

Frontiers of Network Science Fall 2021

Class 4: Graph Theory II (Chapter 2 in Textbook) Boleslaw Szymanski

Section 2.2

Reference Networks

NETWORK	NODES	LINKS	DIRECTED UNDIRECTED	N	L
Internet	Routers	Internet connections	Undirected	192,244	609,066
WWW	Webpages	Links	Directed	325,729	1,497,134
Power Grid	Power plants, transformers	Cables	Undirected	4,941	6,594
Mobile Phone Calls	Subscribers	Calls	Directed	36,595	91,826
Email	Email addresses	Emails	Directed	57,194	103,731
Science Collaboration	Scientists	Co-authorship	Undirected	23,133	93,439
Actor Network	Actors	Co-acting	Undirected	702,388	29,397,908
Citation Network	Paper	Citations	Directed	449,673	4,689,479
E. Coli Metabolism	Metabolites	Chemical reactions	Directed	1,039	5,802
Protein Interactions	Proteins	Binding interactions	Undirected	2,018	2,930

Degree, Average Degree and Degree Distribution

BRIEF STATISTICS REVIEW

Four key quantities characterize a sample of N values x_1, \dots, x_N :

Average (mean):

$$\langle x \rangle = \frac{x_1 + x_2 + \dots + x_N}{N} = \frac{1}{N} \sum_{i=1}^N x_i$$

The n^{th} moment:

$$\langle x^n \rangle = \frac{x_1^n + x_2^n + \dots + x_N^n}{N} = \frac{1}{N} \sum_{i=1}^N x_i^n$$

Standard deviation:

$$\sigma_x = \sqrt{\frac{1}{N} \sum_{i=1}^N (x_i - \langle x \rangle)^2}$$

Distribution of x :

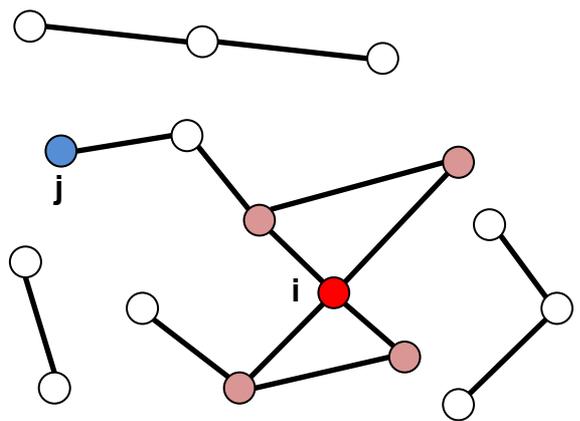
$$p_x = \frac{1}{N} \sum_i \delta_{x, x_i}$$

where p_x follows

$$\sum_i p_x = 1 \quad \left(\int p_x dx = 1 \right)$$

AVERAGE DEGREE

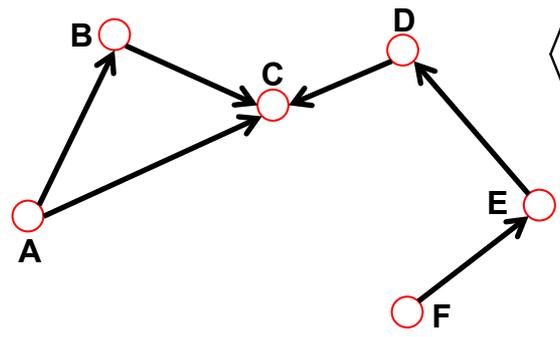
Undirected



$$\langle k \rangle \equiv \frac{1}{N} \sum_{i=1}^N k_i \quad \langle k \rangle \equiv \frac{2L}{N}$$

N – the number of nodes in the graph

Directed



$$\langle k^{in} \rangle \equiv \frac{1}{N} \sum_{i=1}^N k_i^{in}, \quad \langle k^{out} \rangle \equiv \frac{1}{N} \sum_{i=1}^N k_i^{out}, \quad \langle k^{in} \rangle = \langle k^{out} \rangle$$

$$\langle k \rangle \equiv \frac{L}{N}$$

DEGREE DISTRIBUTION

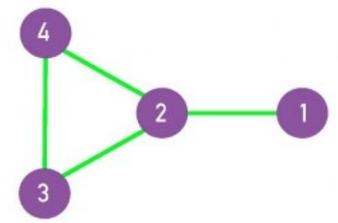
Degree distribution

$P(k)$: probability that a randomly chosen node has degree k

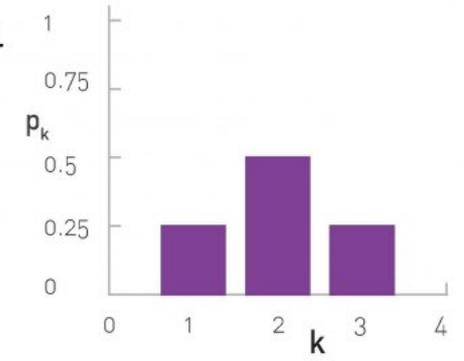
$N_k = \#$ nodes with degree k

$P(k) = N_k / N$  plot

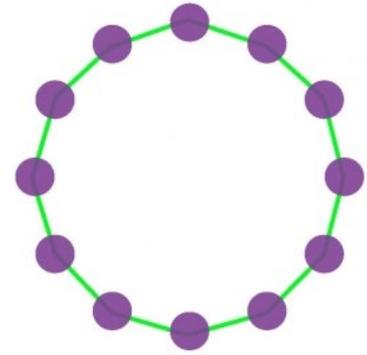
a.



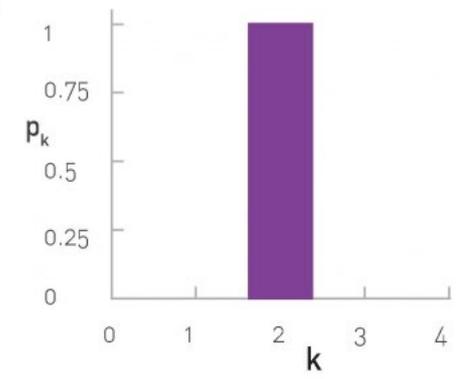
b.



c.



d.



DEGREE DISTRIBUTION

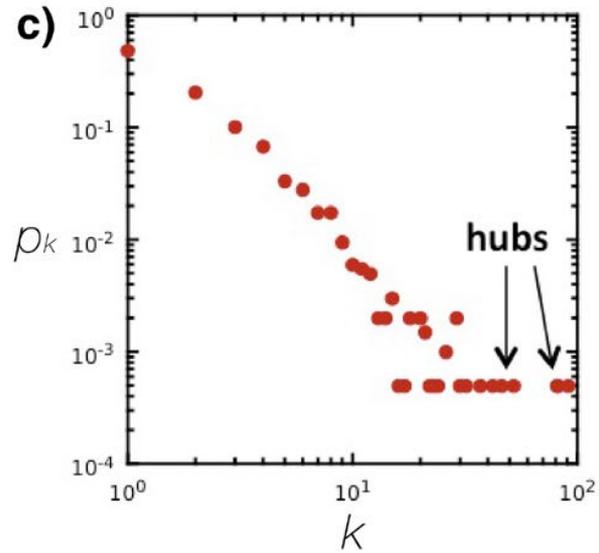
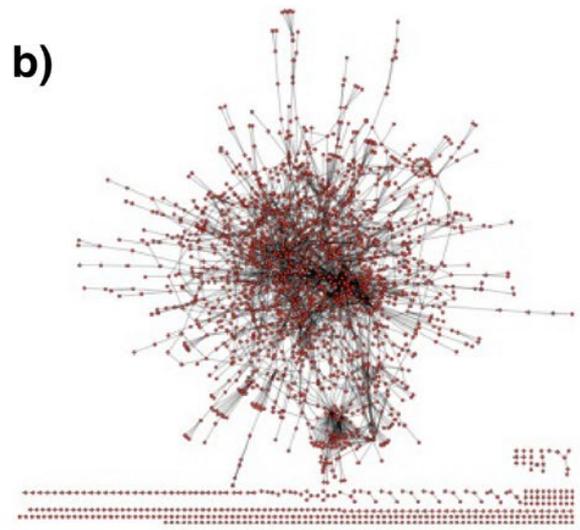
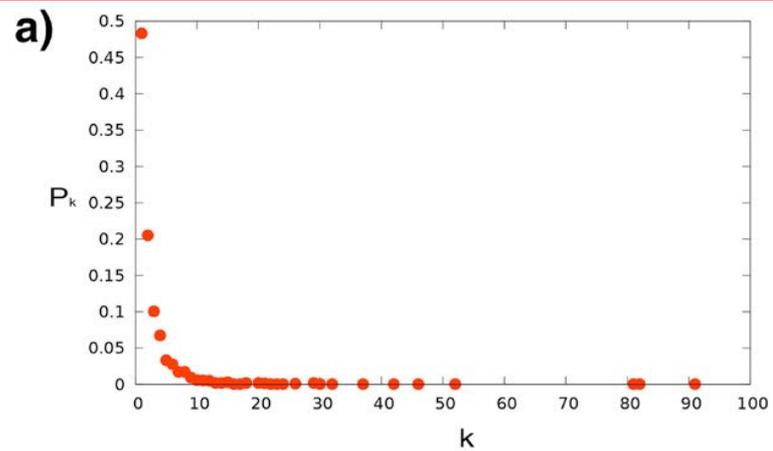


Image 2.4b

DEGREE DISTRIBUTION

Discrete Representation: p_k is the probability that a node has degree k .

