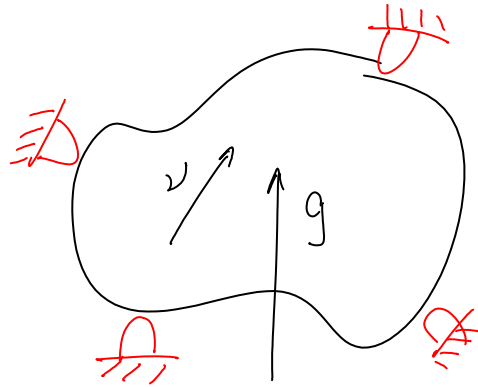


# Form Closure: First-Order

Thursday, January 31, 2008  
9:12 AM

Definition 1: a grasp has form closure iff it cannot be moved when the fingers are locked

$$G_n^T v \geq 0 \Rightarrow v = 0$$



Definition 2: a grasp has form closure iff it is possible for frictionless fingers to generate any wrench.

$$\left. \begin{matrix} G_n \lambda_n = g \\ \lambda_n \geq 0 \end{matrix} \right\} \forall g \in \mathbb{R}^{n_v} \equiv \boxed{\exists \lambda_n > 0 \Rightarrow G \lambda_n = 0}$$

*with use this for computation*

Necessary condition  $\text{rank}(G_n) = n_v$  !

Necessary condition  $n_c \geq n_v + 1$

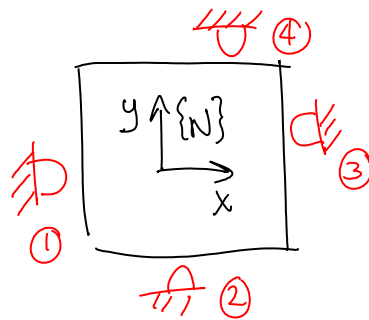
$\therefore$  Need at least 7 contacts to form close a 3D object.

Need at least 4 contacts to form close a 2D object.

Example: Object in  $\mathbb{R}^2$

$$v = \begin{bmatrix} \dot{x} \\ \dot{y} \end{bmatrix} \leftarrow \begin{matrix} \text{assume} \\ \text{translation} \\ \text{only} \end{matrix} v \in \mathbb{R}^2$$

$$G_n = \begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix}$$



Twist Test

$$G_n^T v \geq 0 \Rightarrow \left. \begin{matrix} \dot{x} \geq 0 \\ \dot{y} \geq 0 \\ -\dot{x} \geq 0 \end{matrix} \right\} \Rightarrow \begin{matrix} \dot{x} = 0 \\ \dot{y} = 0 \end{matrix} \Rightarrow \boxed{\text{Form Closure exists!}}$$

$$\underline{\underline{Form Closure}} \Rightarrow \left. \begin{array}{l} \dot{x} \geq 0 \\ \dot{y} \geq 0 \\ -\dot{x} \geq 0 \\ -\dot{y} \geq 0 \end{array} \right\} \Rightarrow \begin{array}{l} \dot{x} = 0 \\ \dot{y} = 0 \end{array} \Rightarrow \boxed{\text{Form Closure exists!}}$$

Wrench Test

$$G_n \lambda_n = \begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix} \begin{bmatrix} f_{1n} \\ f_{2n} \\ f_{3n} \\ f_{4n} \end{bmatrix} = \begin{bmatrix} f_{1n} - f_{3n} \\ f_{2n} - f_{4n} \end{bmatrix} = \begin{bmatrix} -F_x \\ -F_y \end{bmatrix} = -g$$

$$f_{1n}, f_{2n}, f_{3n}, f_{4n} \geq 0$$

Note that given any  $g$ , we can find  $\lambda_n \ni G_n \lambda_n = -g$ !

$\therefore$  Form closure exists.

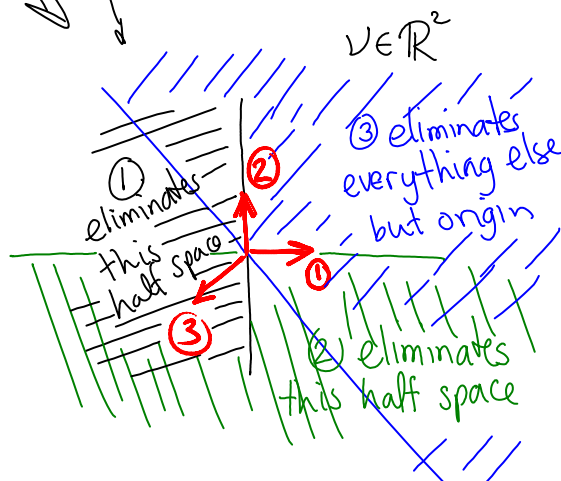
Other interpretations:

Velocity constraints

Each row of  $\tilde{G} v \geq 0$  eliminates a half-space in the twist space

Form closure means that no non-trivial twists are possible!

