## Manipulation

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## Manipulation under plurality rule (lexicographic tie-breaking)



Strategic behavior (of the
agents)

- Manipulation: an agent (manipulator) casts a vote that does not represent her true preferences, to make herself better off
- A voting rule is strategy-proof if there is never a (beneficial) manipulation under this rule


## Using Borda?



- Inverse the tie-breaking order?


## Using STV?



- $\mathrm{N}>\mathrm{M}>\mathrm{O} \rightarrow \mathrm{O}>\mathrm{M}>\mathrm{N}$


## Any strategy-proof voting rule?

-. No reasonable voting rule is strategyproof

- Gibbard-Satterthwaite Theorem [Gibbard Econometrica-73, Satterthwaite JET-75]: When there are at least three alternatives, no voting rules except dictatorships satisfy
- non-imposition: every alternative wins for some profile
- unrestricted domain: voters can use any linear order as their votes
- strategy-proofness
- Axiomatic characterization for dictatorships!
- Randomized version [Gibbard Econometrica-77]


## A few ways out

- Relax non-dictatorship: use a dictatorship
- Restrict the number of alternatives to 2
- Relax unrestricted domain: mainly pursued by economists
- Single-peaked preferences:
- Range voting: A voter submit any natural number between 0 and 10 for each alternative
- Approval voting: A voter submit 0 or 1 for each alternative


## Single-peaked preferences

- There exists a social axis S
- linear order over the alternatives
- Each voter's preferences V are compatible with the social axis $S$
- there exists a "peak" $a$ such that
- [ $b<c<a$ in S] implies [ $c>b$ in V]
- [a>c>b in S] implies [c>b in V]
- alternatives closer to the peak are more preferred
- different voters may have different peaks


## Examples

rank
Single-peaked preferences


## Strategy-proof rules for single-peaked preferences

- The median rule
- given a profile of "peaks"
- choose the median in the social axis
- Theorem. The Median rule is strategy-proof.
- The median rule with phantom voters
- parameterized by a fixed set of "peaks" of phantom voters
- chooses the median of the peaks of the regular voters and the phantom voters
- Theorem. Any strategy-proof rule for single-peaked preferences are median rules with phantom voters
- Talk announcement: Dominik Peters 9/21 3-4pm Sage 3713


## Computational thinking

- Use a voting rule that is too complicated so that nobody can easily predict the winner
- Dodgson
- Kemeny
- The randomized voting rule used in Venice Republic for more than 500 years [Walsh\&Xia AAMAS-12]
- We want a voting rule where
- Winner determination is easy
- Manipulation is hard
- The hard-to-manipulate axiom: manipulation under the given voting rule is NP-hard


## Example 3: Venetian election

 (1268-1797)

- Round 2:

- Round 3 :

> Approval like


## Manipulation: A computational complexity perspective <br> - ${ }_{\text {er }}$ - If it is computationally too hard for a

 manipulator to compute a manipulation, she is best off voting truthfully- Similar as in cryptography
(? For which common voting rules manipulation is computationally hard?


## Unweighted coalitional manipulation (UCM) problem

- Given
- The voting rule $r$
- The non-manipulators' profile $P^{N M}$
- The number of manipulators $n$,
- The alternative $c$ preferred by the manipulators
- We are asked whether or not there exists a profile $P^{M}$ (of the manipulators) such that $c$ is the winner of $P^{N M} \cup P^{M}$ under $r$


## The stunningly big table for UCM

| \#manipulators | One manipulator | At least two |
| :---: | :---: | :---: |
| Copeland | P [BTT SCW-89b] | NPC [FHS AAMAS-08,10] |
| STV | NPC [BO SCW-91] | NPC [BO SCW-91] |
| Veto | P [ZPR AIJ-09] | P [ZPR AIJ-09] |
| Plurality with runoff | P [ZPR AIJ-09] | P [ZPR AIJ-09] |
| Cup | P [CSL JACM-07] | P [CSL JACM-07] |
| Borda | P [BTT SCW-89b] | NPC [DKN+ AAAI-11] <br> [BNW IJCAI-11] |
| Maximin | P [BTT SCW-89b] | NPC [XZP+ IJCAI-09] |
| Ranked pairs | NPC [XZP+ IJCAI-09] | NPC [XZP+ IJCAI-09] |
| Bucklin | P [XZP+ IJCAI-09] | P [XZP+ IJCAI-09] |
| Nanson's rule | NPC [NWX AAA-11] | NPC [NWX AAA-11] |
| Baldwin's rule | NPC [NWX AAA-11] | NPC [NWX AAA-11] |

## What can we conclude?

- For some common voting rules, computational complexity provides some protection against manipulation
- Is computational complexity a strong barrier?
- NP-hardness is a worst-case concept