Continuum Description: $p(k)$ is the pdf of the degrees, where

$$\int_{k_1}^{k_2} p(k) dk$$

represents the probability that a node's degree is between k_1 and k_2 .

Normalization condition:

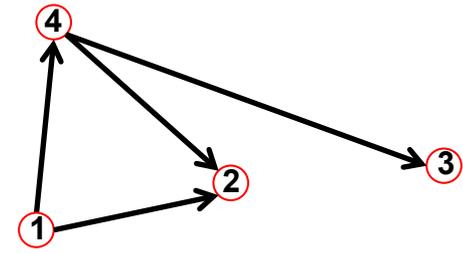
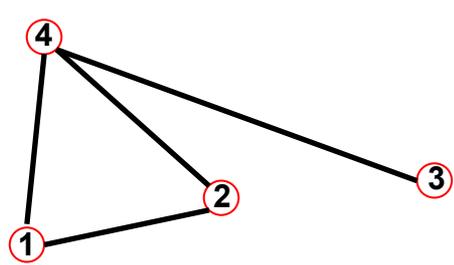
$$\sum_0^{\infty} p_k = 1$$

$$\int_{K_{\min}}^{\infty} p(k) dk = 1$$

where K_{\min} is the minimal degree in the network.

Adjacency matrix

ADJACENCY MATRIX



$A_{ij}=1$ if there is a link between node i and j

$A_{ij}=0$ if nodes i and j are not connected to each other.

$$A_{ij} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$A_{ij} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

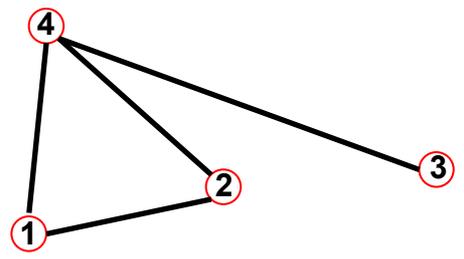
Note that for a directed graph (right) the matrix is not symmetric.

$A_{ij} = 1$ if there is a link pointing from node j and i

$A_{ij} = 0$ if there is no link pointing from j to i .

ADJACENCY MATRIX AND NODE DEGREES

Undirected



$$A_{ij} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

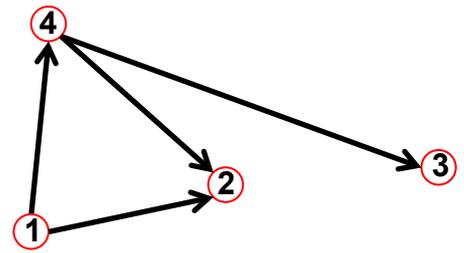
$$A_{ij} = A_{ji}$$
$$A_{ii} = 0$$

$$k_i = \sum_{j=1}^N A_{ij}$$

$$k_j = \sum_{i=1}^N A_{ij}$$

$$L = \frac{1}{2} \sum_{i=1}^N k_i = \frac{1}{2} \sum_{ij} A_{ij}$$

Directed



$$A_{ij} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

$$A_{ij} \neq A_{ji}$$
$$A_{ii} = 0$$

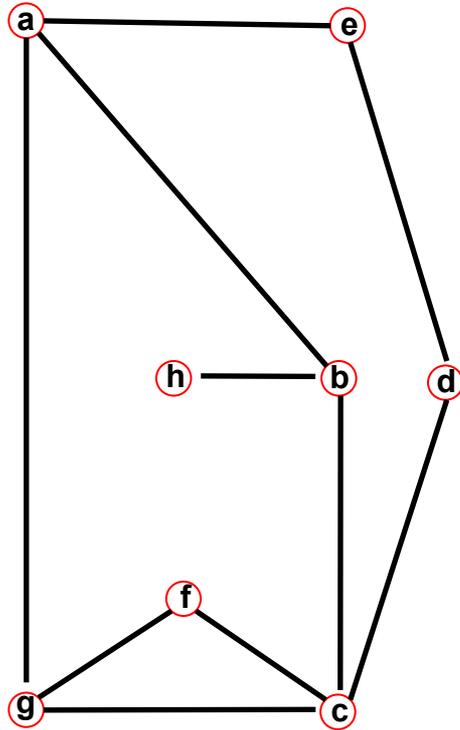
$$k_i^{in} = \sum_{j=1}^N A_{ij}$$

$$k_j^{out} = \sum_{i=1}^N A_{ij}$$

$$L = \sum_{i=1}^N k_i^{in} = \sum_{j=1}^N k_j^{out} = \sum_{i,j} A_{ij}$$



ADJACENCY MATRIX

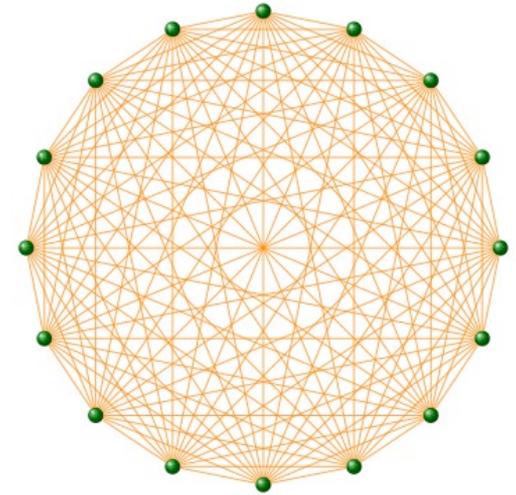


	a	b	c	d	e	f	g	h
a	0	1	0	0	1	0	1	0
b	1	0	1	0	0	0	0	1
c	0	1	0	1	0	1	1	0
d	0	0	1	0	1	0	0	0
e	1	0	0	1	0	0	0	0
f	0	0	1	0	0	0	1	0
g	1	0	1	0	0	0	0	0
h	0	1	0	0	0	0	0	0

Real networks are sparse

COMPLETE GRAPH

The maximum number of links a network of N nodes can have is: $L_{\max} = \binom{N}{2} = \frac{N(N-1)}{2}$



A graph with degree $L=L_{\max}$ is called a **complete graph**, and its average degree is $\langle k \rangle = N-1$

Most networks observed in real systems are sparse:

