

Quiz 1

60 Minutes

First Name: SOLUTIONS

Last Name: _____

RIN: _____

Section: _____

NO COLLABORATION or electronic devices.

Any violations will result in an **F**.

No questions allowed during the test unless you think there is a mistake.

GOOD LUCK!

Circle at most one answer per question.

10 points for each correct answer.

You **MUST** show **CORRECT** work to get credit.

Correct answers with no explanation will get a 0.

Final Score: _____ / 200

1. What is a simpler expression for the set $A = \{n \mid n = k + m, k \in \mathbb{Z}, m \in \mathbb{N}\}$?

- A $A = \{n \mid n \in \mathbb{N}\}$
- B $A = \{n \mid n \in \mathbb{Z}\}$
- C $A = \{n \mid n = 1 + k, k \in \mathbb{N}\}$
- D $A = \{n \mid n = -1 + k, k \in \mathbb{N}\}$
- E None of the above.

Pick any $z \in \mathbb{Z}$.
 We know $z = \underbrace{(z-1)}_{\in \mathbb{Z}} + \underbrace{1}_{\in \mathbb{N}}$

2. Define sets $A = \{n \mid n = 2k, k \in \mathbb{Z}\}$ and $B = \{n \mid n \in \mathbb{N}\}$. Which of the following is true?

- A $A \subseteq B$ $-2 \notin B$
- B $B \subseteq A$ $1 \notin A$
- C $A = B$
- D $A \cap B = \{n \mid n = 2k, k \in \mathbb{N}\}$
- E None of the above.

3. What is the value of the expression $(p \wedge q) \rightarrow (\neg p \vee q)$?

- A Depends on the value of p .
- B Depends on the value of q .
- C Always true.
- D Always false.
- E None of the above.

We know
 $(p \Rightarrow q) \equiv \neg p \vee q$
 $(p \wedge q) \Rightarrow (\neg p \vee q) \equiv$
 $\equiv \neg p \vee \underbrace{\neg q}_{\vee \neg p} \vee \underbrace{q}_{\vee q} \equiv T$

4. IF it snows on a given day, THEN the temperature on that day is below 35. There is snow on the ground and the temperature today is -15. What do we know?

- A It is snowing today.
- B It snowed yesterday.
- C On the day when it snowed, the temperature was below 35.
- D It snows every day.
- E None of the above.

5. What is another way to write the expression $(p \wedge q) \rightarrow \neg p$?

- A $\neg p \wedge q$
- B $\neg p \wedge \neg q$
- C $\neg p \vee q$
- D $\neg p \vee \neg q$
- E None of the above.

$(p \wedge q) \Rightarrow \neg p \equiv \neg p \vee \neg q \vee \neg p$

6. Suppose I want to prove $p \rightarrow q$. Which of the following proof techniques will work?

- A Assume p and show that it leads to q being false.
- B Assume q is true and show that p must be true.
- C Assume q is false and show that p must be true.
- D Use a proof by contradiction to show that q can never be true.
- E None of the above.

| p | q | $p \Rightarrow q$ |
|-----|-----|-------------------|
| F | F | T |
| F | T | T |
| T | F | F |
| T | T | T |

Avoid

7. Which of the following claims is NOT true?

- A $\forall x \in \mathbb{N} : \exists y \in \mathbb{R} : \sqrt{x} = y$
- B $\forall x \in \mathbb{Z} : \exists y \in \mathbb{R} : (x \geq 0 \rightarrow \sqrt{x} = y)$
- C $\forall x \in \mathbb{Q} : \exists y \in \mathbb{R} : (x \geq 0 \rightarrow \sqrt{x} = y)$
- D $\forall x \in \mathbb{N} : \exists y \in \mathbb{N} : \sqrt{x} = y$
- E They are all true.

$2 \in \mathbb{N}, \sqrt{2} \notin \mathbb{N}$

8. Consider the set $\mathbb{Q}_+ = \{x \mid x > 0, x \in \mathbb{Q}\}$ of positive rational numbers. Suppose I use the following argument to prove that \mathbb{Q}_+ has a minimum element: since each rational number can be written as p/q where $q \in \mathbb{N}$, from the principle of well-ordering it follows that \mathbb{Q}_+ has a minimum element. What is wrong with this argument?

