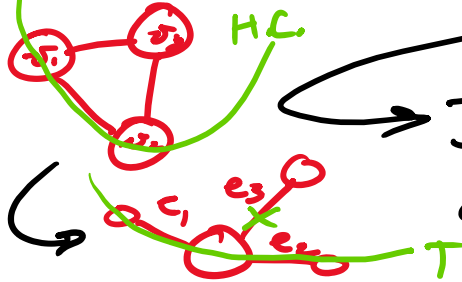


( $\Rightarrow$ ) Consider the H.C. on  $L(G)$  as a traversal of  $V(L(G)) \rightarrow \{v_0 v_1 \dots v_n v_0\}$   
 $\rightarrow$  These vertices are the edges of  $G$  per the definition of line graphs

SO  $\{v_0 v_1 \dots v_n v_0\} \rightarrow \{e_0 e_1 \dots e_n e_0\}$



$\rightarrow$  If it is possible to traverse edges in this order  $\rightarrow$  we are done

Otherwise: we have some  $e_i$  incident on a vertex shared with  $e_{i-1}$  and  $e_{i+1}$   
 $\rightarrow$  we remove  $e_i$  but it is still covered

$\Rightarrow$  we have a closed tour that is trivially a cover  $\checkmark$

( $\Leftarrow$ ) We construct a H.C. on  $L(G)$   
ALGORITHMICALLY

Consider our closed trail on  $G$

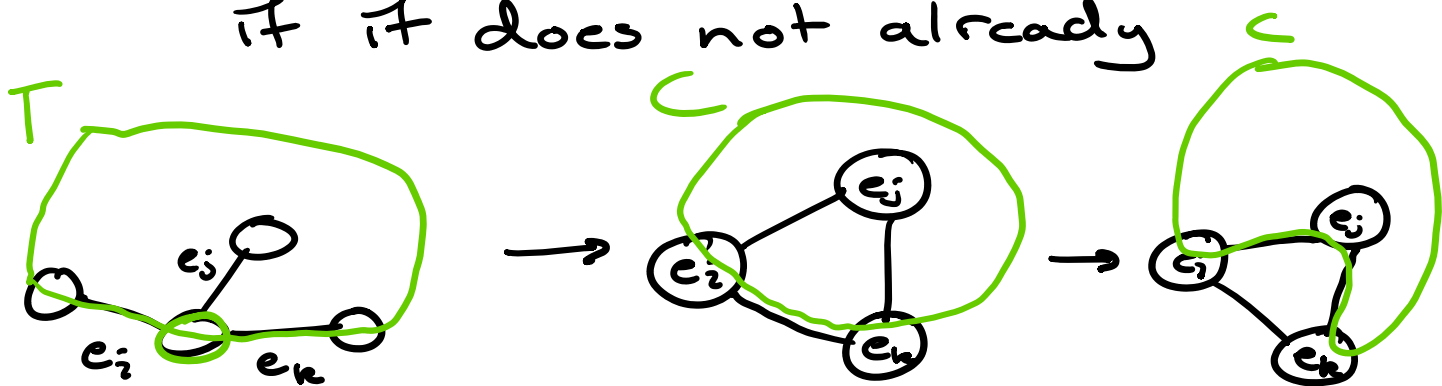
$$\Gamma = \{e_0 e_1 \dots e_i e_0\}$$

$$T = \{e_0, e_1, \dots, e_{i-1}, e_{i+1}, \dots, e_{k-1}, e_k\}$$

Note: this does not contain all edges of  $G$  (necessarity)

BUT: it contains a set of vertices that are incident to all edges

Directly translate  $T$  to a cycle on  $L(G)$ , and consider expanding the cycle to contain all  $V(L(G))$  if it does not already



As any  $e_j$  not in  $T$  is incident on a shared vertex in  $V(T)$ , it is part of a clique with some  $e_i, e_k$  in  $L(G)$

→ we can trivially add  $e_j$  to  $C$   
via the edges in that clique

→ if there are multiple  $e_j$  on  
the same vertex, we can also  
iterate this process

⇒ we will eventually add all  $e_j$   
to  $C$ , and since  $V(T)$  is a cover,  
 $C$  will include all of  $V(G)$   
and therefore be a H.C.  $\square$