$$L \ll L_{\max}$$

or

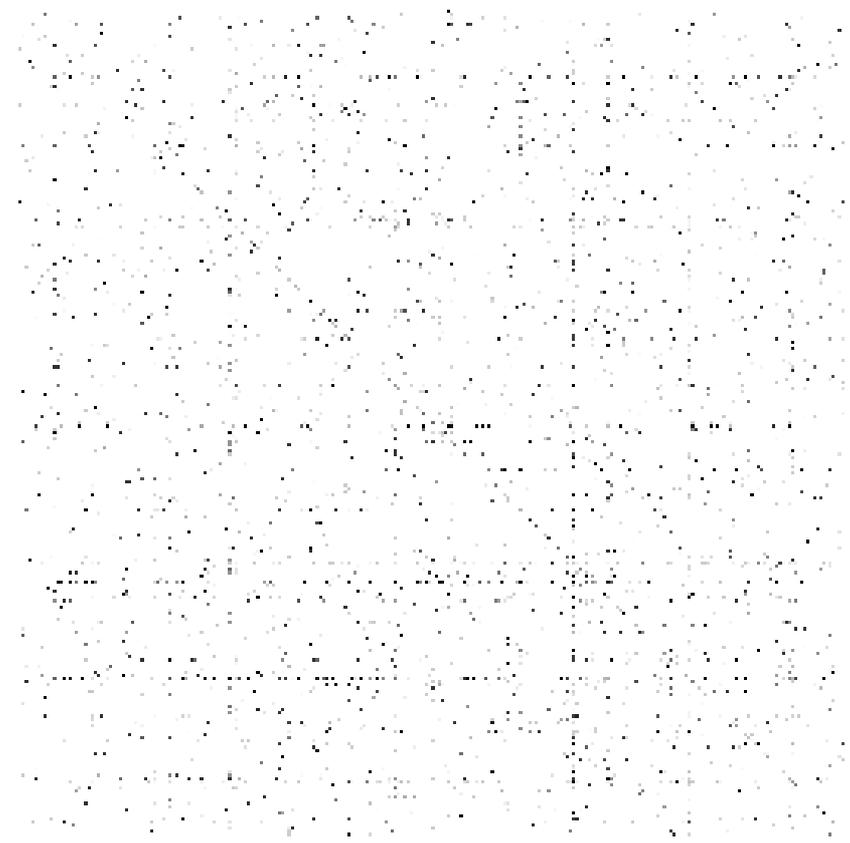
$$\langle k \rangle \ll N-1.$$

WWW (ND Sample):	$N=325,729;$	$L=1.4 \cdot 10^6$	$L_{\max}=10^{12}$	$\langle k \rangle=4.51$
Protein (<i>S. Cerevisiae</i>):	$N= 1,870;$	$L=4,470$	$L_{\max}=10^7$	$\langle k \rangle=2.39$
Coauthorship (Math):	$N= 70,975;$	$L=2 \cdot 10^5$	$L_{\max}=3 \cdot 10^{10}$	$\langle k \rangle=3.9$
Movie Actors:	$N=212,250;$	$L=6 \cdot 10^6$	$L_{\max}=1.8 \cdot 10^{13}$	$\langle k \rangle=28.78$

(Source: Albert, Barabasi, RMP2002)



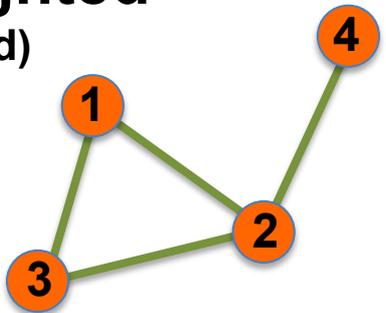
ADJACENCY MATRICES ARE SPARSE



WEIGHTED AND UNWEIGHTED NETWORKS

$$A_{ij} = w_{ij}$$

Unweighted (undirected)



$$A_{ij} = \begin{pmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}$$

$$A_{ii} = 0$$

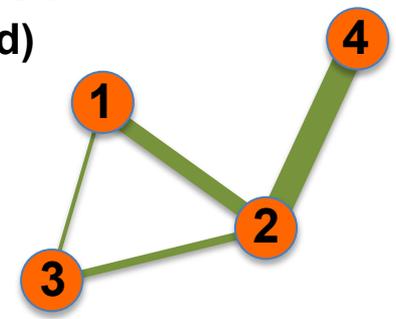
$$A_{ij} = A_{ji}$$

$$L = \frac{1}{2} \sum_{i,j=1}^N A_{ij}$$

$$\langle k \rangle = \frac{2L}{N}$$

protein-protein interactions, www

Weighted (undirected)



$$A_{ij} = \begin{pmatrix} 0 & 2 & 0.5 & 0 \\ 2 & 0 & 1 & 4 \\ 0.5 & 1 & 0 & 0 \\ 0 & 4 & 0 & 0 \end{pmatrix}$$

$$A_{ii} = 0$$

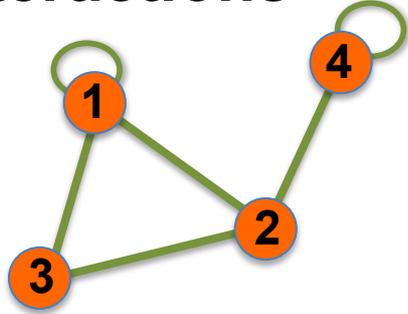
$$A_{ij} = A_{ji}$$

$$L = \frac{1}{2} \sum_{i,j=1}^N \text{nonzero}(A_{ij})$$

$$\langle k \rangle = \frac{2L}{N}$$

Call Graph, metabolic networks

Self-interactions



$$A_{ij} = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix}$$

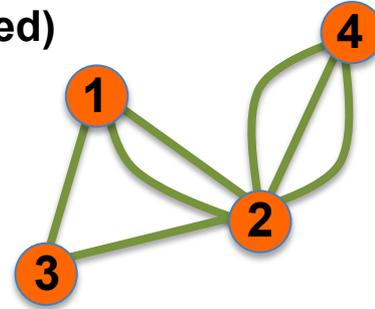
$$A_{ii} \neq 0$$

$$A_{ij} = A_{ji}$$

$$L = \frac{1}{2} \sum_{i,j=1, i \neq j}^N A_{ij} + \sum_{i=1}^N A_{ii} \quad ?$$

Protein interaction network, www

Multigraph (undirected)



$$A_{ij} = \begin{pmatrix} 0 & 2 & 1 & 0 \\ 2 & 0 & 1 & 3 \\ 1 & 1 & 0 & 0 \\ 0 & 3 & 0 & 0 \end{pmatrix}$$

$$A_{ii} = 0$$

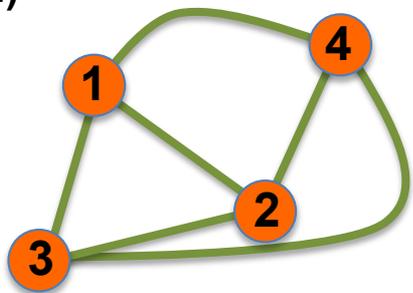
$$A_{ij} = A_{ji}$$

$$L = \frac{1}{2} \sum_{i,j=1}^N \text{nonzero}(A_{ij}) \quad \langle k \rangle = \frac{2L}{N}$$

Social networks, collaboration networks

Complete Graph

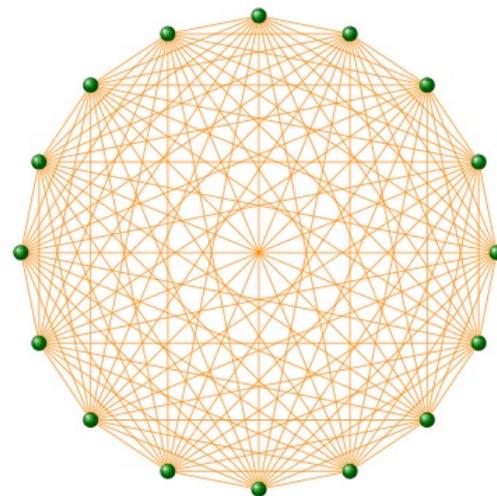
(undirected)



$$A_{ij} = \begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$A_{ii} = 0 \qquad A_{i \neq j} = 1$$