- A Not all elements of \mathbb{Q}_+ can be written as p/q .
- B The principle of well-ordering does not apply to infinite sets.
- C Although the principle of well-ordering can be used to conclude that there is a minimum denominator q , this does not imply that \mathbb{Q}_+ has a minimum.
- D The principle of well-ordering is an axiom and cannot be used in proofs.
- E Nothing is wrong: \mathbb{Q}_+ has a minimum element!

9. Suppose I try to prove $n^2 + 10 \geq 2^n, \forall n \geq 1$ using induction. What goes wrong?

- A The base case is false.
- B I need three base cases.
- C I cannot prove $P(n) \rightarrow P(n+1), \forall n \geq 1$.
- D I cannot prove $P(n+1) \rightarrow P(n), \forall n \geq 1$.
- E Nothing goes wrong because the claim is true.

Exponentials grow faster than polynomials.

10. How would you *disprove* the claim: $n^4 > 2^n, \forall n \geq 3$?

- A Show that $n^4 > 2^n$ for some $n < 3$.
- B Show that $n^4 > 2^n$ for all $n \geq 3$.
- C Show that $n^4 \leq 2^n$ for some $n < 3$.
- D Show that $n^4 \leq 2^n$ for some $n \geq 3$.
- E None of the above.

The negation of $\forall n \geq 3: n^4 > 2^n$ is $\exists n \geq 3: n^4 \leq 2^n$

$$S_{100} = \frac{1}{2} \cdot 100 \cdot 102 = 50 \times 51 = 2550$$

11. Let $S_n = 2 + 4 + \dots + n$ be the sum of all even numbers up to n . What is S_{100} ?

- A 50×99
- B 5050
- C 5100
- D 2550
- E None of the above.

Use Gauss's technique.

$$S_n = 2 + 4 + \dots + n-2 + n$$

$$S_n = n + n-2 + \dots + 2$$

$$2S_n = (n+2) \cdot \frac{n}{2} \rightarrow S_n = \frac{1}{2} \cdot n \cdot (n+2)$$

12. Consider the set S defined as follows: (1) Base case: $4 \in S$; (2) Constructor: $x \in S \rightarrow x^2 \in S$. Which of the following is another definition of the set S ?

- A $S = \{s \mid s = 2^n, n \in \mathbb{N}\}$
- B $S = \{s \mid s = 2^n, n \in \mathbb{Z}\}$
- C $S = \{s \mid s = 2^{2^n}, n \in \mathbb{N}\}$
- D $S = \{s \mid s = 2^{2^n}, n \in \mathbb{N}\}$
- E None of the above.

Base case: $4 = 2^2 = 2^{2^1}$

Tinker: $2^2 \in S \Rightarrow 2^4 \in S \Rightarrow 2^8 \in S$.

Inductive step: $P(n) : 2^{2^n} \in S$

$$(2^{2^n})^2 = 2^{2 \cdot 2^n} = 2^{2^{n+1}} \quad (P(n+1)). \quad \checkmark$$

13. Define the predicate $P(n) : n + 5 \leq n^2$. For which n is $P(n)$ true?

- A $n \geq 1$
- B $n \geq 2$
- C $n \geq 3$
- D Only for $3 \leq n \leq 300$
- E None of the above.

$$6 \not\leq 1$$

$$7 \not\leq 4$$

Base case: $P(3) : 8 \leq 9. \quad \checkmark$

Inductive step: $P(n) \Rightarrow P(n+1)$

$$n+5+1 \leq n^2 + 1 \leq n^2 + 2n + 1$$

for $n \geq 3$.

14. As part of a proof by induction, I showed $P(1)$ and $P(2)$ are true, and I also showed that for all $n \geq 3$, there exist $i, j < n$ such that $(P(i) \wedge P(j)) \rightarrow P(n)$. Is this sufficient proof that $\forall n \geq 1 : P(n)$?

- A No, because I didn't prove $P(n) \rightarrow P(n+1)$
- B No, because i and j may not always exist
- C No, because I didn't prove the base case
- D Yes, it follows from strong induction. Define $Q(n) = P(1) \wedge \dots \wedge P(n)$. I essentially showed $Q(n) \rightarrow Q(n+1)$ since $P(i)$ and $P(j)$ are part of the induction hypothesis.
- E None of the above.