Assume Triangle can only translate in the plane.



Wrench generation

Non negative span of

columns of  $G_n$   
must equal  $\mathbb{R}^{n \times 2}$ .

Example:

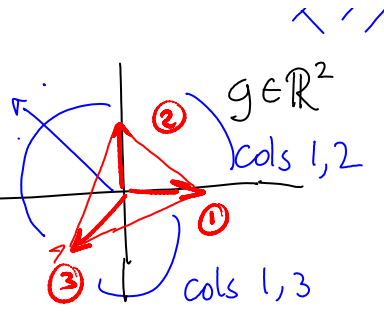
positive span of 1 column  
of  $G_n$

$$G_n \lambda_m, \lambda_m \geq 0$$

this is a ray labeled ①

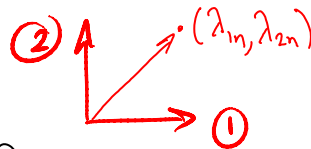
$$\begin{bmatrix} 1 \\ 0 \end{bmatrix} \lambda_m, \lambda_m \geq 0$$

nonnegative  
combinations  
of columns 2,3



Positive span of 2 columns

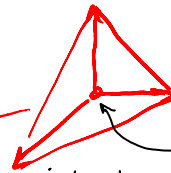
$$\left. \begin{array}{l} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} \lambda_{1n} \\ \lambda_{2n} \end{bmatrix} \\ \lambda_{1n}, \lambda_{2n} \geq 0 \end{array} \right\} = \text{any point in first quadrant of wrench space}$$



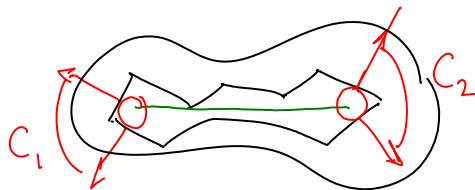
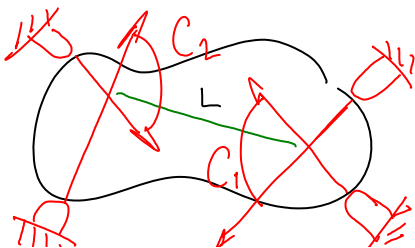
So the positive span includes the two rays  
and everything between them!

Geometric interpretation  
in Wrench Space

Convex hull of columns of  $G_n$  strictly includes origin.



Geometric interpretation  
in work space (planar only)





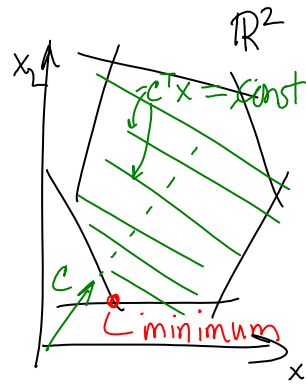
If  $\exists$  Cones,  $C_1 \neq C_2$ , formed with pairs of normals, and line segment  $L$  connecting the cone apexes  $\ni L$  lies entirely in  $C_1 \cap C_2$  or  $-C_1 \cap -C_2$ , then the grasp has form closure.

### Computation Tests for Form Closure.

Linear program

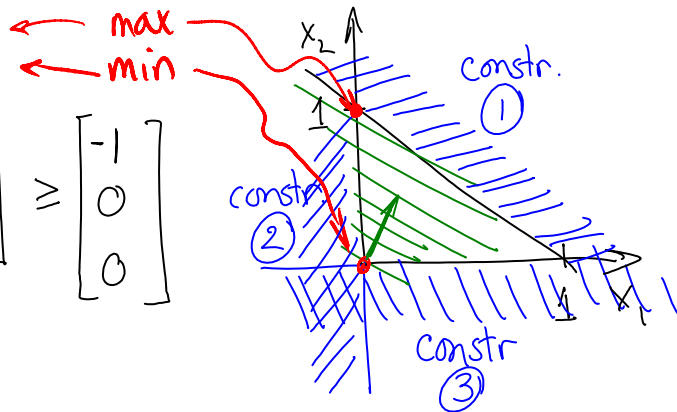
Matlab has good solver.

$$\begin{aligned} \min_x \quad & c^T x \\ \text{s.t.} \quad & Ax \geq b \end{aligned}$$



Example:

$$\begin{aligned} \text{Min} \quad & [1 \ 2] \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \\ \text{s.t.} \quad & \begin{bmatrix} -1 & -1 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \geq \begin{bmatrix} -1 \\ 0 \\ 0 \end{bmatrix} \end{aligned}$$



Form closure requires  $\text{rank}(G) = n_v$  and the existence of  $\lambda_n > 0 \ni G_n \lambda_n = 0$

Second part

We must satisfy  $G_n \lambda_n = 0$   
 $\lambda_n > 0$

How do we "encourage"  $\lambda_n$  to be positive?

$$G_n \lambda_n = 0$$

$$\lambda_n - d \geq 0 \quad \leftarrow \text{"slack variable"}$$

$$d \geq 0$$

$$\lambda_n \geq 0$$

Maximize the slack variable,  $d$ .