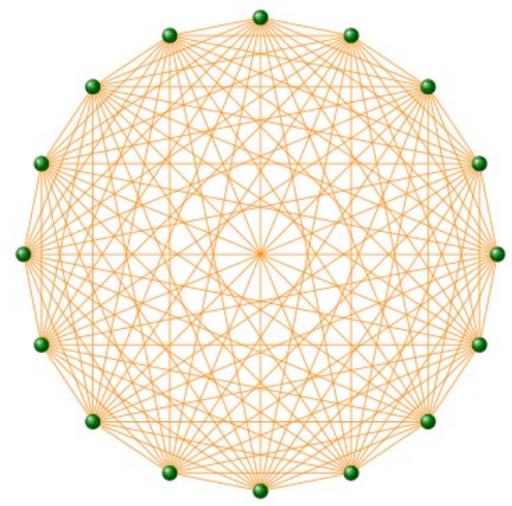
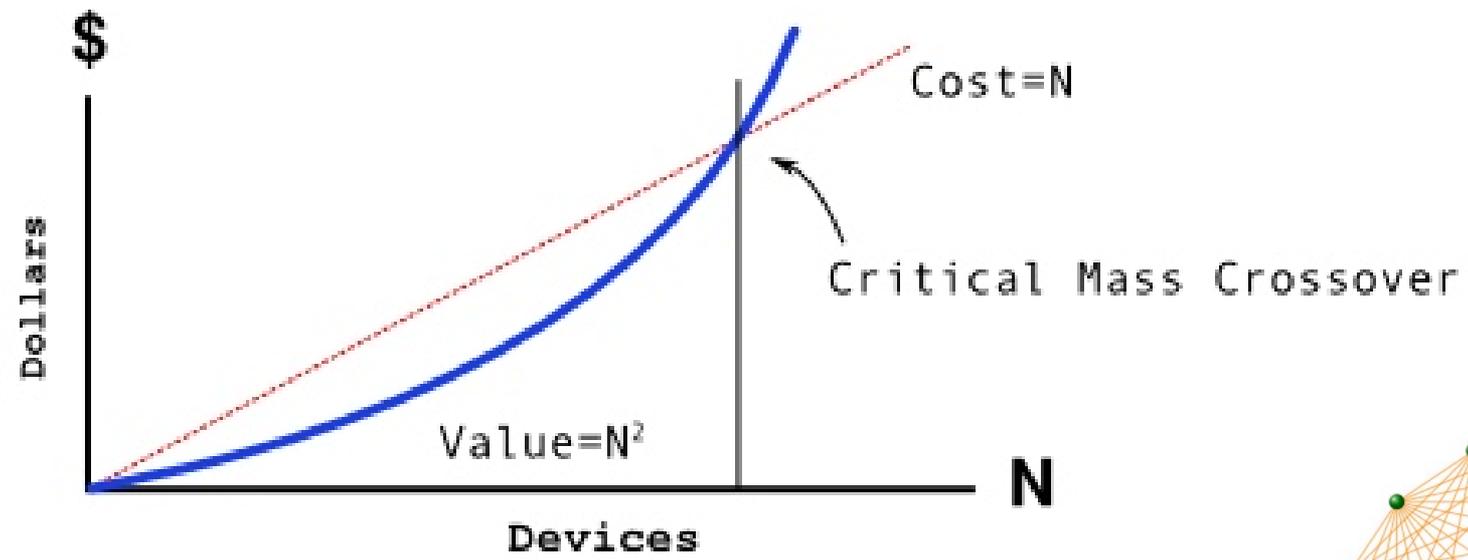
$$L = L_{\max} = \frac{N(N-1)}{2} \qquad \langle k \rangle = N-1$$



Actor network, protein-protein interactions



METCALFE'S LAW



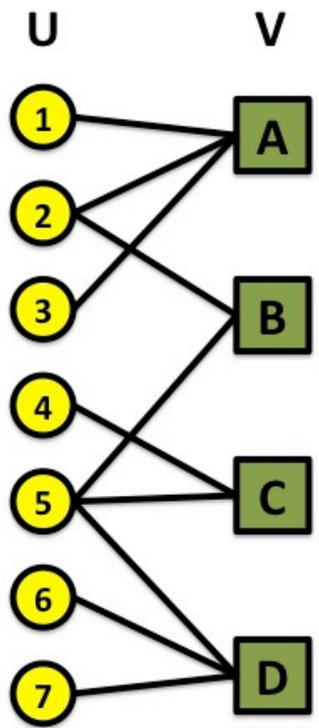
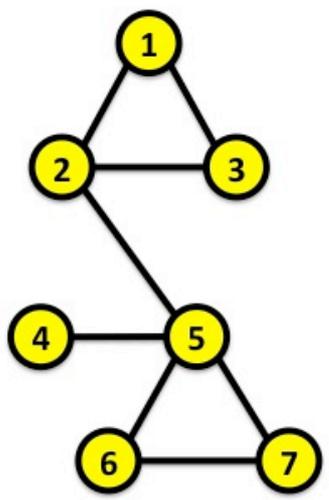
The maximum number of links a network of N nodes can have is:
$$L_{\max} = \binom{N}{2} = \frac{N(N-1)}{2}$$

BIPARTITE NETWORKS

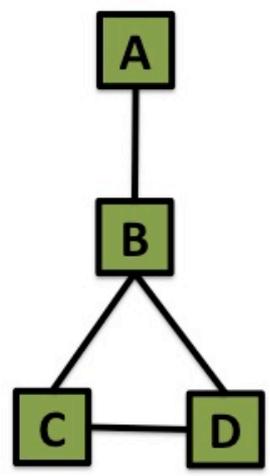
BIPARTITE GRAPHS

bipartite graph (or **bigraph**) is a [graph](#) whose nodes can be divided into two [disjoint sets](#) U and V such that every link connects a node in U to one in V ; that is, U and V are [independent sets](#).

Projection U



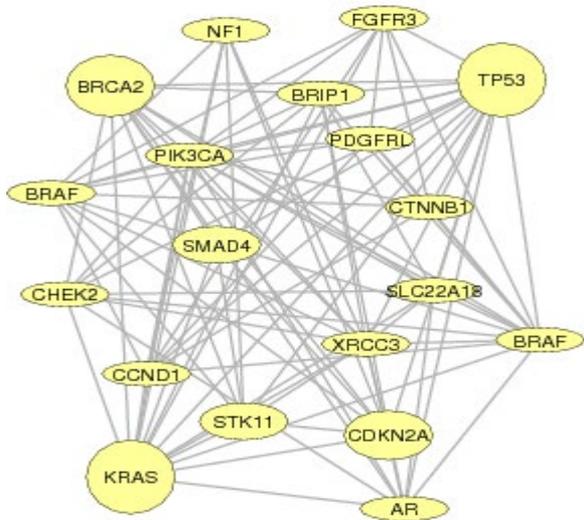
Projection V



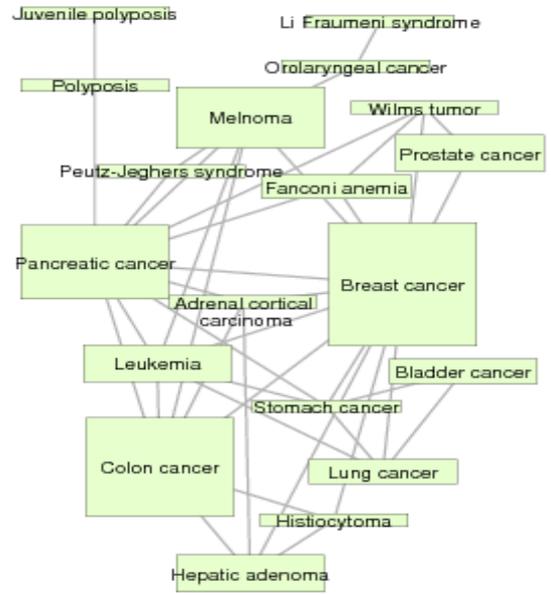
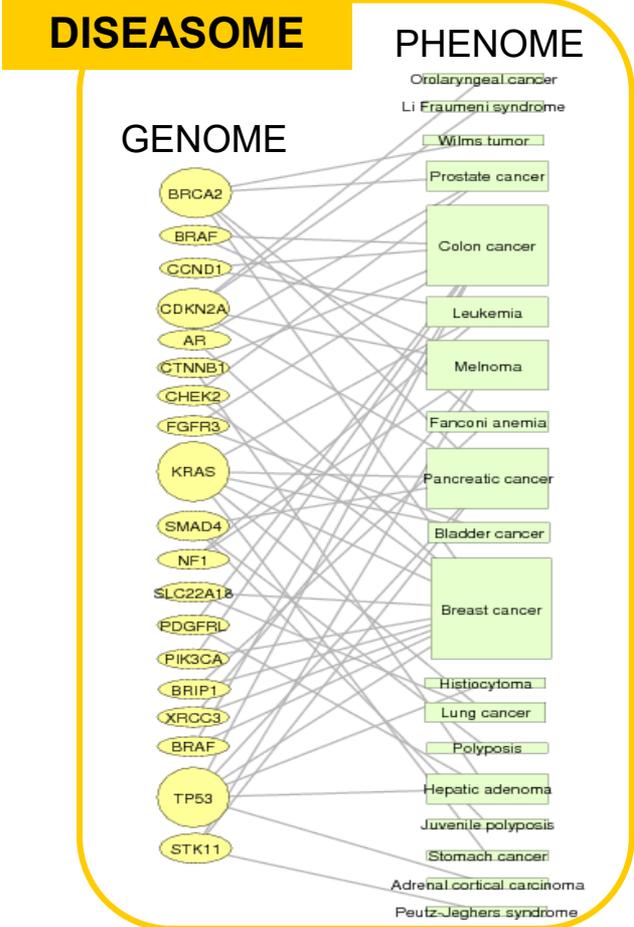
Examples:

- Hollywood actor network
- Collaboration networks
- Disease network (diseasome)

