Project 2 – Yacc

Return Day: Friday, April 25.
Teams: You must form teams of 2 or 3 persons.

1 Overview

In this assignment you will use the parser generator yacc to construct an interpreter for the SmILLY programming language which we describe below. The interpreter executes the statements of a SmILLY program in sequence as they appear in the program.

In Section 2 we describe the SmILLY programming language. In Section 3 we give the grammar of the SmILLY language. In Section 4 we describe what your yacc code will do. In Section 5 are instructions to hand-in your assignment.

2 SmILLY Programming Language

SmILLY is a very simple programming language. The body of a SmILLY program consists from a sequence of statements. There are four kinds of statements: assign, print, if, and do. A basic component for all kinds of statements is the expression. The expression and the statements are described below.

• expression

An expression is any mathematical expression made from identifiers, integers, parenthesis, the arithmetic operators

+ - * / +=

and the comparison operators

< > <= >= == !=

For example, this is a valid expression:

10 + 20 * (10 < 3)

The value of an expression is obtained by executing all the arithmetic operations in the expression. The result of a comparison operation is 1 if the comparison result is true, and 0 otherwise. For example, the above expression has value 10 (since, 10 < 3 = 0).

The value of an identifier is the last value assigned to it in an assign statement. An identifier which hasn’t been assigned a value before cannot be used inside an expression and in this case you should report an error message.
• **assign statement**
  The assign statement has the form:
  
  ```
  identifier = expression;
  ```
  
  For example, this is a valid assign statement:
  
  ```
  var1 = 20 - 3*2;
  ```
  
  In the assign statement the identifier gets the value of the expression. As an example, in the above assign statement the new value of variable `var1` is 14.

• **print statement**
  The print statement print messages on the screen. The print statement has one of following forms:
  
  ```
  print ‘string’;  //prints the string
  print newline;  //prints a newline character
  print expression;  //prints the expression value
  ```
  
  For example, the execution of the following statements:
  
  ```
  print ‘The value of 10*5 is ’;
  print 10*5;
  ```
  
  produces the output:
  
  The value of 10*5 is 50

• **if statement**
  An if statement has two forms:
  
  ```
  if expression then  /*if-then statement*/
  statement
  statement
  ...
  /*more statements */
  endif
  ```
  
  ```
  if expression then  /*if-then-else statement*/
  statement
  ...
  /*more statements*/
  else
  statement
  ...
  /*more statements*/
  endif
  ```
The “if-then” statement means that if the expression value is not 0 then the sequence of statements between then and endif will be executed. The “if-then-else” statement means that if the expression value is not 0 then the statements between then and else will be executed, and otherwise, if the expression value is 0, the statements between else and endif will be executed. For example, the following is a valid if statement:

```plaintext
if (x < 10) then
    print 'x is smaller than ten';
    x = x - y + 20;
else
    x = 10*y;
endif
```

- do statement

A do statement has the following form:

```plaintext
do while expression begind
    statement
    statement
    ...
dendo
```

The do statement implements a loop which executes the statements between begind and enddo for as long as the expression value is not 0. As an example the following while statement will iterate for five times:

```plaintext
v = 1;
do while v <= 5 begind
    print v;       /* print the value of v*/
    print newline;
    v = v + 1;      /* increase v by 1 */
dendo
```

A SmILLY program is a sequence of statements and has the following general form:

```plaintext
statement
statement
...
statement
```

We can have comments in a SmILLY program with /* ... */ as in C. An simple example SmILLY program is the following:
v = 10;
I = 0;

do while i <= v begin do
   print i * i; /* print the square of i */
   if i == v/2 then /* is i the half of v */
      print newline; /* yes */
   else
      print ‘--’; /* no */
   endif
   i = i + 1;
end do

print newline;
print ‘‘end of execution’’;
print newline;

The output of the program is:
0--1--4--9--16--25
36--49--64--81--100--
end of execution

3 SmILLY Grammar

All the SmILLY programs can be described by the the context-free grammar of Figure 1. The start variable is program, the grammar variables are in small letters, and the terminals in capital letters. Notice that although this grammar is ambiguous in the expr variable, all the ambiguities can be removed using the precedence rules of yacc.

4 Yacc Code

You will write a yacc code which implements the interpreter for SmILLY programs. The main part of your yacc code will consist from the SmILLY grammar. You will add actions to the grammar so that your interpreter does the following for any input SmILLY program:

1. builds the derivation tree of the program, and then
2. “executes” the derivation tree.