L.P.  $\max_{\lambda_n, d} d \leftarrow d \text{ is a measure of how far the grasp is from being form closure}$

s.t:  $G_n \lambda_n = 0$

$$I \lambda_n - 1 d \geq 0$$

$$d \geq 0$$


$$1^T \lambda_n \leq n_c \quad \leftarrow \text{prevent unboundedness}$$

Test for form closure.

- 1 Rank( $G_n$ ) =  $n_v$   
 If not, stop. Form closure does not exist.  
 If yes, continue
- 2 Compute solution to LP1  
 If  $d^* > 0$ , then form closure exists,

Simple example - bead on wire  $\rightarrow$

$$G_n = \begin{bmatrix} 1 & -1 \end{bmatrix}$$



$$\lambda_n = \begin{bmatrix} \lambda_{1n} \\ \lambda_{2n} \end{bmatrix} = \begin{bmatrix} f_{1n} \\ f_{2n} \end{bmatrix}$$

$$\max_{f_{1n}, f_{2n}, d} d$$

s.t.  $f_{1n} - f_{2n} = 0$

$$f_{1n} - d \geq 0$$

Feasible Solutions

$$\begin{array}{l}
 f_{2n} - d \geq 0 \\
 d \geq 0 \\
 f_{1n} + f_{2n} \leq 2
 \end{array}$$

$t_{1n} = t_{2n}$	domain of $d$
0	0
0.5	$[0, 0.5]$
1	$[0, 1.0]$
<del>1.1</del>	<del><math>[0, 1.1]</math></del>

Without last constraint,  $f_{in}$  could increase without bound - infinitely tight squeezing  
 Form closure requires ability to squeeze, which is evident without letting  $d \rightarrow \infty$ .

Without the bound ( $\mathbf{1}^T \lambda_n \leq n_c$ ),  $d^*$  for all form closure grasps would be  $\infty$  and then  $d^*$  could not be used to compare different form closure grasps.

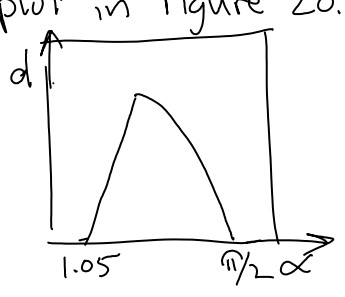
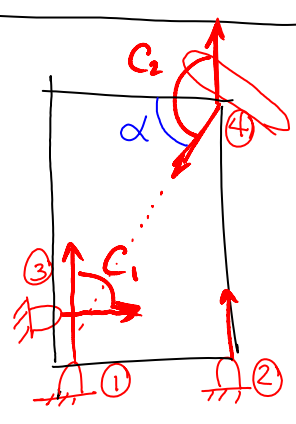
Planar Example

Figure 28.17

Form closure exists for  $1.052 < \alpha < \pi/2$

No form closure if  $\alpha = \pi/2$

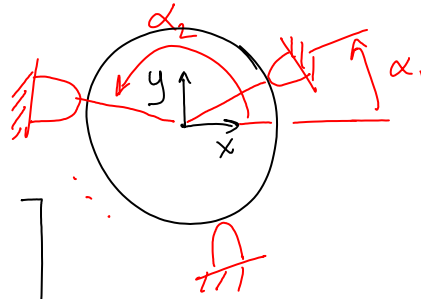
See plot in Figure 28.18



Partial Form Closure - can't prevent all twists, but can prevent a sub-

space of all twists.

Example: Disk  $\rightarrow$



$$G_n = \begin{bmatrix} -\cos(\alpha_1) & -\cos(\alpha_2) & \dots \\ -\sin(\alpha_1) & -\sin(\alpha_2) & \dots \\ 0 & 0 & \dots \end{bmatrix}$$

$\max(\text{Rank}(G_n)) = 2 \implies$  greatest number of d.o.f. constrained is 2.

For this problem, as long as the angle between each pair of contacts is less than  $\pi$ , the disk is form closed in the  $(v_x, v_y)$  subspace of  $v$ .

Can't form close  $w_z$  dimension  
with any # of contacts.