GENE NETWORK – DISEASE NETWORK



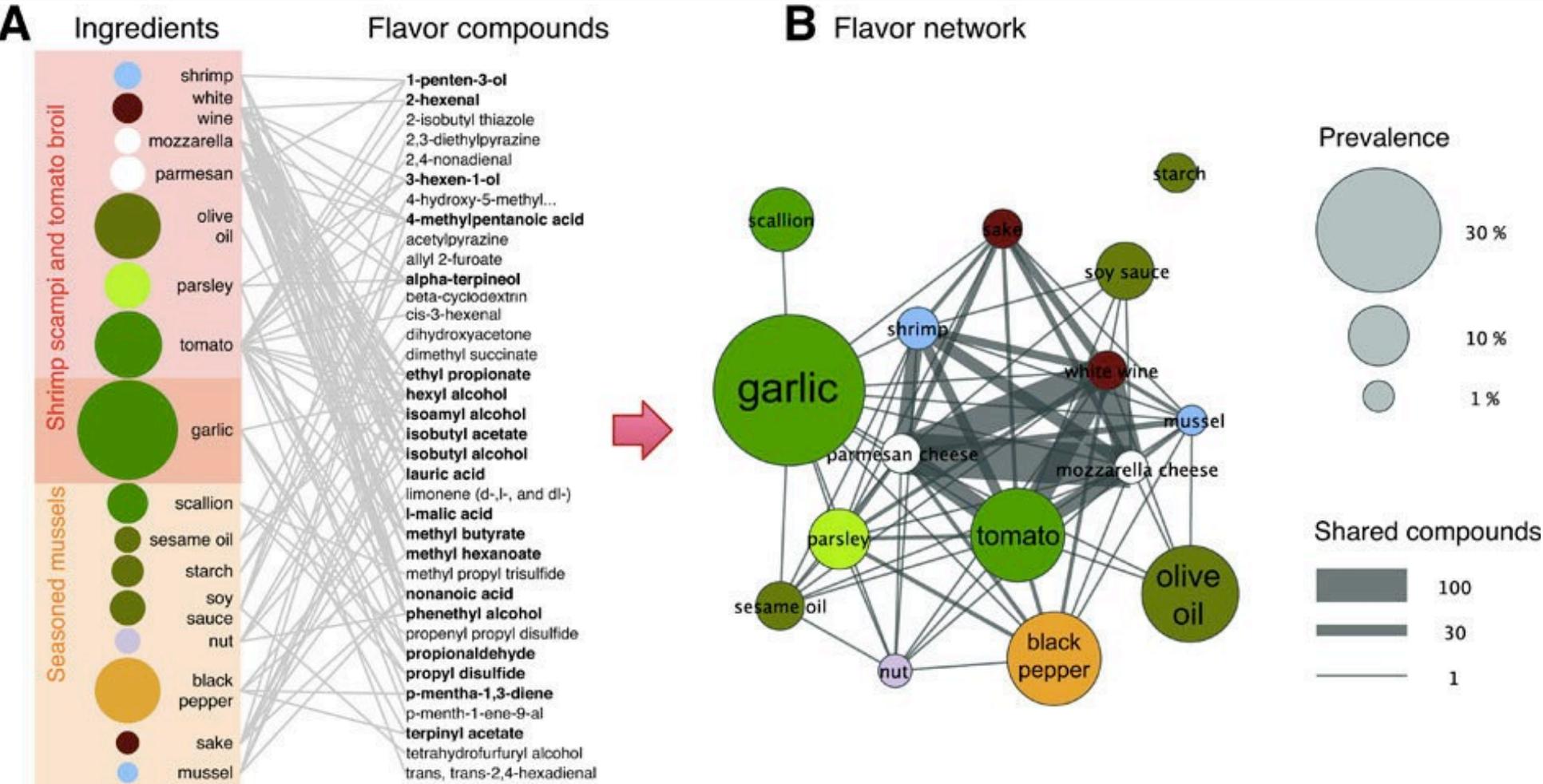
Gene network



Disease network

Goh, Cusick, Valle, Childs, Vidal & Barabási, PNAS (2007)

Ingredient-Flavor Bipartite Network



Y.-Y. Ahn, S. E. Ahnert, J. P. Bagrow, A.-L. Barabási ^[1] Flavor network and the principles of food pairing, *Scientific Reports* 196, (2011).

Categories

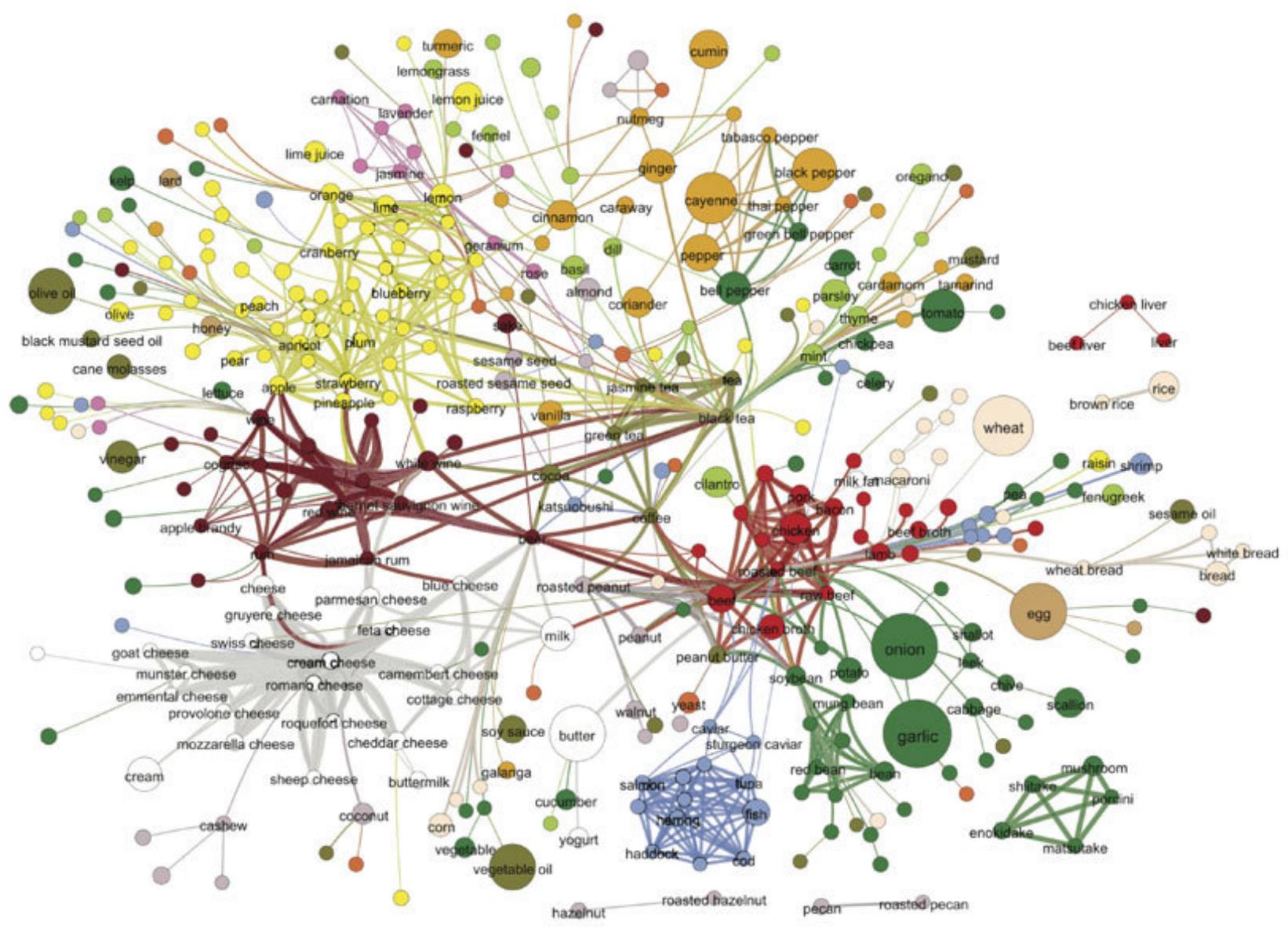
- fruits
- dairy
- spices
- alcoholic beverages
- nuts and seeds
- seafoods
- meats
- herbs
- plant derivatives
- vegetables
- flowers
- animal products
- plants
- cereal

