CS 2400 – Models of Computation

Project 2 – Yacc

Return Day: Friday, November 30.
Teams: You must form teams of 2 or 3 persons.

1 Overview

In this assignment you will use the parser generator yacc to construct an interpreter for the Snail programming language which we describe below. The interpreter executes the statements of a Snail program in sequence as they appear in the program.

In Section 2 we describe the Snail programming language. In Section 3 we give the grammar of the Snail language. In Section 4 we describe what your yacc code will do. In Section 5 are instructions to hand-in your assignment.

2 Snail Programming Language

Snail is a very simple programming language. The body of a Snail program consists from a sequence of statements. There are four kinds of statements: assign, print, if, and while. A basic component for all kinds of statements is the expression. The expression and the statements are described below.

- expression
  - An expression is any mathematical expression made from identifiers, integers, parenthesis, the arithmetic operators
  
    + - * /

  - and the comparison operators

    < > <= >= == !=

  - For example, this is a valid expression:

    10 + 20 * (10 < 3)

  - The value of an expression is obtained by executing all the arithmetic operations in the expression. The result of a comparison operation is 1 if the comparison result is true, and 0 otherwise. For example, the above expression has value 10 (since, 10 < 3 = 0).

  - The value of an identifier is the last value assigned to it in an assign statement. An identifier which hasn’t been assigned a value before cannot be used inside an expression and in this case you should report an error message.
• **assign statement**
  The assign statement has the form:

  ```plaintext
  identifier = expression;
  ```

  For example, this is a valid assign statement:

  ```plaintext
  var1 = 20 - 3*2;
  ```

  In the assign statement the identifier gets the value of the expression. As an example, in the above assign statement the new value of variable `var1` is 14.

• **print statement**
  The print statement print messages on the screen. The print statement has one of following forms:

  ```plaintext
  print `string`; //prints the string
  print newline;  //prints a newline character
  print expression; //prints the expression value
  ```

  For example, the execution of the following statements

  ```plaintext
  print `The value of 10*5 is ';
  print 10*5;
  ```

  produces the output:

  The value of 10*5 is 50

• **if statement**
  An if statement has two forms:

  ```plaintext
  if expression then //if-then statement
  statement
  statement
  ...
  endif

  if expression then //if-then-else statement
  statement
  ...
  endif
  ```
The “if-then” statement means that if the expression value is not 0 then
the sequence of statements between then and endif will be executed. The
“if-then-else” statement means that if the expression value is not 0 then
the statements between then and else will be executed, and otherwise, if
the expression value is 0, the statements between else and endif will be
executed. For example, the following is a valid if statement:

```c
if (x < 10) then
    print ‘x is smaller than ten’;
    x = x - y + 20;
else
    x = 10* y;
endif
```

- **while statement**

A while statement has the following form:

```c
while expression do
    statement
    statement
    ...
endwhile
```

The while statement implements a loop which executes the statements
between do and endwhile for as long as the expression value is not 0. As
an example the following while statement will iterate for five times:

```c
v = 1;
while v <= 5 do
    print v;       // print the value of v
    print newline;
    v = v + 1;     // increase v by 1
endwhile
```

A Snail program is a sequence of statements and has the following general form:

```c
statement
statement
statement
...
statement
```

We can have comments in a Snail program right after “//” (as in a C++ pro-
gram). An simple example Snail program is the following:
v = 10;
i = 0;

while i <= v do
    print i*i;  // print the square of i
    if i == v/2 then  // is i the half of v?
        print newline;  //yes
    else
        print ‘-’;  //no
    endif
    i = i + 1;
endwhile

print newline;
print ‘end of execution’;
print newline;

The output of the program is:
0--1--4--9--16--25
36--49--64--81--100--
end of execution

3 Snail Grammar

All the Snail programs can be described by the the context-free grammar of Figure 1. The start variable is program, the grammar variables are in small letters, and the terminals in capital letters. Notice that although this grammar is ambiguous in the expr variable, all the ambiguities can be removed using the precedence rules of yacc.

4 Yacc Code

You will write a yacc code which implements the interpreter for Snail programs. The main part of your yacc code will consist from the snail grammar. You will add actions to the grammar so that your interpreter does the following for any input Snail program:

1. builds the derivation tree of the program, and then
2. “executes” the derivation tree.

