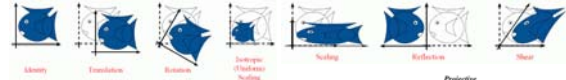


# Adjacency Data Structures

material from Justin Legakis

## Last Time?

- Simple Transformations



- Classes of Transformations

- Representation

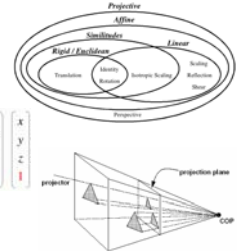
- homogeneous coordinates

- Composition

- not commutative

- Orthographic & Perspective Projections

$$\begin{bmatrix} x' \\ y' \\ z' \\ 1 \end{bmatrix} = \begin{bmatrix} a & b & c & d \\ e & f & g & h \\ i & j & k & l \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x \\ y \\ z \\ 1 \end{bmatrix}$$



## Today

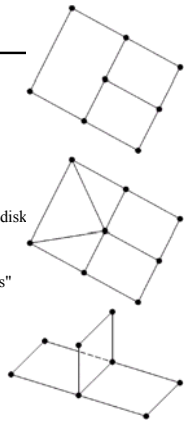
- Surface Definitions
- Simple Data Structures
- Fixed Storage Data Structures
- Fixed Computation Data Structures
- Mesh Simplification

## Today

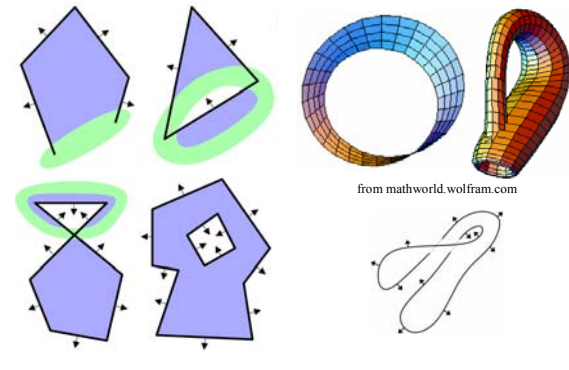
- **Surface Definitions**
  - Well-Formed Surfaces
  - Orientable Surfaces
  - Computational Complexity
- Simple Data Structures
- Fixed Storage Data Structures
- Fixed Computation Data Structures
- Mesh Simplification

## Well-Formed Surfaces

- Components Intersect "Properly"
  - Faces are: disjoint, share single Vertex, or share 2 Vertices and the Edge joining them
  - Every edge is incident to exactly 2 vertices
  - Every edge is incident to exactly 2 faces
- Local Topology is "Proper"
  - Neighborhood of a vertex is homeomorphic to a disk (permits stretching and bending, but not tearing)
  - Called a 2-manifold
  - Boundaries: half-disk, "manifold with boundaries"
- Global Topology is "Proper"
  - Connected
  - Closed
  - Bounded

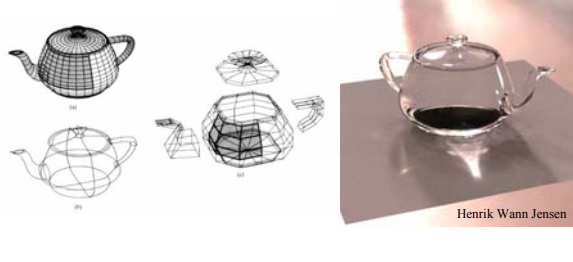


## Orientable Surfaces?



## Closed Surfaces and Refraction

- Original Teapot model is not "watertight":  
intersecting surfaces at spout & handle, no bottom,  
a hole at the spout tip, a gap between lid & base
- Requires repair before ray tracing with refraction



## Computational Complexity

- Adjacent Element Access Time
  - linear, constant time average case, or constant time?
  - requires loops/recursion/if?
- Memory
  - variable size arrays or constant size?
- Maintenance
  - ease of editing
  - ensuring consistency

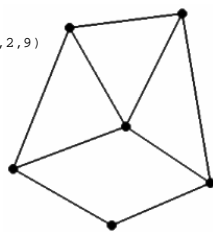
## Questions?

