Fall 2012, Test 1, Introduction to Algorithms

Name:
Section:
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29th September, 2014

Answer all Six questions. You have 90 minutes to complete the exam. You can use a calculator - You are allowed to use your text books and notes.
1. **Recursion and Design and Analysis**

(a) Write an algorithm (C++ like program) to find the smallest integer greater than or equal to \( \sqrt{n} \) of a given input positive integer, \( n \). (You are not allowed to use built in mathematical libraries. You are allowed to use any of the standard (+, -, *, /) arithmetical operators)

For example if \( n = 3 \), the output should be 2, if \( n = 16 \), the output should be 4 and if \( n = 69 \), the output should be 9. What is the running time of your algorithm in \( O \) notation. **[5 points]**

(b) Given a set, \( S \), of \( n \) integer elements (e.g., \( n = 6 \) and \( S = \{-2, -92, 12, 4, -15, -3\} \), the numbers -2, -2, 4 add up to 0. ), write an algorithm to find whether any three elements (not necessarily distinct) add up to 0. **[A pseudo code will suffice]** What is the complexity of your algorithm. **[5 points]**
2. Recurrence Relations, Master Theorem

(a) Match the following Recurrence Relations with the solutions given below. (give tight bounds) [5 Points]

i. $P(n) = 27P\left(\frac{n}{3}\right) + 2, \quad P(1) = 1$

ii. $Q(n) = Q(n-1) + n, Q(1) = 1$

iii. $R(n) = R\left(\frac{n}{2}\right) + \frac{1}{2}, R(1) = 1$

iv. $S(n) = 4S\left(\frac{n}{4}\right) + n, \quad S(1) = 1$

v. $T(n) = T(n/2) + n, T(1) = 1$

Solutions

i. $O(\log(n))$

ii. $O(n^3 \log(n))$

iii. $O(n)$

iv. $O(n^2)$

v. $O(n^3)$

vi. $O(n \log(n))$

(b) How many print statements (cout) will be execute by the following program when $n = 2, n = 3$ and $n = 7$. Write a recurrence relation (not solve) for the number of times print statements are executed for a general $n$. [5 Points]
void csci2300(int n)
{
    if (n<=1) cout << "'Scientia et Instantia'" << endl;
    else
    {
        cout << "'Knowledge and Thoroughness'" << endl;
        csci2300(n-2);
        cout << "'Why not Change the World'" << endl;
        csci2300(n-2);
        cout << "'Let us Go Red'" << endl;
    }
    return;
}
3. **Divide and Conquer and Design of Algorithm**

Given an array of integers. Find a peak element in it. An array element is peak if it is NOT smaller than its neighbors. For corner elements, we need to consider only one neighbor. For example, for input array {15, 16, 22, 18}, 22 is the only peak element. For input array {11, 21, 16, 3, 24, 91, 68}, there are two peak elements: 21 and 91. Note that we need to return any one peak element. A linear time algorithm will get you only 5 points.

(a) Provide a pseudo code description of an algorithm to find the peak element. [5 points]

(b) What is the complexity of your algorithm. [5 points]
4. Algorithms with Numbers

(a) Find $x$ and $y$ such that $87 \times x + 19 \times y = \gcd(87, 19)$ and use it compute multiplicative inverse of 19 mod (87) [6 points]

(b) what is the value of $3^{8^{128}} \mod 85$ (Hint: 85 is a product of two prime numbers 5 and 17). You have to show your work (just writing the answer gives 2 points) and what theorems/lemmas you used. [4 points]
5. **Recursive full binary Trees**

A full binary tree is a tree that has either 0 or 2 children (both left and right) at each vertex. (This is a **Full Binary tree**). Let \( B_n \) denote the number of full Binary trees with \( n \) vertices. Naturally \( B_1 = 1 \). By definition \( B_2 = 0 \). In fact the number of full Binary Trees with even number of nodes is 0, i.e., \( B_{2n} = 0 \).

(a) Draw Binary trees with 5 and 7 vertices, determine the exact values of \( B_5, B_7 \) [6 points].

(b) For general \( 2n + 1 \), derive a recurrence relation (and not solve) for \( B_{2n+1} \) [4 points] (Hint: There are \( 2i+1 \) nodes in the left subtree, \( 2n - 1 - 2i \) nodes in the right subtree)
6. Graphs and Graph Algorithms

(a) For a directed graph $G$ with node set $\{2,3,4,5,6,7,8,9,10,11,12\}$ with a directed edge from node $i$ to node $j$, if $i$ properly (remainder is zero) divides $j$. (By definition there are no edges from $i$ to $i$ - it is not allowed for this problem). What are the out-degrees of node 2 and node 4. What are the indegrees of node 12 and node 11. List all nodes with in-degree 0. [6 points]

(b) Construct a graph $G$, with 6 nodes and 5 edges such that the pre and post numbers for all but two of the nodes are differ by at least 5 (i.e. for 4 of the 6 nodes, post-numbering is at least 5 more than the pre-numbering. i.e, for a node $n$, $\text{post}(n) > \text{pre}(n)+4$). (You can choose a starting node and the order of traversal - show which is your starting node and the order of traversal) [4 points].