15. What can we say about this claim: $\forall x \in \mathbb{R} : \exists y \in \mathbb{R} : xy = y$?

- A True
- B False
- C Depends on the value of x
- D Depends on the value of y
- E None of the above.

$\forall x \in \mathbb{R}$, pick $y = 0$.

Then $0x = 0$.

16. What is another expression for the set $A \cap (A \cup B)$?

- A $A \cap B$
- B $A \cup B$
- C A
- D $A \cap \bar{B}$
- E None of the above.

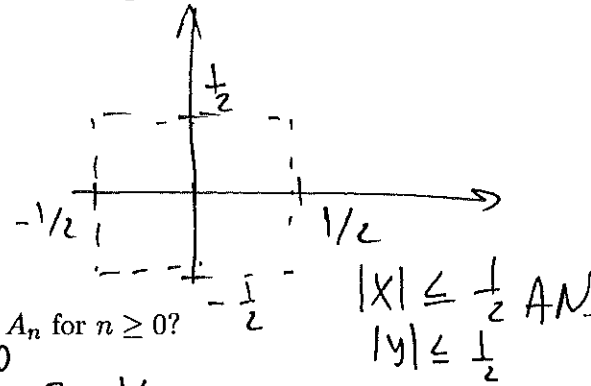
subset of

$$A \cap (A \cup B) = \overbrace{A \cap B}^A \cup \underbrace{A \cap A}_{A}$$

Union of A and a subset of A is just A .

17. What is a formal way to define a square of side 1 centered at the origin of \mathbb{R}^2 ?

- A $S = \{(x, y) \mid x \leq 1, y \leq 1, (x, y) \in \mathbb{R}^2\}$
- B $S = \{(x, y) \mid |x| \leq 1, |y| \leq 1, (x, y) \in \mathbb{R}^2\}$
- C $S = \{(x, y) \mid x \leq 1/2, y \leq 1/2, (x, y) \in \mathbb{R}^2\}$
- D $S = \{(x, y) \mid |x| \leq 1/2, |y| \leq 1/2, (x, y) \in \mathbb{R}^2\}$
- E None of the above.



18. Let $A_n = A_{n-1}^2$ and $A_0 = 3$. What is a general formula for A_n for $n \geq 0$?

- A $A_n = 2^{3^n}$
- B $A_n = 3^{2^n}$
- C $A_n = 2^{2^n}$
- D $A_n = 3^{2^n}$
- E None of the above.

Base case: $P(0): 3^{2^0} = 3$. \checkmark

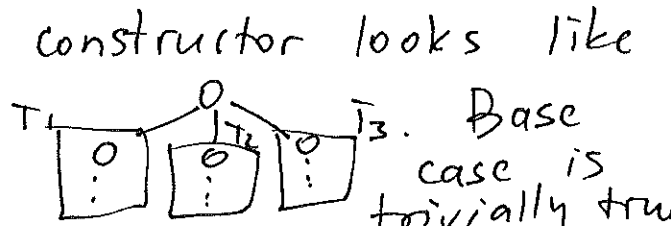
Inductive step: $P(n): A_n = 3^{2^n}$
 $A_{n+1} = (3^{2^n})^2 = 3^{2 \cdot 2^n} = 3^{2^{n+1}} = P(n+1)$

19. Suppose I create a new type of rooted tree, called rooted tertiary tree (RTT), as follows:

- ① The tree with one vertex is an RTT. [base case]
- ② If T_1, T_2 and T_3 are RTTs with roots r_1, r_2 and r_3 , then linking r_1, r_2 and r_3 to a new root r gives a new RTT with root r . [constructor]
- ③ No other tree is an RTT. [minimality]

What do we know about RTTs (recall RBT stands for a rooted binary tree)? No empty trees. So

- A If a vertex has children, then that vertex has three children.
- B All RBTs are also RTTs.
- C All RTTs have an odd number of vertices.
- D All RTTs have an even number of vertices.
- E All of the above.



20. Suppose a rooted tertiary tree has 7 vertices. How many links does the tree have?

- A 6
- B 7
- C 8
- D 9
- E It cannot be determined.

#links = #vertices - 1. Base case: 1 vertex, 0 links

Inductive step: add 3 links and 1 vertex: $L = (n_1 - 1) + (n_2 - 1) + (n_3 - 1) + 3 = n - 1 - 3 + 3 = n - 1$

SHOW WORK

SHOW WORK

Scratch