Logic Programming

(PLP 11, CTM 9.2, 9.4, 12.1-12.2)

Constraint Satisfaction Problems,
Natural Language Parsing (Definite Clause Grammars)

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Constraint Satisfaction Example*

• Given six Italian words:
  – astante, astoria, baratto, cobalto, pistola, statale.

• They are to be arranged, crossword puzzle fashion, in the following grid:

C. Varela, *Example from Learn Prolog Now! by Blackburn et al. (Exercise 2.4)
The following knowledge base represents a lexicon containing these words:

```
word(astante, a,s,t,a,n,t,e).
word(astoria, a,s,t,o,r,i,a).
word(baratto, b,a,r,a,t,t,o).
word(cobalto, c,o,b,a,l,t,o).
word(pistola, p,i,s,t,o,l,a).
word(statale, s,t,a,t,a,l,e).
```

Write a predicate `crossword/6` that tells us how to fill in the puzzle. The first three arguments should be the vertical words from left to right, and the last three arguments the horizontal words from top to bottom.
Constraint Satisfaction Example(3)*

- Try solving it yourself before looking at this solution!

```prolog
crossword(V1, V2, V3, H1, H2, H3) :-
    word(V1, _, H1V1, _, H2V1, _, H3V1, _),
    word(V2, _, H1V2, _, H2V2, _, H3V2, _),
    word(V3, _, H1V3, _, H2V3, _, H3V3, _),
    word(H1, _, H1V1, _, H1V2, _, H1V3, _),
    word(H2, _, H2V1, _, H2V2, _, H2V3, _),
    word(H3, _, H3V1, _, H3V2, _, H3V3, _).
```

C. Varela, *Example from Learn Prolog Now! by Blackburn et al. (Exercise 2.4)
Constraint Satisfaction Example (Oz)

- The following relation represents the lexicon:

```prolog
fun {Word}
  choice astante#a#s#t#a#n#t#e
  [] astoria#a#s#t#o#r#i#a
  [] baratto#b#a#r#a#t#t#o
  [] cobalto#c#o#b#a#l#t#o
  [] pistola#p#i#s#t#o#l#a
  [] statale#s#t#a#t#a#l#e
end
end
```

- Write a predicate `Crossword/1` that tells us how to fill in the puzzle.
Constraint Satisfaction Example (Oz)

```prolog
proc {Crossword S}
    H1V1 H2V1 H3V1 V1 H1
    H1V2 H2V2 H3V2 V2 H2
    H1V3 H2V3 H3V3 V3 H3

in
    S = [V1 V2 V3 H1 H2 H3]
    {Word V1#_#H1V1#_#H2V1#_#H3V1#_}
    {Word V2#_#H1V2#_#H2V2#_#H3V2#_}
    {Word V3#_#H1V3#_#H2V3#_#H3V3#_}
    {Word H1#_#H1V1#_#H1V2#_#H1V3#_}
    {Word H2#_#H2V1#_#H2V2#_#H2V3#_}
    {Word H3#_#H3V1#_#H3V2#_#H3V3#_}

end
```

C. Varela, *Example from Learn Prolog Now! by Blackburn et al. (Exercise 2.4)
Crossword is a *relation* that corresponds to a query, represented as a one-argument procedure (or equivalent function).

Oz’s Search module can produce a *lazy* list of solutions:

- especially useful when there are infinite answers, or when computation of all answers would take too long.

Solutions can be accessed via a *search engine object*:

```
% search engine
E = {New Search.object script(Crossword)}

% calculate and display one at a time
{Browse {E next($)}}
```

C. Varela, *Example from Learn Prolog Now! by Blackburn et al. (Exercise 2.4)"
The Crossword query relation can also be used directly by the Search module:

% Finding one solution
{Browse {Search.base.one Crossword}}

% Finding all solutions
{Browse {Search.base.all Crossword}}
Generate and Test Example

• We can use the relational computation model to generate all digits:

```plaintext
fun {Digit}
  choice 0 [] 1 [] 2 [] 3 [] 4 [] 5 [] 6 [] 7 [] 8 [] 9 end
end

{Browse {Search.base.all Digit}}
% displays [0 1 2 3 4 5 6 7 8 9]
```
Finding digit pairs that add to 10

