

# Higher-Order Programming:

Iterative computation (CTM Section 3.2)

Closures, procedural abstraction, genericity, instantiation,  
embedding (CTM Section 3.6.1)

Carlos Varela

RPI

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Adapted with permission from:

Seif Haridi

KTH

Peter Van Roy

UCL

# Iterative computation

- An iterative computation is one whose execution stack is bounded by a constant, independent of the length of the computation
- Iterative computation starts with an initial state  $S_0$ , and transforms the state in a number of steps until a final state  $S_{final}$  is reached:

$$S_0 \longrightarrow S_1 \longrightarrow \dots \longrightarrow S_{final}$$

# From a general scheme to a control abstraction (1)

```
fun {Iterate  $S_i$ }  
  if {IsDone  $S_i$ } then  $S_i$   
  else  $S_{i+1}$  in  
     $S_{i+1} = \{Transform\ S_i\}$   
    {Iterate  $S_{i+1}$ }  
  end  
end
```

- *IsDone* and *Transform* are problem dependent

# From a general scheme to a control abstraction (2)

```
fun {Iterate S IsDone Transform}
  if {IsDone S} then S
  else S1 in
    S1 = {Transform S}
    {Iterate S1 IsDone Transform}
  end
end
```

```
fun {Iterate  $S_i$ }
  if {IsDone  $S_i$ } then  $S_i$ 
  else  $S_{i+1}$  in
     $S_{i+1} = \{Transform S_i\}$ 
    {Iterate  $S_{i+1}$ }
  end
end
```

# Sqrt using the Iterate abstraction

```
fun {Sqrt X}
  fun {Improve Guess}
    (Guess + X/Guess)/2.0
  end
  fun {GoodEnough Guess}
    {Abs X - Guess*Guess}/X < 0.000001
  end
  Guess = 1.0
in
  {Iterate Guess GoodEnough Improve}
end
```

# Sqrt using the control abstraction

```
fun {Sqrt X}
  {Iterate
    1.0
    fun {$ G} {Abs X - G*G}/X < 0.000001 end
    fun {$ G} (G + X/G)/2.0 end
  }
end
```

Iterate could become a linguistic abstraction

# Sqrt in Haskell

```
let sqrt x = head (dropWhile (not . goodEnough) sqrtGuesses)
```

```
  where
```

```
    goodEnough guess = (abs (x - guess*guess))/x < 0.00001
```

```
    improve guess = (guess + x/guess)/2.0
```

```
    sqrtGuesses = 1:(map improve sqrtGuesses)
```

This sqrt example uses infinite lists enabled by lazy evaluation, and the map control abstraction.

# Higher-order programming

- **Higher-order programming** = the set of programming techniques that are possible with procedure values (lexically-scoped closures)
- Basic operations
  - **Procedural abstraction**: creating procedure values with lexical scoping
  - **Genericity**: procedure values as arguments
  - **Instantiation**: procedure values as return values
  - **Embedding**: procedure values in data structures
- Higher-order programming is the foundation of component-based programming and object-oriented programming



# Procedural abstraction

- Procedural abstraction is the ability to convert any statement into a procedure value
  - A procedure value is usually called a **closure**, or more precisely, a **lexically-scoped closure**
  - A procedure value is a pair: it combines the procedure code with the environment where the procedure was created (the contextual environment)
- Basic scheme:
  - Consider any statement  $\langle s \rangle$
  - Convert it into a procedure value:  $P = \text{proc } \{ \$ \} \langle s \rangle \text{ end}$
  - Executing  $\{P\}$  has **exactly the same effect** as executing  $\langle s \rangle$

# Procedural abstraction

```
fun {AndThen B1 B2}  
  if B1 then B2 else false  
  end  
end
```

# Procedural abstraction

```
fun {AndThen B1 B2}  
  if {B1} then {B2} else false  
  end  
end
```

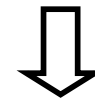
# A common limitation

- Most popular imperative languages (C, Pascal) do **not** have procedure values
- They have only **half** of the pair: variables can reference procedure code, but there is no contextual environment
- This means that **control abstractions cannot be programmed** in these languages
  - They provide a predefined set of control abstractions (for, while loops, if statement)
- Generic operations are still possible
  - They can often get by with just the procedure code. The contextual environment is often empty.
- The limitation is due to **the way memory is managed** in these languages
  - Part of the store is put on the stack and deallocated when the stack is deallocated
  - This is supposed to make memory management simpler for the programmer on systems that have no garbage collection
  - It means that contextual environments cannot be created, since they would be full of dangling pointers
- Object-oriented programming languages can use objects to encode procedure values by making external references (contextual environment) instance variables.

# Genericity

- Replace specific entities (zero 0 and addition +) by function arguments
- The same routine can do the sum, the product, the logical or, etc.

```
fun {SumList L}
  case L
  of nil then 0
  [] X|L2 then X+{SumList L2}
  end
end
```



```
fun {FoldR L F U}
  case L
  of nil then U
  [] X|L2 then {F X {FoldR L2 F U}}
  end
end
```

# Instantiation

```
fun {FoldFactory F U}
  fun {FoldR L}
    case L
    of nil then U
    [] X|L2 then {F X {FoldR L2}}
    end
  end
in
  FoldR
end
```

- Instantiation is when a procedure returns a procedure value as its result
- Calling {FoldFactory fun {\$ A B} A+B end 0} returns a function that behaves identically to SumList, which is an « instance » of a folding function

# Embedding

- Embedding is when procedure values are put in data structures
- Embedding has many uses:
  - **Modules**: a module is a record that groups together a set of related operations
  - **Software components**: a software component is a generic function that takes a set of modules as its arguments and returns a new module. It can be seen as **specifying** a module in terms of the modules it needs.
  - **Delayed evaluation** (also called **explicit lazy evaluation**): build just a small part of a data structure, with functions at the extremities that can be called to build more. The consumer can control explicitly how much of the data structure is built.

# Exercises

15. CTM Exercise 3.10.2 (page 230)
16. CTM Exercise 3.10.3 (page 230)
17. Develop a control abstraction for iterating over a list of elements.
18. CTM Exercise 3.10.5 (page 230)
19. Suppose you have two sorted lists. Merging is a simple method to obtain an again sorted list containing the elements from both lists. Write a Merge function that is generic with respect to the order relation.



# Exercises

20. Instantiate the FoldFactory to create a ProductList function to multiply all the elements of a list.
21. Create an AddFactory function that takes a list of numbers and returns a list of functions that can add by those numbers, e.g. {AddFactory [1 2]} => [Inc1 Inc2] where Inc1 and Inc2 are functions to increment a number by 1 and 2 respectively, e.g., {Inc2 3} => 5.
22. Implement exercises 18-21 in both Oz and Haskell.