The derivation tree (see Chapter 5 in Book) has a node for each variable and terminal. At the root of the tree is the variable program. An example Snail program and derivation tree is shown in Figure 2.
program -> stmt_list

stmt_list -> stmt_list stmt
    | stmt

stmt -> assign_stmt
    | print_stmt
    | if_stmt
    | while_stmt

assign_stmt -> ID = expr ;

print_stmt -> PRINT expr ;
    | PRINT string ;
    | PRINT NEWLINE ;

if_stmt -> IF expr THEN stmt_list ENDIF
    | IF expr THEN stmt_list ELSE stmt_list ENDIF

while_stmt -> WHILE expr DO stmt_list ENDWHILE

expr -> ( expr )
    | expr + expr
    | expr - expr
    | expr * expr
    | expr / expr
    | expr < expr
    | expr > expr
    | expr <= expr
    | expr >= expr
    | expr == expr
    | expr != expr
    | - expr
    | INT
    | ID

Figure 1: The snail grammar
Program: print 10+5;

Figure 2: A small Snail program and its derivation tree
You will build the tree bottom up starting from the leaves. (The reason of
the bottom up construction is that yacc gives a bottom up parser.) To build
the derivation tree you need a special routine, e.g. add_node, which you will
invoke at each production of your grammar and will add a node (or nodes)
in the derivation tree. Your nodes of your tree must be general enough to
accommodate all the different kinds of productions, variables and terminals
in the grammar. You need a mechanism to distinguish between the various kinds
of nodes. (for example, you can have a variable kind inside each node). You
don’t need to have a node for each terminal in your grammar, e.g. you don’t
need nodes for semicolons ‘;’.

By “executing the tree” we mean that we traverse the tree recursively from
the root to the leaves and execute the code that corresponds to each node of
the tree. For this you will need to write a special routine, e.g. execute_tree,
which you will invoke after you build the derivation tree. (Normally you will
invoke the execute_tree routine at an action of the program variable of your
grammar.) The main part of execute_tree is a big switch statement for the
various kinds of nodes. The pseudo-code for execute_tree is as follows:

// tree is the derivation tree

execute_tree(tree) {
   root = root(tree);
   left_child = root.left_child;
   middle_child = root.middle_child
   right_child = root.right_child;

   switch (root.kind) {
      case expr_plus: execute_tree(left_child);
                      execute_tree(right_child);
                      root.value = left_child.value + right_child.value;

      case expr_times: ...... 
                     ...... 

      case print_string: execute_tree(middle_child);
                         printf("%s", middle_child.value);

      case print_expr : execute_tree(middle_child);
                         printf("%d", middle_child.value);
                     ...... 
   }
}

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Each \texttt{expr} node must have a value variable which holds the current value of the expression. The \texttt{execute_tree} routine computes the \texttt{expr} values recursively, by computing the values of the children \texttt{expr} first.

The value of an \texttt{ID} (identifier) node can be stored in the symbol table. For an \texttt{assign} node we update the value of the \texttt{ID} child in the symbol table to be equal to the value of the \texttt{expr} child.

For a \texttt{print} node we just print the contents (or value) of the middle child, which can be either a \texttt{STRING}, a \texttt{expr} or \texttt{NEWLINE}.

For an \texttt{if} node, we first execute the \texttt{expr} child and then if the value of \texttt{expr} is not 0 we execute the if-then \texttt{stmt_list} child. Otherwise, we execute the if-else \texttt{stmt_list} child.

For a \texttt{while} node, we repeatedly do the following: first we execute the \texttt{expr} child and if the value of \texttt{expr} is not 0 we execute the \texttt{stmt_list} node. When the value of \texttt{expr} is zero the execution of the \texttt{while} node has finished.

If you find a syntax error in the input program then report an error message with the line number where you found the error and abort the program.

Your yacc program will use the lexical analyzer of the first project and for this you need to modify appropriately the lex code.

5 Hand-In

You should hand-in in paper your lex and yacc code. Also you should hand-in the output of your interpreter for five Snail programs which are given in the course web page.

In the course web page you can also find example yacc programs (together with lex programs) that may help you to get started with your project.