- Using generate and test to do combinatorial search:

```plaintext
fun {PairAdd10}
    D1 D2 in
    D1 = {Digit} % generate
    D2 = {Digit} % generate
    D1+D2 = 10 % test
    D1#D2
end

{Browse {Search.base.all PairAdd10}}
% displays [1#9 2#8 3#7 4#6 5#5 6#4 7#3 8#2 9#1]
```
Finding digit pairs that add to 10 (Prolog)

• Using generate and test to do combinatorial search:

```prolog
digit(D) :- between(0,9,D).

pairAdd10(D1,D2) :-
    digit(D1),
    digit(D2),
    D1 + D2 =:= 10.

allPairs(L) :-
    findall(p(D1,D2),pairAdd10(D1,D2),L).
```

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Finding palindromes

• Find all four-digit palindromes that are products of two-digit numbers:

```haskell
fun {Palindrome}
    X in
    X = (10*{Digit}+{Digit})*(10*{Digit}+{Digit}) % generate
    (X>=1000) = true % test
    (X div 1000) mod 10 = (X div 1) mod 10 % test
    (X div 100) mod 10 = (X div 10) mod 10 % test
    X
end

{Browse {Search.base.all Palindrome}} % 118 solutions
```

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Finding palindromes (Prolog)

• Find all four-digit palindromes that are products of two-digit numbers:

\[
\text{palindrome}(S) :-
\]

\[
\text{digit}(D1), \text{digit}(D2), \text{digit}(D3), \text{digit}(D4), \quad \% \text{generate}
\]

\[
S \text{ is } (10\times D1 + D2) \times (10\times D3 + D4),
\]

\[
S \geq 1000, \quad \% \text{test}
\]

\[
\text{mod}(\text{div}(S, 1000), 10) =:= \text{mod}(S, 10), \quad \% \text{1st = 4th}
\]

\[
\text{mod}(\text{div}(S, 100), 10) =:= \text{mod}(\text{div}(S, 10), 10). \quad \% \text{2nd = 3rd}
\]

\[
\text{allPalindromes}(S, L) :- \text{findall}(P, \text{palindrome}(P), S), \text{length}(S, L).
\]
Propagate and Search

• The *generate and test* programming pattern can be very inefficient (e.g., Palindrome program explores 10000 possibilities).

• An alternative is to use a *propagate and search* technique.

Propagate and search filters possibilities during the generation process, to prevent combinatorial explosion when possible.
Propagate and search approach is based on three key ideas:

• *Keep partial information*, e.g., “in any solution, X is greater than 100”.

• *Use local deduction*, e.g., combining “X is less than Y” and “X is greater than 100”, we can deduce “Y is greater than 101” (assuming Y is an integer.)

• *Do controlled search*. When no more deductions can be done, then search. Divide original CSP problem \( P \) into two new problems: \((P \land C)\) and \((P \land \neg C)\) and where \( C \) is a new constraint. The solution to \( P \) is the union of the two new sub-problems. Choice of \( C \) can significantly affect search space.
Find two digits that add to 10, multiply to more than 24:

\[ D_1::0\#9 \quad D_2::0\#9 \quad \% \text{ initial constraints} \]

\{Browse D1\} \quad \{Browse D2\} \quad \% \text{ partial results}

\[ D_1 + D_2 =: 10 \quad \% \text{ reduces search space from 100 to 81 possibilities} \]

\% D1 and D2 cannot be 0.

\[ D_1 \times D_2 \geq: 24 \quad \% \text{ reduces search space to 9 possibilities} \]

\% D1 and D2 must be between 4 and 6.

\[ D_1 <: D_2 \quad \% \text{ reduces search space to 4 possibilities} \]

\% D1 must be 4 or 5 and D2 must be 5 or 6.