Prevalence

- 50 %
- 30 %
- 10 %
- 1 %

Shared compounds

- 150
- 50
- 10



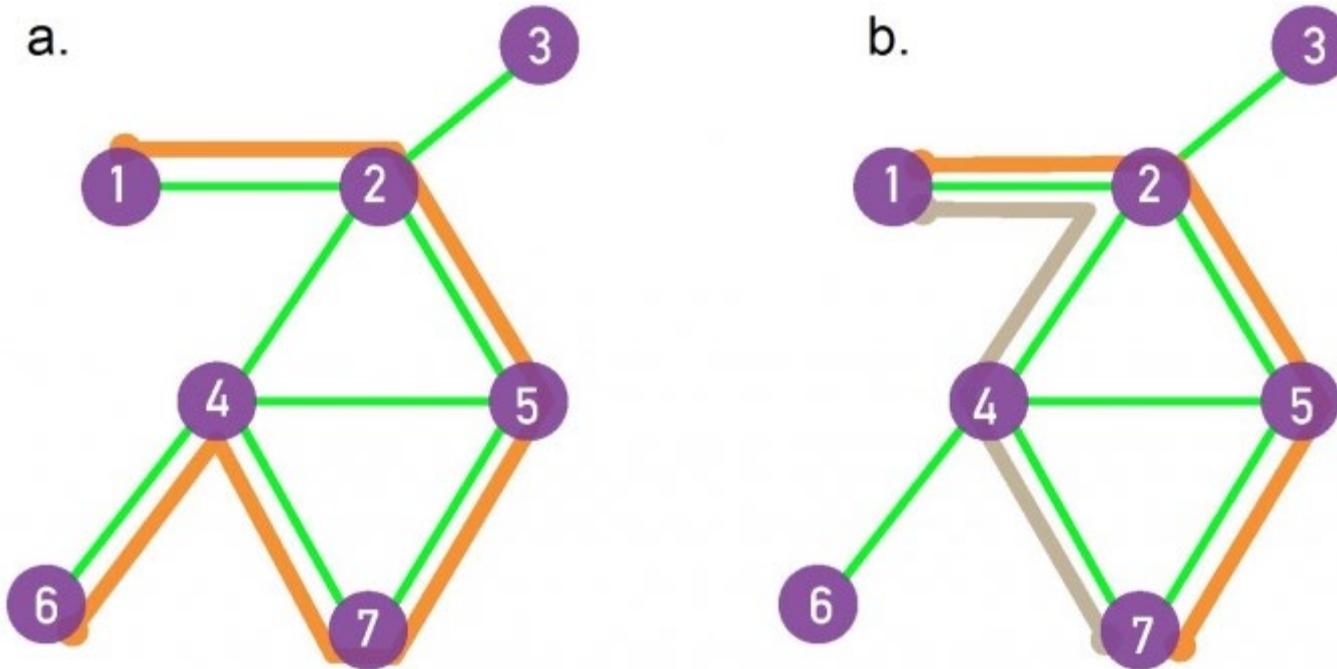
PATHOLOGY

PATHS

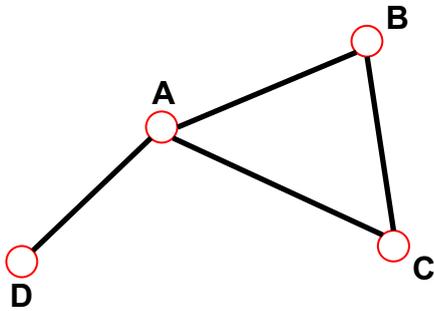
A *path* is a sequence of nodes in which each node is adjacent to the next one

P_{i_0, i_n} of length n between nodes i_0 and i_n is an ordered collection of $n+1$ nodes and n links

$$P_n = \{i_0, i_1, i_2, \dots, i_n\} \quad P_n = \{(i_0, i_1), (i_1, i_2), (i_2, i_3), \dots, (i_{n-1}, i_n)\}$$

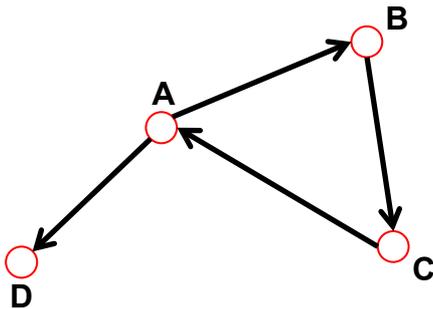


- In a directed network, the path can follow only the direction of an arrow.



The *distance (shortest path, geodesic path)* between two nodes is defined as the number of edges along the shortest path connecting them.

*If the two nodes are disconnected, the distance is infinity.



In **directed graphs** each path needs to follow the direction of the arrows.

Thus in a digraph the distance from node A to B (on an AB path) is generally different from the distance from node B to A (on a BCA path).

N_{ij} , number of paths between any two nodes i and j :

Length $n=1$: If there is a link between i and j , then $A_{ij}=1$ and $A_{ij}=0$ otherwise.

Length $n=2$: If there is a path of length two between i and j , then $A_{ik}A_{kj}=1$, and $A_{ik}A_{kj}=0$ otherwise.

The number of paths of length 2:

$$N_{ij}^{(2)} = \sum_{k=1}^N A_{ik}A_{kj} = [A^2]_{ij}$$

Length n : In general, if there is a path of length n between i and j , then $A_{i_1k_1} \dots A_{i_nj_n}=1$ and $A_{i_1k_1} \dots A_{i_nj_n}=0$ otherwise.

The number of paths of length n between i and j is*

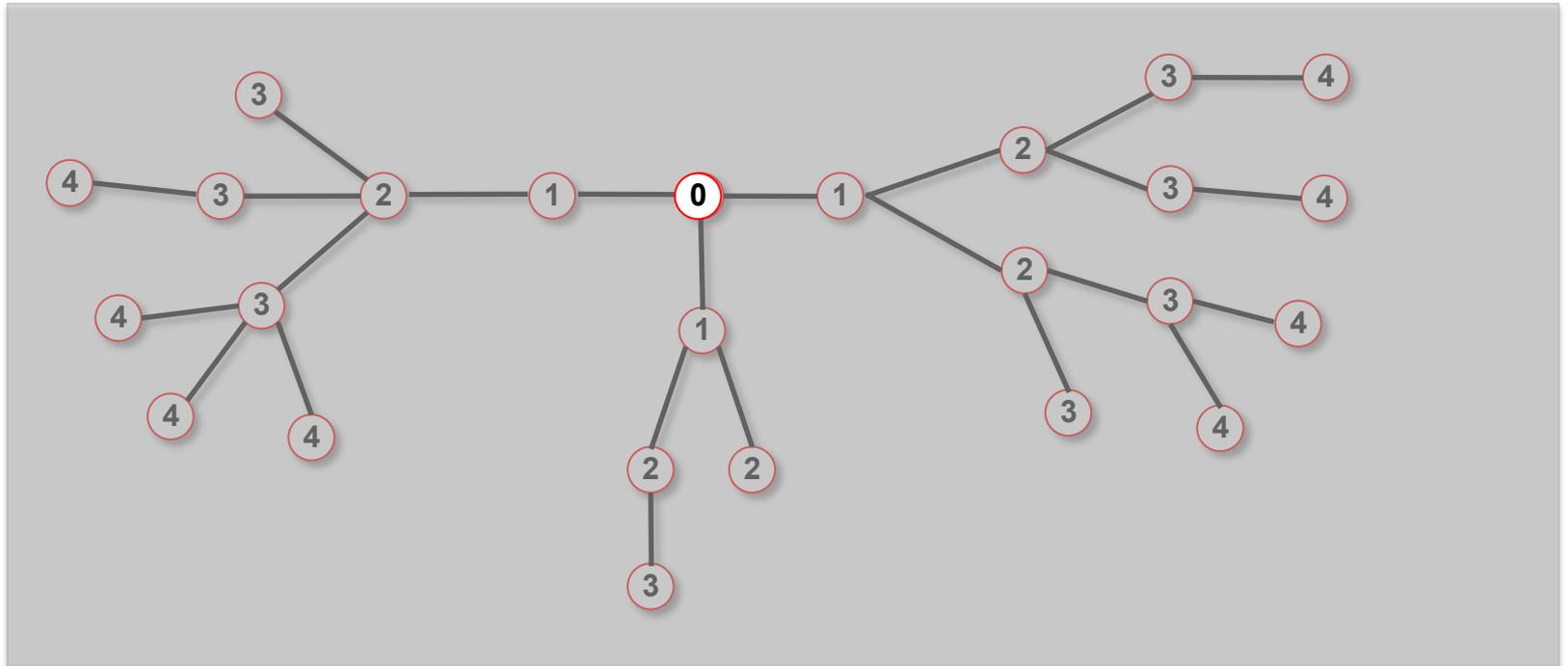
$$N_{ij}^{(n)} = [A^n]_{ij}$$

* holds for both directed and undirected networks.

