\}$$

rotational symmetry too!

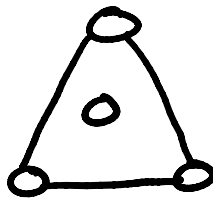
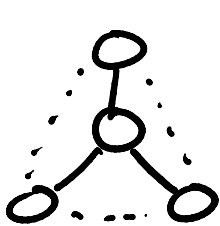
More definition

complement of  $G \rightarrow \bar{G}$   
(assume  $G$  is simple)

$$V(\bar{G}) = V(G)$$

$$E(\bar{G}) = \{ \forall u, v \in V(G) : (u, v) \notin E(G) \}$$

↑  
not in



$G$

$\bar{G}$

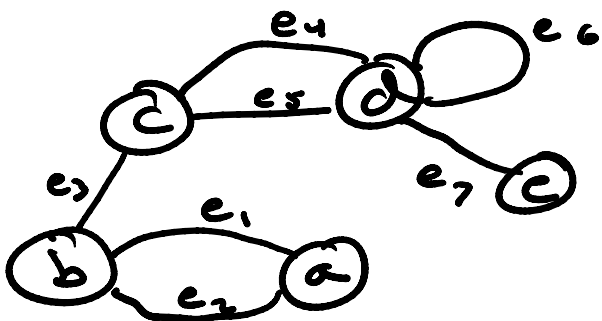
Time ⌚ to take ..

Time ⌚ to take a stroll

walk: a list of vertices and edges  
 s.t. each listing is adjacent/incident to the listing preceding and proceeding

trail: as above, except edges do not repeat

path: as above, except edges and vertices do not repeat



$W: \{a e_1 b e_3 c e_2 b e_3 c e_4 d e_5 d e_7 e\}$

$w: \{e_1 e_3 e_3 e_3 e_4 e_6 e_7\}$

$\Gamma: \{e_1 e_3 e_4 e_6 e_7\}$

$P: \{e_1 e_2 e_4 e_7\}$

↑

a, e-path

↑  
end

→



A walk/trail/path that starts and ends at the same vertex is closed otherwise, they are open

Note: a closed path is really just a cycle

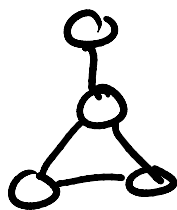
Length: number of edges traversed

Hop: traversal of a single edge

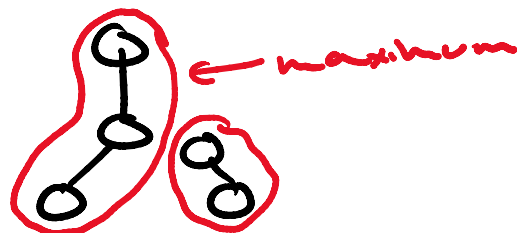
Let's get connected

Recall:  $G$  is connected if

$\forall u, v \in V(G): \exists u, v\text{-path}$



$G$  is connected



$G'$  is disconnected

connected component: a maximal  
connected subgraph of  $G$

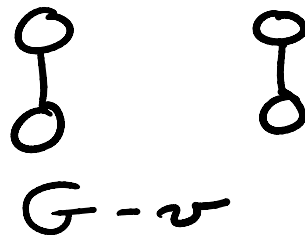
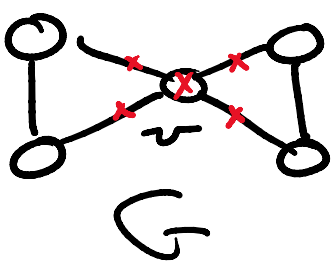
maximal: can not be made larger

maximum: the largest possible

Note: same for minimal/minimum  
but smaller/smallest

cut vertex: some  $v \in V(G)$  s.t.