\% It does not find unique solution D1=4 and D2=6
Propagate and Search Example(2)

- Find a rectangle whose perimeter is 20, whose area is greater than or equal to 24, and width less than height:

```plaintext
fun {Rectangle}
  W H in W::0#9  H::0#9
  W+H =: 10
  W*H >=: 24
  W <: H
  {FD.distribute naive rect(W H)}
  rect(W H)
end

{Browse {Search.base.all Rectangle}}
% displays [rect(4 6)]
```
Propagate and Search Example (Prolog)

- Find two digits that add to 10, multiply to more than 24:

```prolog
:- use_module(library(clpfd)).
q(D1,D2) :-
    D1 in 0..9, D2 in 0..9, % initial constraints
    D1+D2 #= 10, % D1 and D2 cannot be 0.
    D1*D2 #>= 24, % D1 and D2 must be between 4 and 6.
    D1 #< D2. % D1 must be 4 or 5 and
               % D2 must be 5 or 6.
    % It does not find unique solution D1=4 and D2=6.
```

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Propagate and Search Example (2)

- Find a rectangle whose perimeter is 20, whose area is greater than or equal to 24, and width less than height:

\[
\begin{align*}
\text{rectangle}(\left[ W,H \right]) : & - \\
& W \text{ in } 0..9, \ H \text{ in } 0..9, \\
& W+H \neq 10, \\
& W\times H \geq 24, \\
& W < H.
\end{align*}
\]

\[
\begin{align*}
\text{rectangleSolve}(\text{rect}(W,H)) : & - \\
& \text{rectangle}(\left[ W,H \right]), \\
& \text{label}(\left[ W,H \right]).
\end{align*}
\]

?- rectangleSolve(S).
S = rect(4, 6).
Finding palindromes (revisited)

- Find all four-digit palindromes that are products of two-digit numbers:

```lisp
fun {Palindrome}  
  A B C X Y in 
  A::1000#9999 B::0#99 C::0#99 
  A == B*C  
  X::1#9 Y::0#9 
  A == X*1000+Y*100+Y*10+X 
  {FD.distribute ff [X Y]}  
  A 
end 
{Browse {Search.base.all Palindrome}}  % 36 solutions 
```
Finding palindromes (revisited in Prolog)

- Find all four-digit palindromes that are products of two-digit numbers:

\[
\text{palindrome}(A,B,C,X,Y) :- \\
A \text{ in } 1000..9999, \ B \text{ in } 0..99, \ C \text{ in } 0..99, \\
A \neq B \times C, \\
X \text{ in } 1..9, \ Y \text{ in } 0..9, \\
A \neq X \times 1000 + Y \times 100 + Y \times 10 + X.
\]

\[
\text{palindromeSolve}(A) :- \\
\text{palindrome}(A,\_\_\_,X,Y), \\
\text{labeling([ff], [X,Y])}.
\]

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word(article,a).
word(article,every).
word(noun,criminal).
word(noun,'big kahuna burger').
word(verb,eats).
word(verb,likes).

sentence(Word1,Word2,Word3,Word4,Word5) :-
   word(article,Word1),
   word(noun,Word2),
   word(verb,Word3),
   word(article,Word4),
   word(noun,Word5).
fun {Word}
  choice
    article#a
    [] article#every
    [] noun#criminal
    [] noun#"big kahuna burger"
    [] verb#eats
    [] verb#likes
end
end

proc {Sentence S}
  Word1 Word2 Word3 Word4 Word5
in
  S = [Word1 Word2 Word3 Word4 Word5]
  {Word article#Word1}
  {Word noun#Word2}
  {Word verb#Word3}
  {Word article#Word4}
  {Word noun#Word5}
end
• *Definite Clause Grammars (DCG)* are useful for natural language parsing.

• Prolog can load DCG rules and convert them automatically to Prolog parsing rules.
DCG Syntax

-->  

DCG operator, e.g.,

sentence --> subject, verb, object.

Each goal is assumed to refer to the head of a DCG rule.

{prolog_code}

Include Prolog code in generated parser, e.g.,

subject --> modifier, noun, {write('subject')}.

[terminal_symbol]

Terminal symbols of the grammar, e.g.,

noun --> [cat].
Natural Language Parsing
(example rewritten using DCG)

sentence --> article, noun, verb, article, noun.

article --> [a] | [every].

noun --> [criminal] | ['big kahuna burger'].

verb --> [eats] | [likes].
Let us look at Prolog’s generated Horn clause for the `sentence` non-terminal:

```prolog
?- listing(sentence).
sentence(A, F) :-
    article(A, B),
    noun(B, C),
    verb(C, D),
    article(D, E),
    noun(E, F).
```

A-F is a *difference list*. B, C, D, and E are *accumulators*. Possible usage:

```prolog
?- sentence([a,criminal,likes,every,'big kahuna burger'],[]).
true
```
Natural Language Parsing (3)
(example rewritten using DCG)

Now, let us look at Prolog’s generated Horn clause for the verb non-terminal:

?- listing(verb).
verb(A, B) :-
    (   A=[eats|B]
    ;   A=[likes|B]
    ).