FINDING DISTANCES: BREADTH FIRST SEARCH

Distance between node 0 and node 4:

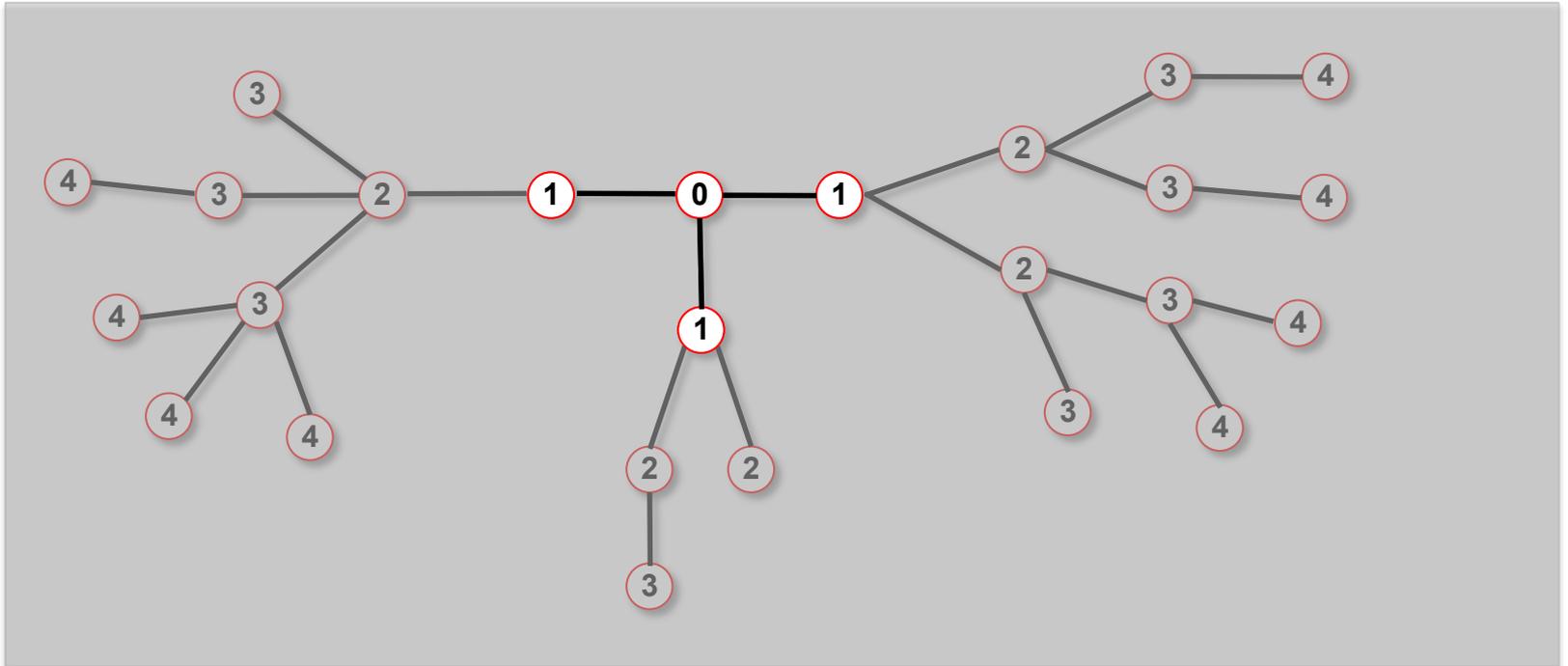
1. Start at 0.



FINDING DISTANCES: BREADTH FIRST SEARCH

Distance between node 0 and node 4:

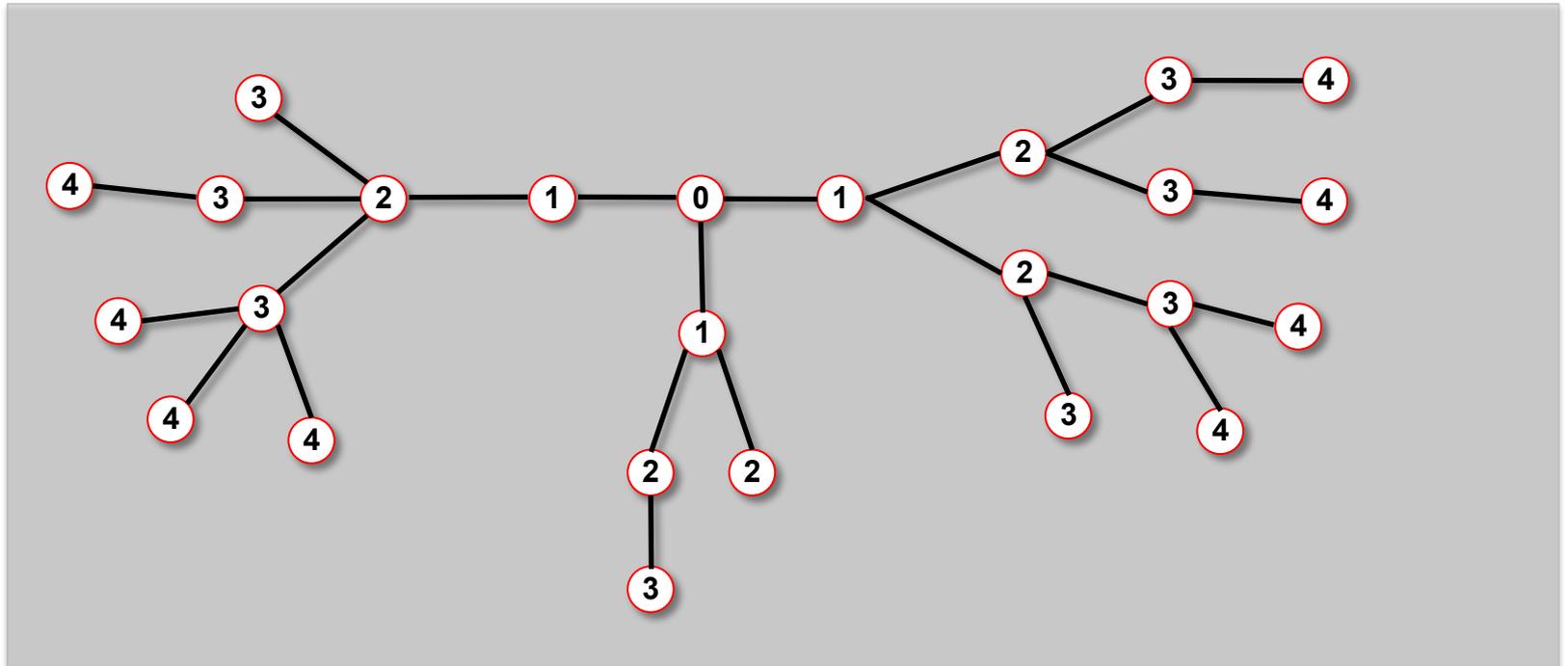
1. Start at 0.
2. Find the nodes adjacent to 1. Mark them as at distance 1. Put them in a queue.



FINDING DISTANCES: BREADTH FIRST SEARCH

Distance between node 0 and node 4:

- 1.Repeat until you find node 4 or there are no more nodes in the queue.
- 2.The distance between 0 and 4 is the label of 4 or, if 4 does not have a label, infinity.



NETWORK DIAMETER AND AVERAGE DISTANCE

Diameter: d_{max} the maximum distance between any pair of nodes in the graph.

