

# **FINAL: 180 Minutes**

Last Name: \_\_\_\_\_

First Name: \_\_\_\_\_

RIN: \_\_\_\_\_

Section: \_\_\_\_\_

Answer **ALL** questions.

**NO COLLABORATION** or electronic devices. Any violations result in an F.

**NO questions** allowed during the test. Interpret and do the best you can.

## **GOOD LUCK!**

You **MUST** show **CORRECT** work to get full credit.

**When in doubt, TINKER.**

<b>1-12</b>	<b>Total</b>
<b>10 each</b>	<b>120</b>

- 1 Determine the type of proof and prove: a right triangle with integer sides cannot be isosceles.

**2** Prove that the product of any 5 consecutive natural numbers is divisible by  $5!$ .

- 3** Let  $G_0 = 0$ ,  $G_1 = 1$ , and  $G_n = 7G_{n-1} - 12G_{n-2}$  for  $n > 1$ . Compute  $G_5$ . Show  $G_n = 4^n - 3^n$  for  $n \geq 0$ .

**4** Determine and prove the order-relationships between  $\ln n$ ,  $\ln(n^2 + 1)$ , and  $\ln(2n)$ .

**5**    Prove that  
  
implies that

$$xk \equiv yk \pmod{d}$$

$$x \equiv y \pmod{(d/\gcd(k, d))}.$$

- 6** Adam, Barb, Charlie, and Doris each independently choose a random number uniformly distributed in  $\{1, 2, 3, 4, 5\}$ . What is the probability that some pair chooses the same number? What if there are  $k$  people and  $n$  numbers?

- 7** Let  $\mathcal{V}$  be a set of  $n$  vertices, and let the edge set  $\mathcal{E}$  be initially empty. For each pair of vertices  $i \neq j$ , add the edge  $(i, j)$  to  $\mathcal{E}$  with probability  $p$ . Give the pdf for the degrees of the nodes of this graph.



- 8 Voltage in the US has a mean of 120V and a standard deviation of 5V. A device's operating voltage is 112–128. Use Chebyshev's inequality to bound the probability that the device will not be damaged when turned on.

- 9** Give a DFA for strings whose even digits alternate between 0 and 1.

**10** For CFG  $S \rightarrow 0S|S1|0|1$ , prove that no string has 10 as a substring.

- 11** Consider the language of palindromes  $\mathcal{L} = \{\omega \mid \omega \in \{0,1\}^*, \omega = \omega^R\}$ . Give well-written high-level pseudocode for a decider for this language.

- 12** Given an ultimate-debugger  $\mathcal{D}$  that takes  $\langle M \rangle \# \omega$  and decides if TM  $M$  halts on input  $\omega$ , show that every recognizer of a language  $\mathcal{L}$  can be converted into a decider for the language.

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