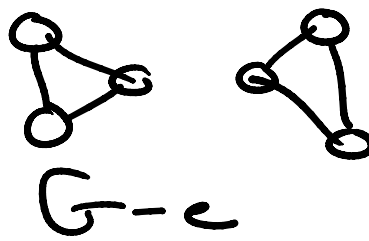
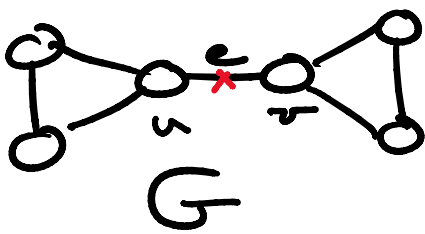
$G - v$  has more connected  
vertex deletion components than  
 $G$   
→ remove  $v$  and  
all incident edges on  $v$



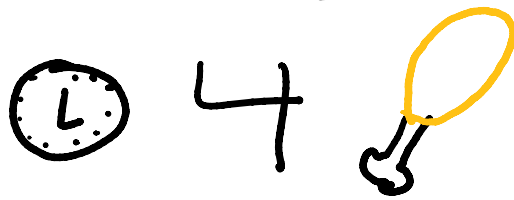
cut edge: some  $e \in E(G)$  s.t.

$G - e$  has more components  
than  $G$   
→ deletion

edge deletion than  $G$   
 $\rightarrow$  just remove  $e$



Time for the meat  
of graph theory



aka  
proovin' stuff

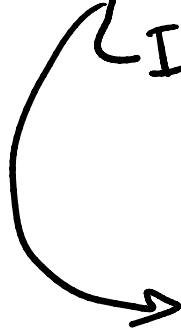
weak induction  $\Downarrow$

Prove:  $2^1 + 2^2 + 2^3 \dots + 2^n = 2^{n+1} - 2$

Basis:  $P(1) \Rightarrow 2^1 = 2^{1+1} - 2 = 2 \checkmark$

Inductive step:  $P(n=k+1)$

Inductive Hypothesis: we assume what  
we're trying to prove holds  
for some  $P(k)$



$c_1 \dots n_1 \dots$



↑  
basis

↑  
assume holds  
via I.H.  
for \*all\*  $1 \leq k < n$

↓  
show it  
holds

Example Proof:

length is odd  
↓

Show every closed odd walk  
contains an odd cycle

↙ Induction on the length of the walk

Basis:  $P(1) = \emptyset \checkmark$

Inductive step:  $P(n > k \geq 1)$

↳ Assume we have  $P(n)$ , which is  
of length  $n$ , odd, closed

Consider the cases

Generally: what possible configurations  
we can use to simplify our  
proof

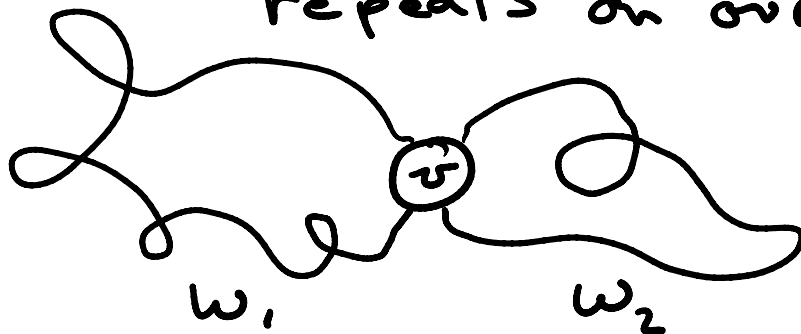
Case 1: no vertices repeat  
on our walk

$\Rightarrow$  trivial case, our odd walk



$\Rightarrow$  trivial case, our odd walk  
 $\Rightarrow$  just a cycle  $\checkmark$

Case 2: at least one vertex  $v$   
repeats on our walk  $w$



$$|w| = |w_1| + |w_2|$$

$\% \text{ integer parity } \%$   
 $\% \qquad \qquad \qquad \%$

this means  $\text{odd} + \text{odd} = \text{even}$   
 $\text{odd} + \text{even} = \text{odd}$   
 $\text{even} + \text{even} = \text{even}$

this implies w.l.o.g.  $|w_1| = \text{odd}$   
without loss  
of generality

Note:  $|w_1| < |w|$

$\hookrightarrow$  we define  $P(k) = w_1$   
and use our I.H. on  $P(k)$

$\Rightarrow \exists$  same odd cycle on  $P(k)$

Bring it back to  $P(n)$

$\hookrightarrow$  show our property still holds on  $P(n)$

In this case:

adding  $w_2$  back to  $w_1$ ,

to create  $w = P(n)$  will not delete that odd cycle on  $w$   $\square$