## Today

- Surface Definitions
- Simple Data Structures
  - List of Polygons
  - List of Edges
  - List of Unique Vertices & Indexed Faces:
  - Simple Adjacency Data Structure
- Fixed Storage Data Structures
- Fixed Computation Data Structures
- Mesh Simplification

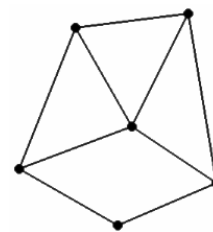
## List of Polygons:

(3, -2, 5), (3, 6, 2), (-6, 2, 4)  
 (2, 2, 4), (0, -1, -2), (9, 4, 0), (4, 2, 9)  
 (1, 2, -2), (8, 8, 7), (-4, -5, 1)  
 (-8, 2, 7), (-2, 3, 9), (1, 2, -7)



## List of Edges:

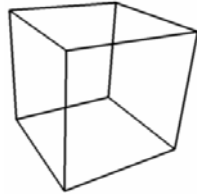
(3, 6, 2), (-6, 2, 4)  
 (2, 2, 4), (0, -1, -2)  
 (9, 4, 0), (4, 2, 9)  
 (8, 8, 7), (-4, -5, 1)  
 (-8, 2, 7), (1, 2, -7)  
 (3, 0, -3), (-7, 4, -3)  
 (9, 4, 0), (4, 2, 9)  
 (3, 6, 2), (-6, 2, 4)  
 (-3, 0, -4), (7, -3, -4)



## List of Unique Vertices & Indexed Faces:

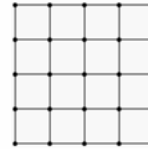
Vertices: (-1, -1, -1)  
(-1, -1, 1)  
(-1, 1, -1)  
(-1, 1, 1)  
(1, -1, -1)  
(1, -1, 1)  
(1, 1, -1)  
(1, 1, 1)

Faces: 1 2 4 3  
5 7 8 6  
1 5 6 2  
3 4 8 7  
1 3 7 5  
2 6 8 4

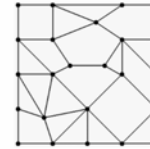


## Problems with Simple Data Structures

- No Adjacency Information
- Linear-time Searches



Structured



Unstructured

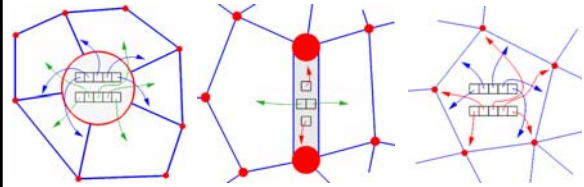
- Adjacency is implicit for structured meshes, but what do we do for unstructured meshes?

## Mesh Data

- So, in addition to:
  - Geometric Information (position)
  - Attribute Information (color, texture, temperature, population density, etc.)
- **Let's store:**
  - Topological Information (adjacency, connectivity)

## Simple Adjacency

- Each element (vertex, edge, and face) has a list of pointers to all incident elements
- Queries depend only on local complexity of mesh
- Data structures do not have fixed size
- Slow! Big! Too much work to maintain!



## Questions?

## Today

- Surface Definitions
- Simple Data Structures
- **Fixed Storage Data Structures**
  - **Winged Edge (Baumgart, 1975)**
- Fixed Computation Data Structures
- Mesh Simplification

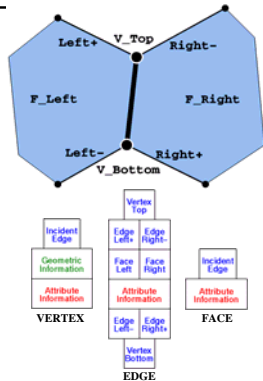
## Winged Edge (Baumgart, 1975)

- Each edge stores pointers to 4 Adjacent Edges
- Vertices and Faces have a single pointer to one incident Edge
- Data Structure Size?

Fixed

- How do we gather all faces surrounding one vertex?

Messy, because there is no consistent way to order pointers

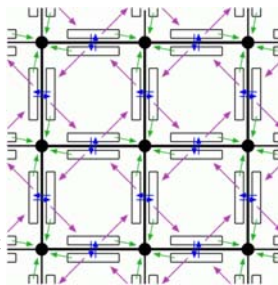


## Today

- Surface Definitions
- Simple Data Structures
- Fixed Storage Data Structures
- Fixed Computation Data Structures
  - HalfEdge (Eastman, 1982)
  - SplitEdge
  - Corner
  - QuadEdge (Guibas and Stolfi, 1985)
  - FacetEdge (Dobkin and Laszlo, 1987)

## HalfEdge (Eastman, 1982)

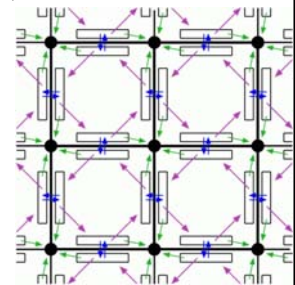
- Every edge is represented by two directed HalfEdge structures
- Each HalfEdge stores:
  - vertex at end of directed edge
  - symmetric half edge
  - face to left of edge
  - next points to the HalfEdge counter-clockwise around face on left
- Orientation is essential, but can be done consistently!