The derivation tree (see Chapter 5 in Book) has a node for each variable and terminal. At the root of the tree is the variable program. An example SmILLY program and derivation tree is shown in Figure 2.
program -> stmt_list

stmt_list -> stmt_list stmt
            | stmt

stmt -> assign_stmt
      | print_stmt
      | if_stmt
      | do_stmt

assign_stmt -> ID = expr ;

print_stmt -> PRINT expr ;
             | PRINT string ;
             | PRINT NEWLINE ;

if_stmt -> IF expr THEN stmt_list ENDIF
          | IF expr THEN stmt_list ELSE stmt_list ENDIF

do_stmt -> DO WHILE expr BEGINDO stmt_list ENDDO

expr -> ( expr )
      | expr + expr
      | expr - expr
      | expr * expr
      | expr / expr
      | expr < expr
      | expr > expr
      | expr <= expr
      | expr >= expr
      | expr == expr
      | expr != expr
      | - expr
      | INT
      | ID

Figure 1: The SmILLY grammar
Program: print 10+5;

Figure 2: A small SmILLY program and its derivation tree
You will build the tree bottom up starting from the leaves. (The reason of the bottom up construction is that yacc gives a bottom up parser.) To build the derivation tree you need a special routine, e.g. add_node, which you will invoke at each production of your grammar and will add a node (or nodes) in the derivation tree. Your nodes of your tree must be general enough to accommodate all the different kinds of productions, variables and terminals in the grammar. You need a mechanism to distinguish between the various kinds of nodes. (for example, you can have a variable kind inside each node). You don’t need to have a node for each terminal in your grammar, e.g. you don’t need nodes for semicolonns ‘;’.

By “executing the tree” we mean that we traverse the tree recursively from the root to the leaves and execute the code that corresponds to each node of the tree. For this you will need to write a special routine, e.g. execute_tree, which you will invoke after you build the derivation tree. (Normally you will invoke the execute_tree routine at an action of the program variable of your grammar.) The main part of execute_tree is a big switch statement for the various kinds of nodes. The pseudo-code for execute_tree is as follows:

```c
/* tree is the derivation tree */

execute_tree(tree) {
    root = root(tree);
    left_child = root.left_child;
    middle_child = root.middle_child
    right_child = root.right_child;

    switch (root.kind) {
    case expr_plus:   execute_tree(left_child);
                     execute_tree(right_child);
                     root.value = left_child.value + right_child.value;

        case expr_times: ......
                          ......

        case print_string: execute_tree(middle_child);
                             printf(‘%s’, middle_child.value);

        case print_expr:  execute_tree(middle_child);
                             printf(‘%d’, middle_child.value);

                          ......
    }
    }
```
Each $\mathit{expr}$ node must have a value variable which holds the current value of the expression. The $\text{execute\_tree}$ routine computes the $\mathit{expr}$ values recursively, by computing the values of the children $\mathit{expr}$ first.

The value of an $\mathit{ID}$ (identifier) node can be stored in the symbol table. For an $\mathit{assign}$ node we update the value of the $\mathit{ID}$ child in the symbol table to be equal to the value of the $\mathit{expr}$ child.

For a $\mathit{print}$ node we just print the contents (or value) of the middle child, which can be either a $\mathit{STRING}$, a $\mathit{expr}$ or $\mathit{NEWLINE}$.

For an $\mathit{if}$ node, we first execute the $\mathit{expr}$ child and then if the value of $\mathit{expr}$ is not 0 we execute the if-then $\mathit{stmt\_list}$ child. Otherwise, we execute the if-else $\mathit{stmt\_list}$ child.

For a $\mathit{do}$ node, we repeatedly do the following: first we execute the $\mathit{expr}$ child and if the value of $\mathit{expr}$ is not 0 we execute the $\mathit{stmt\_list}$ node. When the value of $\mathit{expr}$ is zero the execution of the $\mathit{do}$ node has finished.

If you find a syntax error in the input program then report an error message with the line number where you found the error and abort the program.

Your yacc program will use the lexical analyzer of the first project and for this you need to modify appropriately the lex code.

5 Hand-In

You should hand-in in paper your lex and yacc code. Also you should hand-in the output of your interpreter for five SmILLY programs (Proj2test1, Proj2test2, Proj2test3, Proj2test4, Proj2test5) which are given in the course web page.

In the course web page you can also find example yacc programs (together with lex programs) that may help you to get started with your project.