A-B is a difference list. Possible usage:

?- verb([likes],[]).
true.

?- verb([likes,cats],[cats]).
true.

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Let us look at an Oz relation for the sentence non-terminal:

\[
\text{proc } \text{Sentence } S \ Sn\{\text{Sentence } S \ Sn\}
    \text{S1 S2 S3 S4}
\text{in}
    \{\text{Article } S \ S1\}
    \{\text{Noun } S1 \ S2\}
    \{\text{Verb } S2 \ S3\}
    \{\text{Article } S3 \ S4\}
    \{\text{Noun } S4 \ Sn\}
\text{end}
\]

S-Sn is a \textit{difference list}. S1, S2, S3, and S4 are \textit{accumulators}. Possible usage:

\[
\text{proc } \text{Query } S\{\text{Sentence } S \ nil\} \text{ end}
\{\text{Browse } \{\text{Search.base.all Query}\}\}
\]
Now, let us look at Oz relation for the verb non-terminal:

\[
\text{proc \{Verb S Sn\}} \\
\text{ \hspace{1em} choice} \\
\text{ \hspace{2em} S = eats|Sn} \\
\text{ \hspace{2em} [] \hspace{0.5em} S = likes|Sn} \\
\text{ \hspace{1em} end} \\
\text{end}
\]

S-Sn is a \textit{difference list}. Possible usage:

\[
\{\text{Browse \{Search.base.all}} \\text{proc \{$ V\}} \\
\text{ \hspace{1em} \{Verb V nil\}} \\
\text{ \hspace{1em} end\}}}\]
Natural Language Parsing and Information Extraction

sentence(V) --> subject, verb(V), subject.
sentence(V) --> subject, verb(V).

subject --> article, noun.

article --> [a] | [every].

noun --> [criminal]
      | ['big kahuna burger']
      | [dog].

verb(eats) --> [eats].
verb(likes) --> [likes].
Prolog’s generated Horn clauses for the sentence non-terminal:

?- listing(sentence).

sentence(B, A, E) :-
    subject(A, C),
    verb(B, C, D),
    subject(D, E).

sentence(B, A, D) :-
    subject(A, C),
    verb(B, C, D).

A-E and A-D are difference lists. B is the extracted information (which could be a parse tree). C and D are accumulators. Possible usage:

?- sentence(Verb, [a,dog,eats],[[]]).
Verb = eats.
Now, let us look at Prolog’s generated Horn clause for the verb non-terminal:

?- listing(verb).
verb(eats, [eats|A], A).
verb(likes, [likes|A], A).

Possible usage:

?- verb(Verb, [eats], []).
Verb = eats.

?- verb(Verb, S, T).
Verb = eats,
S = [eats|T] ;
Verb = likes,
S = [likes|T].
Let us look at an Oz relation for the sentence non-terminal:

```
proc {Sentence V S Sn}
  S1 S2
in
  choice
    {Subject S S1}
    {Verb V S1 S2}
    {Subject S2 Sn}
  []
    {Subject S S1}
    {Verb V S1 Sn}
  end
end
```

S-Sn is a *difference list*. S1, and S2 are *accumulators*.

Possible usage:

```
fun {Query}
  V S
in
  {Sentence V S nil}
  V#S
end
{Browse {Search.base.all Query}}
```
Now, let us look at Oz relation for the verb non-terminal:

```oz
proc {Verb V S Sn}
  choice
  V = eats
  S = eats|Sn
  V = likes
  S = likes|Sn
end
end
```

S-Sn is a *difference list*. V is the extracted verb. Possible usage:

```oz
{Browse {Search.base.all
  proc {$ V}
    {Verb V _ _}
  end}}
```

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83. How would you translate DCG rules into Prolog/Oz rules?
84. PLP Exercise 11.8 (pg 571).
85. PLP Exercise 11.14 (pg 572).
86. CTM Exercise 12.6.2 (pg 774).