## HalfEdge (Eastman, 1982)

- Starting at half edge HE, how do we find:

- the other vertex of the edge?
- the other face of the edge?
- the clockwise edge around the face at the left?
- all the edges surrounding the face at the left?
- all the faces surrounding the vertex?



## HalfEdge (Eastman, 1982)

- Loop around a Face:
 

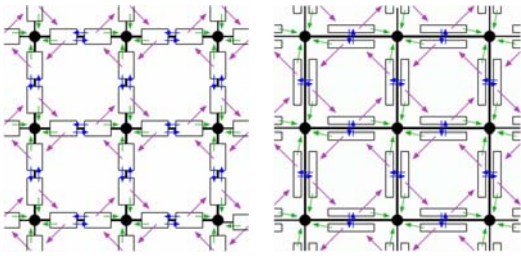
```
HalfEdgeMesh::FaceLoop(HalfEdge *HE) {
    HalfEdge *loop = HE;
    do {
        loop = loop->Next;
    } while (loop != HE);
}
```
- Loop around a Vertex:
 

```
HalfEdgeMesh::VertexLoop(HalfEdge *HE) {
    HalfEdge *loop = HE;
    do {
        loop = loop->Next->Sym;
    } while (loop != HE);
}
```

## HalfEdge (Eastman, 1982)

- Data Structure Size?
  - Fixed
- Data:
  - geometric information stored at Vertices
  - attribute information in Vertices, HalfEdges, and/or Faces
  - topological information in HalfEdges only!
- Orientable surfaces only (no Mobius Strips!)
- Local consistency everywhere implies global consistency
- Time Complexity?
  - linear in the amount of information gathered

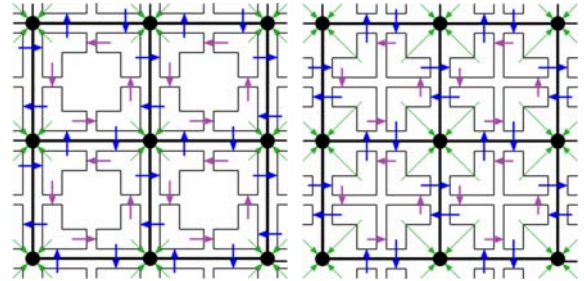
## SplitEdge Data Structure:



- HalfEdge and SplitEdge are dual structures!
- ```
SplitEdgeMesh::FaceLoop() = HalfEdgeMesh::VertexLoop()
SplitEdgeMesh::VertexLoop() = HalfEdgeMesh::FaceLoop()
```

## Corner Data Structure:

- The Corner data structure is its own dual!



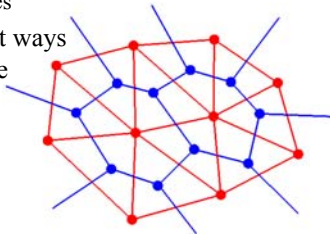
## Questions?

## Today

- Surface Definitions
- Simple Data Structures
- Fixed Storage Data Structures
- Fixed Computation Data Structures
  - HalfEdge (Eastman, 1982)
  - SplitEdge
  - Corner
  - QuadEdge (Guibas and Stolfi, 1985)
  - FacetEdge (Dobkin and Laszlo, 1987)
- Mesh Simplification

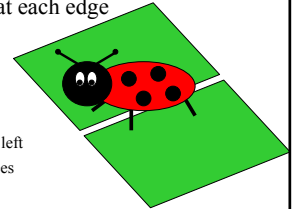
## QuadEdge (Guibas and Stolfi, 1985)

- Consider the Mesh and its *Dual* simultaneously
  - Vertices and Faces switch roles, we just re-label them
  - Edges remain Edges
- Now there are eight ways to look at each edge
  - Four ways to look at primal edge
  - Four ways to look at dual edge