Average path length/distance, $\langle d \rangle$, for a **connected graph**:

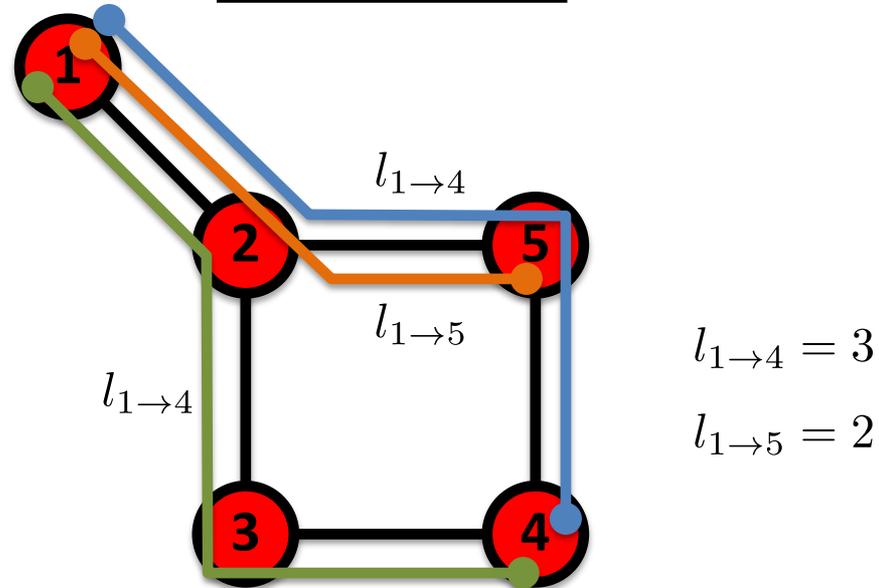
where d_{ij} is the distance from node i to node j

$$\langle d \rangle \equiv \frac{1}{2L_{\max}} \sum_{i,j \neq i} d_{ij}$$

In an *undirected graph* $d_{ij} = d_{ji}$, so we only need to count them once:

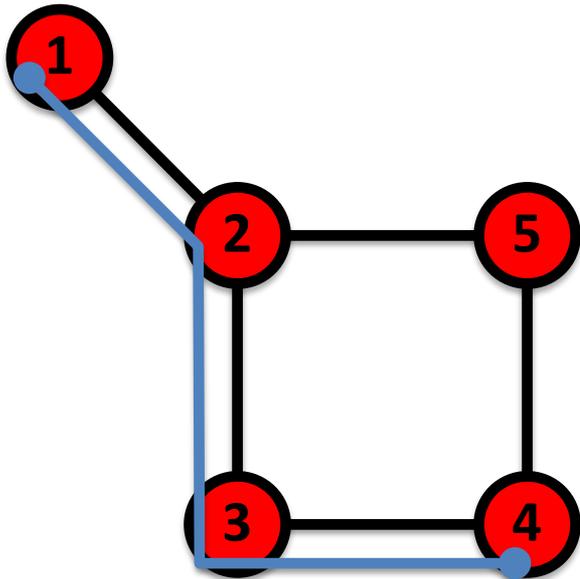
$$\langle d \rangle \equiv \frac{1}{L_{\max}} \sum_{i,j > i} d_{ij}$$

Shortest Path



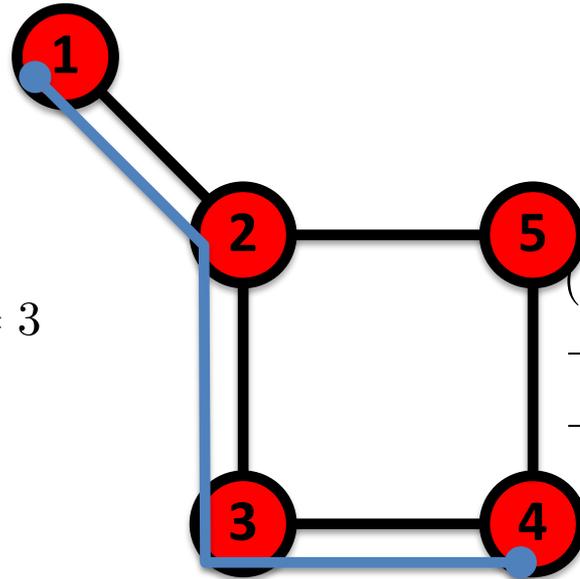
The path with the shortest length between two nodes (distance).

Diameter



The longest shortest path in a graph

Average Path Length

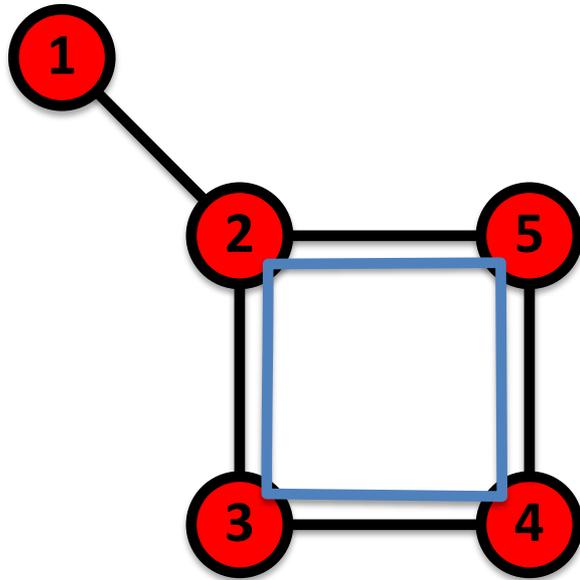


$$l_{1 \rightarrow 4} = 3$$

$$(l_{1 \rightarrow 2} + l_{1 \rightarrow 3} + l_{1 \rightarrow 4} + l_{1 \rightarrow 5} + l_{2 \rightarrow 3} + l_{2 \rightarrow 4} + l_{2 \rightarrow 5} + l_{3 \rightarrow 4} + l_{3 \rightarrow 5} + l_{4 \rightarrow 5}) / 10 = 1.6$$

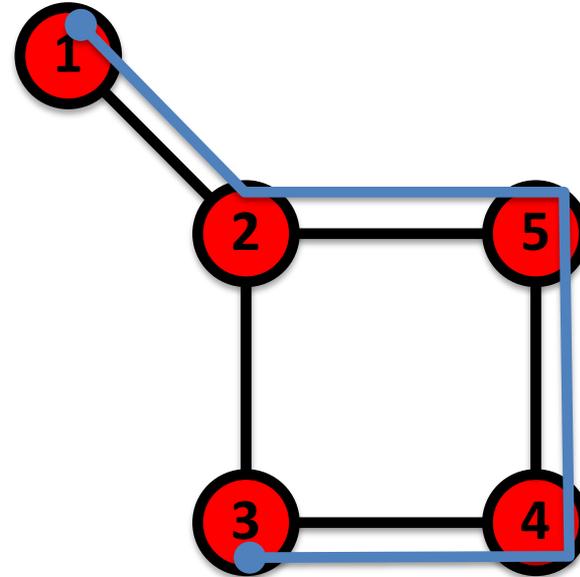
The average of the shortest paths for all pairs of nodes.

Cycle



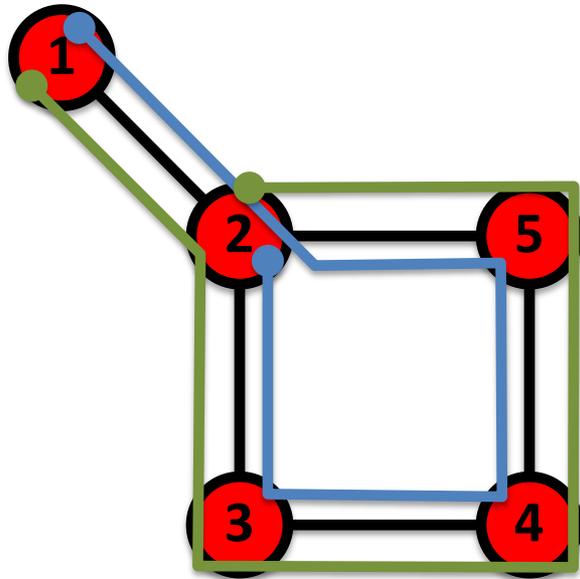
A path with the same start and end node.

Self-avoiding Path



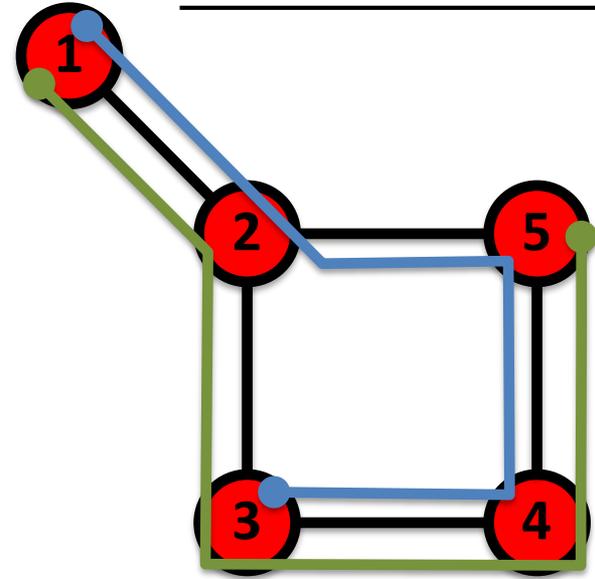
A path that does not intersect itself.

Eulerian Path



A path that traverses each link exactly once.

Hamiltonian Path



A path that visits each node exactly once.