

Possible proofs (of many)

For all, we first consider w.l.o.g. a component of G where $|E| \geq |V|$ holds

→ one must obviously exist, as all cannot have $|E| < |V|$ while G has $|E(G)| \geq |V(G)|$

We can also assume G is simple

→ self loops and multi-edges are trivially cyclic

①

Contrapositive

G has no cycles $\Rightarrow |V(G)| > |E(G)|$

Last WP we showed

{connected acyclic simple} $\Rightarrow |V| - 1 = |E|$
(i.e., is a tree)

Note: trees are maximally acyclic

→ adding an edge creates a cycle
since all vertices are connected

\Rightarrow so any acyclic graph must have at most as many edges as a tree

$$\Rightarrow |E| \leq |V| - 1 \quad \square$$
$$|E| < |V|$$

② Strong induction  arm

Basis: $C_n \rightarrow |E| = |V| \rightarrow |E| \geq |V|$ ✓
(Note: if we use graph class \mathcal{G} , we need to show it holds $\forall G \in \mathcal{G}$)

Assume we have $P(n)$ with n edges
and $n \geq |V(P(n))|$

Note: $\exists v \in V(P(n)) : d(v) = 1$

\rightarrow otherwise, the minimum degree of G is 2 and it must have a cycle
(proved in class via **Extremal argument**)

$$P(k) = P(n) - v$$

\rightarrow we delete exactly one vertex
and edge, so $|E(P(k))| = |E(P(n))| - 1$
 $|V(P(k))| = |V(P(n))| - 1$

and the inequality still holds

I.H. $\rightarrow P(k)$ has a cycle

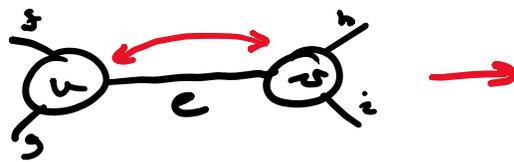
I.H. $\rightarrow P(k)$ has a cycle

$P(k) + v \rightarrow P(n)$

\rightarrow Adding a vertex does not
remove a cycle

$\Rightarrow P(n)$ is cyclic \square

③ Can do same proof above using
edge contractions as construction



Basis could
just be
 $P(1) = \emptyset = C_1$

④ Can do weak inductive proof on the
contra-positive 

Basis $P(0)$: \emptyset \leftarrow single vert $|V| > |E|$

I.H.: Assume $P(k)$ is acyclic and $|V| > |E|$

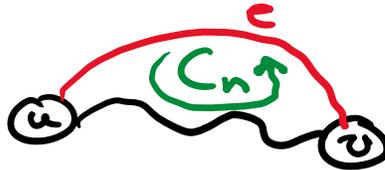
$P(k+1) = P(k) + e$

(Note our above assumption about
the graph being connected)

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the graph being connected)

Consider the cases

1. e connects to 2 new vertices \times
 \rightarrow this makes $P(k+1)$ disconnected
2. e connects to 2 existing vertices \times
 \rightarrow this creates a cycle, as those
2 vertices are already connected



3. e must connect a new degree-1
(leaf) vertex to an existing vertex

$$\rightarrow |V(P(k+1))| = |V(P(k))| + 1$$

$$|E(P(k+1))| = |E(P(k))| + 1$$

\Rightarrow and the inequality holds \square