## QuadEdge (Guibas and Stolfi, 1985)

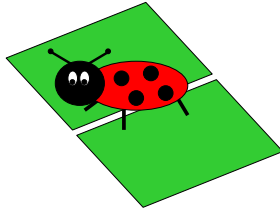
- Relations Between Edges: Edge Algebra
- Elements in Edge Algebra:
  - Each of 8 ways to look at each edge
- Operators in Edge Algebra:
  - Rot: Bug rotates 90 degrees to its left
  - Sym: Bug turns around 180 degrees
  - Flip: Bug flips up-side down
  - Onext: Bug rotates CCW about its origin (either Vertex or Face)



## QuadEdge (Guibas and Stolfi, 1985)

- Some Properties of Flip, Sym, Rot, and Onext:

- $e \text{ Rot}^1 = e$
- $e \text{ Rot}^2 \neq e$
- $e \text{ Flip}^2 = e$
- $e \text{ Flip Rot Flip Rot} = e$
- $e \text{ Rot Flip Rot Flip} = e$
- $e \text{ Rot Onext Rot Onext} = e$
- $e \text{ Flip Onext Flip Onext} = e$
- $e \text{ Flip}^{-1} = e \text{ Flip}$
- $e \text{ Sym} = e \text{ Rot}^2$
- $e \text{ Rot}^{-1} = e \text{ Rot}^3$
- $e \text{ Rot}^{-1} = e \text{ Flip Rot Flip}$
- $e \text{ Onext}^{-1} = e \text{ Rot Onext Rot}$
- $e \text{ Onext}^{-1} = e \text{ Flip Onext Flip}$

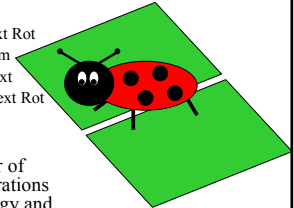


## QuadEdge (Guibas and Stolfi, 1985)

- Other Useful Definitions:

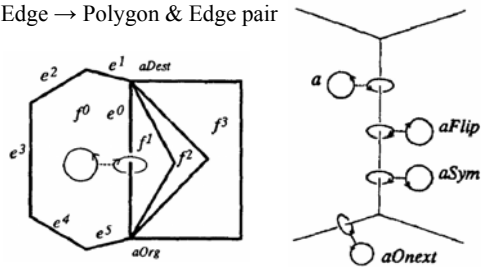
- $e \text{ Lnext} = e \text{ Rot}^{-1} \text{ Onext Rot}$
- $e \text{ Rnext} = e \text{ Rot Onext Rot}^{-1}$
- $e \text{ Dnext} = e \text{ Sym Onext Sym}^{-1}$
- $e \text{ Oprev} = e \text{ Onext}^{-1} = e \text{ Rot Onext Rot}$
- $e \text{ Lprev} = e \text{ Lnext}^{-1} = e \text{ Onext Sym}$
- $e \text{ Rprev} = e \text{ Rnext}^{-1} = e \text{ Sym Onext}$
- $e \text{ Dprev} = e \text{ Dnext}^{-1} = e \text{ Rot}^{-1} \text{ Onext Rot}$

- All of these functions can be expressed as a constant number of Rot, Sym, Flip, and Onext operations independent of the local topology and the global size and complexity of the mesh.



## FacetEdge (Dobkin and Laszlo, 1987)

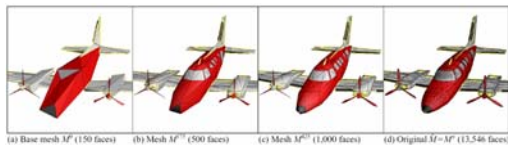
- QuadEdge (2D, surface)  $\rightarrow$  FacetEdge (3D, volume)
- Faces  $\rightarrow$  Polyhedra / Cells
- Edge  $\rightarrow$  Polygon & Edge pair



## Questions?

## Today

- Surface Definitions
- Simple Data Structures
- Fixed Storage Data Structures
- Fixed Computation Data Structures
- **Mesh Simplification**



Hugues Hoppe "Progressive Meshes" SIGGRAPH 1996

## Progressive Meshes

- Mesh Simplification
  - vertex split / edge collapse
  - geometry & discrete/scalar attributes
  - priority queue
- Level of Detail
  - geomorphs
- Progressive Transmission
- Mesh Compression
- Selective Refinement
  - view dependent

## Selective Refinement

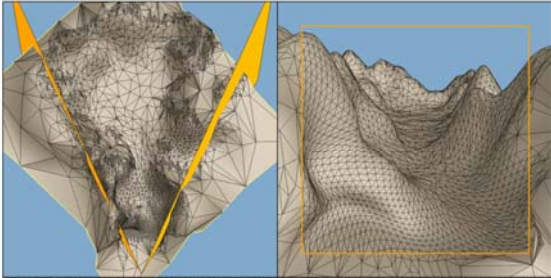


Figure 10: Selective refinement of a terrain mesh taking into account view frustum, silhouette regions, and projected screen size of faces (7,438 faces).

## Preserving Discontinuity Curves

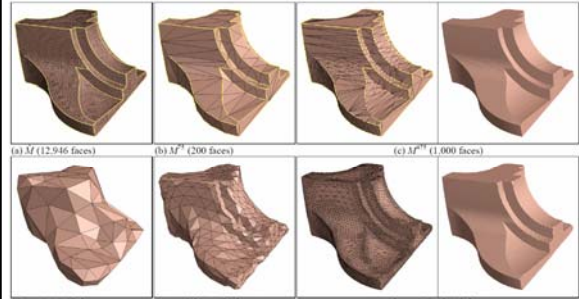


Figure 12: Approximations of a mesh  $M$  using (b-c) the PM representations, and (d-f) the MRA scheme of Eck et al. [7]. As demonstrated, MRA cannot recover  $M$  exactly, cannot deal effectively with surface creases, and produces approximating meshes of inferior quality.

## Other Simplification Strategies

- Remove a vertex & surrounding triangles, re-triangulate the hole



- Merge Nearby Vertices
  - will likely change the topology...

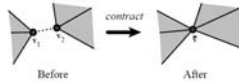


Figure 2: Non-edge contraction. When non-edge pairs are contracted, unconnected sections of the model are joined. The dashed line indicates the two vertices being contracted together.  
from Garland & Heckbert, "Surface Simplification Using Quadric Error Metrics" SIGGRAPH 1997

## When to Preserve Topology?



Figure 3: On the left is a regular grid of 100 closely spaced cubes. In the middle, an approximation built using only edge contractions demonstrates unacceptable fragmentation. On the right, the result of using more general pair contractions to achieve aggregation is an approximation much closer to the original.

from Garland & Heckbert, "Surface Simplification Using Quadric Error Metrics" SIGGRAPH 1997

## Quadric Error Simplification

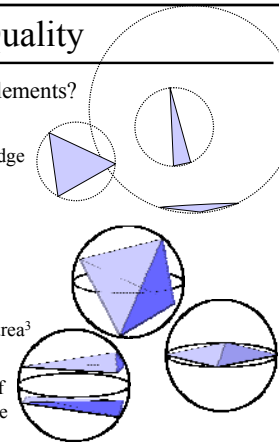
- Contract (merge) vertices  $v_i$  and  $v_j$  if:
  - $(v_i, v_j)$  is an edge, or
  - $\|v_i - v_j\| < t$ , where  $t$  is a threshold parameter
- Track cumulative error by summing 4x4 quadric error matrices after each operation:

$$\begin{aligned} \Delta(v) &= \sum_{p \in \text{triangle}(v)} (v^T p)(p^T v) \\ &= \sum_{p \in \text{triangle}(v)} v^T (pp^T) v \\ &= v^T \left( \sum_{p \in \text{triangle}(v)} K_p \right) v \\ K_p = pp^T &= \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix} \end{aligned}$$

Garland & Heckbert, "Surface Simplification Using Quadric Error Metrics" SIGGRAPH 1997

## Judging Element Quality

- How "equilateral" are the elements?
- For Triangles?
  - Ratio of shortest to longest edge
  - Ratio of area to perimeter<sup>2</sup>
  - Smallest angle
  - Ratio of area to area of smallest circumscribed circle
- For Tetrahedra?
  - Ratio of volume<sup>2</sup> to surface area<sup>3</sup>
  - Smallest *solid* angle
  - Ratio of volume to volume of smallest circumscribed sphere

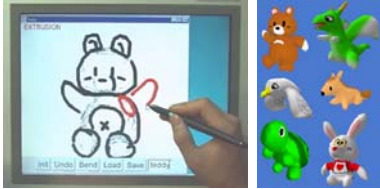




## Reading for Tuesday (1/22)

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- "Teddy: A Sketching Interface for 3D Freeform Design", Igarashi et al., SIGGRAPH 1999



- Post a comment or question on the LMS discussion by 10am on Tuesday 1/22